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THESIS

DESIGN AND IMPLEMENTATION OF A BASIC
CROSS-COMPILER AND VIRTUAL MEMORY MANAGEMENT
SYSTEM FOR THE TI-59 PROGRAMMABLE CALCULATOR

by

Mark R. Kindl
and
James H. W. Inskeep, Jr.
June, 1983

Thesis Advisor: Bruce MacLennan

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for the TI-59 Programmable Calculator

by

Mark R. Kindl
Captain, United States Army
E.S., United States Military Academy, 1974

and

James H. W. Inskape, Jr.
Captain, United States Army
E.S., United States Military Academy, 1974

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June 1983

Authors:

Mark R. Kindl

James H. W. Inskape, Jr.

Approved by:

James H. W. Inskape, Jr.

Thesis Advisor

Q. E. Latta

Second Reader

David K. Hsieh

Chairman, Department of Computer Science

Kenneth T. Marshall

Dean of Information and Policy Sciences

ABSTRACT

The instruction set of the TI-59 Programmable Calculator bears a close similarity to that of an assembler. Though most of the calculator instructions perform primitive data movement and/or sequence control, some can do the work of small high level language procedures. Regardless of this fact, to design and debug TI-59 programs of moderate size can be more difficult than doing the computations themselves. Programming in a higher order language such as BASIC offers many advantages over calculator code. This report presents the design and implementation of a cross-compiler which translates correct BASIC programs into equivalent TI-59 programs. This software package includes a linker which maps calculator instructions to a set of magnetic cards. The cards are then used to implement a manually operated virtual memory system for the calculator. This expands program step capacity, and permits more complex programs to be written in BASIC language for translation into TI-59 instructions.

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I. INTRODUCTION

Hand-held programmable calculators provide an extremely portable means of computation. Designed primarily for small-scale numerical computation, these devices are limited by their small memory capacity, slow processing speed, and inability to perform symbol manipulation. These constraints, however, cannot hide the programmable calculator's usefulness and power as a computational tool. The instruction sets of these machines resemble those of assemblers. Most instruction words are primitive and perform simple data movements or sequence control. Yet, some specialized instructions do the work of small high order language procedures. Typical examples include polar to rectangular coordinate conversion, trigonometric functions, and logarithmic functions. In most assembly languages these operations must be constructed from primitive instructions. Even with these added features, the designing and debugging of calculator programs for non-trivial problems often requires an expenditure of effort which conceals the usefulness of the final product as well as the calculator.

There are a number of reasons for this difficulty. The lack of a sophisticated display mechanism such as a CRT prevents the user from viewing more than a single data item or instruction at any given time. Printing devices can lend assistance, yet most provide little more than a means of dumping memory contents. Furthermore, program debugging can be very difficult if the only error diagnostic is a single flashing display or an incorrect result.

Maximum memory storage capacity constrains even main-frame computers. One solution has been the implementation of virtual memory, whereby a relatively limitless on-call

secondary store is used to back up the primary storage. Programmable calculators usually have secondary storage in the form of magnetic cards. Normally, these cards are used as archival rather than on-call storage. The instruction sets of calculators are generally a cross between assembler instructions and a math function library. Compared to an assembler, the calculator instruction set is small and includes only the most basic sequence controls. Though it is possible to build more sophisticated control constructs from the primitives, such endeavors are often constrained by storage capacity. As a result, if complex programs are to fit into memory, it becomes necessary to learn or invent machine dependent "tricks."

The programmer's inability to use meaningful names for variables can create more difficulties during calculator programming. Numerical register indices (a form of absolute addressing) must be used to reference variables. One of the major advantages of assemblers is that they provide for variable naming. While composing his code, the calculator programmer must either remember the register indices of his variables or continuously refer to his own written symbol table while composing code. Both methods become more error prone as the number of variables in use grows. Programs of any substantial computing power usually require large numbers of variables.

The problems associated with calculator programming are many of the same problems which plagued experienced programmers of large scale computers in the past. How can a beginner be expected to design good, sophisticated programs for a pocket calculator if it can be so difficult for someone with experience? One concludes that the majority of users write small, relatively straight forward programs and never fully utilize the power of such calculators.

The Texas Instruments TI-59 Programmable Calculator is one of the more popular models. Its value as a powerful engineering tool is indicated by its use at the U.S. Naval Postgraduate School and in the U.S. Army Field Artillery. Yet, sophisticated programming of the TI-59 suffers from the very weaknesses mentioned earlier. Why, then, should there be such interest in this device? Perhaps, the best answer to this question is provided by Hamming [Ref. 1]. He feels that such a primitive programming machine offers the user valuable experience because it is easy to operate and allows the beginner direct access to very basic computing hardware. But, he warns that attempting mastery of the TI-59 language is a waste of time. One who must do sophisticated or extensive programming for this calculator should, instead, use a cross-compiler to automate and reduce his effort. This report presents the design and implementation of one such cross-compiler for the TI-59.

II. SOFTWARE REQUIREMENTS

A systematic approach to software development begins with the defining of general requirements. In this case, the basic design goal is the production of an effective software tool which will simplify program development and increase memory capacity for the TI-59 Programmable Calculator. Achievement of this goal should result in several enhancements to the utility and capability of the calculator. There will be an increase in its ability to execute larger and more sophisticated software. Most computations which can be programmed in BASIC and some existing BASIC software (which may require minor modifications to be explained later) will become as portable as the TI-59.

Important requirements for a user-oriented language translation system should include that it be easy to use and easy to learn. The BASIC programming language is an obvious choice for the source language; it is popular, simple, and easily learned. More importantly, many BASIC constructs and key words are similar to those of TI-59. This similarity and the fact that both languages are line-oriented and sequential in nature greatly facilitates translation between them.

Many versions of the BASIC language currently exist. Because of its availability at the Naval Postgraduate School and its having many structured control flow constructs, Waterloo BASIC Version 2.0 (WBASIC) [Ref. 2], was chosen as the specific source language to be implemented by this compiler. The power of the TI-59 compared to the WBASIC language places restrictions upon the set of WBASIC commands which can be translated. The specific WBASIC subset implemented is deferred to the discussion of design issues in the

next chapter. While WBASIC is easy to learn, it should be apparent that subsetting the language will introduce exceptions and restrictions which will tend to complicate learning for the novice and confuse the veteran. It is desirable to maintain as few exceptions as possible, and to require that restrictions be clean and obvious. A construct should be implemented as completely as possible (within obvious limitations, such as the file handling or alphanumeric features) or not at all.

Provision for error detection and debugging is another important requirement of a language system. The intended use of this initial system will be as a supplement to an existing WBASIC interpreter or compiler. As such, the cross-compiler will assume error-free source program input. The only requirement for error detection will be for the compiler to recognize words/constructs which are not implemented, but which are ordinarily legal WBASIC commands. Debugging of TI-59 programs is so much more difficult than debugging of higher level language programs that it is reasonable to assume that a user would prefer to debug his WBASIC program using the WBASIC interpreter/compiler available. Once the program is logically correct, it may be cross-compiled to TI-59 code, at which time it will be checked for subset and calculator capacity errors.

The TI-59's designers have provided it with capabilities which can be roughly equated to the power of higher level language routines. Interchangeable Solid State Software (trademark of Texas Instruments, Inc.) modules allow on-line access to utility program libraries. Program steps required to call them and the exclusive reservation of particular registers are usually the only storage costs paid for use of these library programs. It is required that the power of these modules be harnessed by the translator under design. Additionally, other sophisticated features of the calculator

should be exploited whenever possible in order to maximize and enhance advantages gained by high level language programming.

The linker will statically link the steps of TI-59 programs so that it will not be necessary for a complete program to reside in calculator step storage during execution. Since the swapping in and out of memory modules in the form of magnetic cards can become quite complicated for a running process, it will be necessary to keep this manual system as transparent to the user as can be reasonably expected. The fact that the calculator has a single item display window and associated register will certainly restrict the degree of transparency which might otherwise be possible.

A system must perform static linking if it exceeds program storage available to it during execution. This program will be segmented into overlays according to the size of memory available to it and the portions of code it needs to execute at any given time. That a program be segmented so as to minimize overlay swapping, must be an additional explicit requirement of a linker whose overlay swapping will be supervised manually. It is assumed that we cannot significantly affect the execution speed of the calculator. Thus, the intent of the minimization requirement should be obvious--suppress program segmentation which will tend to involve human thrashing.

The system source code must be portable. In the simplest case, it is desired that the unmodified source code be capable of utilization on any machine which possesses the resources to store, compile, and execute it. Because of operating system variations in such conventions as file naming and handling, transfer and processing of the source code in unmodified form on another machine (with operating system different from that on which it is developed) will be

very unlikely. However, the need for changes should be kept to a minimum and should be localized.

Finally, as with most all major software projects, maintenance and readability are considered paramount. Even after its completion, the system will certainly contain undiscovered bugs, areas to improve, and room for additions. Furthermore, development of a large prototype software system requires a great deal of careful planning for addition and modification. Adherence to the programming principles which support both readability and maintenance is absolutely necessary. Additionally, detailed documentation of the source code will supplement and assist in achieving these goals.

III. PRELIMINARY DESIGN

After requirements definition the next step in the software engineering process is the formulation of a preliminary design. Sound software design principles are applied to previously stated requirements to construct the framework for a software solution. It is during this phase of design that many of the most critical decisions are made. These decisions may be based upon a variety of considerations, each of which directly impacts the software organization. These decisions and the resulting organization are explored in this chapter.

A. PRELIMINARY DESIGN DECISIONS

Before a design can be formalized the engineer must investigate all design options and tools available. The following section summarizes the major decisions which strongly influenced and constrained many aspects of the project. With most large software projects, time is an extremely critical resource. As such, its impact upon preliminary design considerations is usually quite strong. Keeping on schedule is generally cost effective. It will be readily apparent that time played a key role in this design also.

1. Cross-Compiler

The fundamental considerations which most influenced design of the BAX59 (EAsic X-compiler for the TI-59) cross-compiler were the method of parsing it would use, and the language(s) in which it would be written. The availability of several versions of Pascal at the Naval Postgraduate

School and the working experience of the authors of this report with Pascal were, perhaps, the overriding reasons for its early selection as the design language. In addition, the extensibility, strong typing, and block structure of Pascal support modularity, readability, and maintainability. It was at this point that the parsing technique became an issue. The decisions were reduced to a selection between two alternate approaches. Berkeley Pascal was available on a Digital Equipment Corporation VAX-11/780 with Bell Laboratories' Unix operating system. The Unix system included software tools LEX and YACC which are capable of automatically generating a lexical scanner and LALR(1) parser from an input grammar. YACC allows the user to specify code generating actions which will be executed as the productions of the grammar are processed. The alternative system was entirely International Business Machines Corporation. IBM Pascal VS was available on an IBM-3033AP with VM/370 operating system. This system does not have the tools for automatic generation of a compiler front end. Instead, a scanner, recursive descent parser, and code generator would have to be developed from scratch.

While the prospect of automating the development of BAX59 seemed more time efficient and less trouble, it turns out that subtle problems involved are many. A compiler constructed using LEX and YACC is generated in the C language, a kind of structured assembly language. While it is possible to link object code compiled from Berkeley Pascal to object code compiled from the C language, the mixing of source code tends to destroy the portability and maintainability required of the system. Modifications or improvements to the finished system could only be made if the programmer were familiar with Pascal, C, and their interface. Likewise, a machine would be required to have both C and Pascal compilers in order to process the source

code for use. Thus, a recursive descent compiler in pure Pascal VS was the alternative selected. It quickly became apparent that a recursive descent compiler would be far easier to develop in pure Pascal. Using a block structured language which supports recursion, explicit use of parse tables and stacks is unnecessary. The activation record stack resulting from the recursive procedure calls implicitly holds the same information stored in a parsing stack. The advantage of using Pascal VS is the powerful debugging tools which this language system provides. While BAX59 was written in Pascal VS, to the greatest extent possible only those constructs and features which are standard Pascal [Ref. 3] were used.

Another major consideration involved the identification of the particular WBASIC language subset which could be translated to TI-59. Both feasibility and time constrained this selection. Commands and functions which primarily perform character string and file manipulation were quickly eliminated. The TI-59 is weak in alphanumerics and its storage capacity is too small to consider any concept of file handling.

The WBASIC language is rich with matrix and array functions and constructs. The overhead and difficulty of implementing these operations would outweigh any programming benefits they could provide. As a result, functions and constructs involving all composite data types were ruled out. With only slight overhead, it is possible to implement limited size, single-dimensional arrays. However, time restrictions required that this concept remain a suggested improvement.

In order to simplify the translation of a WBASIC source program, it was decided to allow the BAX59 scanner to recognize all WBASIC keywords as reserved words. This provided a clear distinction between real errors and

occurrences of legal keywords which had passed through the WBASIC interpreter but which had not been implemented in the subset. Otherwise, legitimate WBASIC keywords, not implemented in BAX59, would be treated and translated as identifiers. This obvious inconsistency might be very difficult for the user to detect as an error. It should be noted that WBASIC function names are not handled in this way. The reason is that the user can extend the BAX59 built-in function library. Further discussion of this idea is deferred to Chapter IV.

Appendix A is a summary of the WBASIC keywords and functions which have been implemented in BAX59 version 1.0. There are three general categories of keywords recognized by this cross-compiler. Command reserved words are implemented WBASIC keywords which indicate the start of a particular WBASIC construct or statement. Supplemental reserved words are implemented WBASIC keywords which cannot be used to begin a construct or statement, but which can be used (optionally at times) within same to guide the interpreter. Unimplemented reserved words are all WBASIC keywords which have not been implemented in BAX59. The use of this last category of reserved words will result in a fatal subset error during translation.

2. linker

In designing the linker three major problems arose. This first problem involved the fact that the linker is mainly a postprocessor of compiled data. As such, the linker is highly dependent on the compiler portion of the project. If this dependency were allowed, then most work on the linker would have to be deferred until the compiler was formalized and in the implementation phase of design. The second problem involved settling on a strategy to segment compiled code according to the software requirement to

minimize magnetic card reads. Two courses of action were discovered, each of which had advantages over the other. The third problem involved how prompting procedures were to be used to ensure proper execution of the segmented program. Procedures were required to be user-friendly and easily understood.

In the first problem it was decided to make the coupling between the linker and the compiler as loose as possible thereby reducing the dependency. This was achieved by defining a specific "third party" interface between the compiler and linker. This interface was defined to be a text file containing the four coded pieces of information required by the linker to accomplish its task.

This arrangement had several advantages and disadvantages. One advantage was that it allowed for the parallel development of the linker and the compiler. Since the interface was well defined, no other information needed to pass between the linker and the cross-compiler. By using this system, interfacing considerations such as naming conventions were nil since each process was totally independent. Another advantage concerned future implementations. It was envisioned that future versions of the system would be implemented on microcomputers. By having the project divided in half both logically and physically it would be easily adaptable to the more constricting memory requirements of the microcomputer environment. These two advantages alone outweighed the only major disadvantage of the decision. The disadvantage is that the linker needed to be able to regenerate the compiled linked TI-59 code structure which was originally produced by the compiler. As it turned out, this penalty was small when compared to the overall size of the finished linker.

The second problem was a very difficult problem to solve. Due to the limited size of the calculator memory and the cumbersome nature of the magnetic card backing storage system, the software requirements dictate the minimization of magnetic card reads. This requirement mandated the following decision: a code segment-break cannot occur within a backward loop. It would be preposterous to read a magnetic card every time a program encountered a segment break within a thousand iteration backward loop. This led to the following hierarchy of segmentation rules. First priority, segmentation may not occur within a backward loop. Second priority, maintain invoked subroutines with the invoking code. Third priority, keep adjacent sequential code together. To implement these decisions it became necessary to examine the control flow structure of an input program.

The decision as to how to accomplish this control flow examination is the foundation of linker design. Basically, two options were determined. One dealt with the input program as a whole and the other dealt with it as a series of sequential parts.

In dealing with the input program as a whole, the design algorithm would check to see if the program met the memory limitations. If it did not then the algorithm would examine the control structure and determine where to make an optimal break, that is, a break in a sequential portion of the program. It would then check each new segment to ensure that they complied with the memory restriction. The algorithm would continue until all segments met the memory requirements.

In the other method the program was decomposed into a series of sequential segments (a sequential segment is defined as a segment which does not contain a backward loop reference to any instruction other than possibly to the

first instruction of the segment). The algorithm first determined the sequential segments. Next the algorithm combined adjacent segments until the memory limit was encountered. At this point a segmentation occurred. The memory limits were reset and the combining process continued until another limit was encountered or the whole sequentially segmented input program was processed.

The second method was selected for two reasons. The first reason is that the first method eventually required the evaluation of code on a small segment level much like the second to determine a suitable segmentation point. Rather than do this and more, it was decided to just evaluate the small segments and build up rather than down. The second reason was that the second method lent itself to a recursive solution during the recombination process. The recursive solution greatly reduced the length and complexity of the segmentation code.

The third problem involved deciding upon a method for accomplishing the prompting of the user. One method dealt with assigning coded prompt numbers of short length to be built into the code. The other method involved building larger self-explanatory prompts into the code. The second choice was selected. This was done to reduce the number of instruction references that a user might have to make during the execution of the generated calculator program. This was in keeping with the requirement to make the system user-friendly.

B. PRELIMINARY DESIGN ORGANIZATION

Thus far the system design space has been narrowed to the design language, source language subset, and the general techniques for compilation and linking. Organization of the software into functional categories may now begin. This

next phase of development is characterized by a more specialized, yet still preliminary consideration of system components. It should be apparent by now that a natural division into two major functional components, cross-compiler and linker, has been assumed since conception of the system. For a two-man design team, this partitioning appeared to have the greatest potential for success. It allowed the simultaneous development of two independent system components of low coupling [Ref. 4: p. 85] and high cohesion [Ref. 4: p. 106]. The result of this separation was a minimization of programmer interaction, maximization of work time efficiency, and simplification of interfacing. The remainder of this section outlines the preliminary design and organization of the cross-compiler, linker, and the direct interface between them.

1. CROSS-COMPILER

The common form of all versions of BASIC language can be characterized as imperative, line-oriented, and sequential. The design of BAX59 is based upon this fact. Each WBASIC line is parsed, beginning to end, by recursive descent. Equivalent TI-59 code is generated for each line concurrently with the parsing operation. This means that BAX59 will successfully translate a sequence of syntactically correct yet meaningless WBASIC statements into equivalent TI-59 code. However, the TI-59 code will be as meaningless as the source code. For this reason, it is recommended that the user successfully execute his WBASIC source program using a WBASIC interpreter (or compiler/loader) prior to translation with BAX59.

A line-oriented view of the source code provides several advantages to the design. First, there is a direct sequential correlation between the original source program and the translated TI-59 code. As will be seen later, this

allows easy management of the generated code and its associated data. Second, the parse driver routine can be a fairly simple loop, since lines are parsed to end of line one at a time until the end of the source file. Third, viewing the source as a sequence of independent pieces greatly facilitates an iterative enhancement approach [Ref. 5] to the progressive development of BAX59. This approach virtually guarantees a working prototype throughout the coding phase, and supports both reliability and maintainability. Modifications can be quickly tested within the context of the entire compiler system to date. Upon completion, the programmer tends to have greater confidence in his product, because a great deal of testing has already been conducted.

BAX59 was temporally organized into three major functional sections: initialization, translation, and resolution. The primary operations performed by initialization involve setting up data structures and initializing variable values. There are three major data structures manipulated by the translation section: the reserved word table, the symbol table, and the code data structure. Conceptual subdivisions of this section, namely the scanner, the parser, and the code generator, manage each database respectively. While the scanner is a separate routine by itself, the parser and code generator are not as separately defined. These functions are actually performed concurrently by a set of mutually recursive procedures under the direction of the main driver. This driver calls the correct subprogram into execution as its corresponding WBASIC construct is recognized. Once translation has been completed, the resolution section processes the generated code into a form suitable for final output. This includes label insertion, peephole optimization, and absolute address resolution.

2. Linker

The linker was organized into three phases. The design of these phases were the direct results of the preceding design decisions made in the preliminary phase.

The first phase is the direct result of the loose coupling between the linker and compiler and the decision to combine small sequential segments to form memory-size constrained segments. It is the preprocessor phase. This phase processes the interface input file and reconstructs the compiled linked code structure. In addition the preprocessor determines the sequential segments of code and constructs an internal data table called the segment table which is used by the second phase, the segmentation phase. The segmentation phase utilizes the recursive algorithm to recombine sequential code. The postprocessor phase is the result of output design decisions. This phase inserts the prompting code and develops the segmented code lists to be output to the user. It then produces the code in a text format together with specific instructions as to its use.

3. Direct Interface

The design organization is built around the loosely coupled compiler and linker. This coupling is made possible through the rigid definition of the interface text file. The organization of this text file is critical to the design and will be described in more detail.

The text file is the only direct transmittal of data between compiler and linker. Four pieces of data are transmitted. They are the following: a number signifying the next register available for use; generated code list in a numeric formatted form; text containing DATA/READ information; and text containing the mapping of TI-59 registers to BASIC variables. Each piece of information is preceded by a

\$XXX in the first column where XXX is a number. This simple format enables the linker to easily locate the correct information and process it accordingly. Since this is the only explicit interface, the compiler and linker are not as dependent on each other as they would have been in a closely coupled system.

IV. DETAILED DESIGN

The source code contained in Appendices C and D provides high resolution understanding of the system. However, in order to provide rationale behind design issues for a language system comprised of almost 10,000 lines of code and comment, discussions of some detail are in order.

It is not our intention to explain everything. What we wish to do in this chapter is to introduce design details and strategies of the more important concepts and components in the system. This will serve to illustrate software engineering and how it is used in this project.

Upon the completion of preliminary design, the detailed design is begun. It is during this phase that the actual details of full implementation are defined and laid out prior to coding. Categorized under cross-compiler and linker, the general format for the sub-sections of this chapter will include an informal solution strategy for a specific design problem, followed by a discussion of the major data objects and procedures which manipulate those objects. Where appropriate, inter-procedure interfacing criteria are outlined. In several cases, significant problems, their solutions, and possible system improvements are discussed. The last section presents implicit interface design which impacts greatly on the system.

A. CROSS-COMPILER

The fundamental design of BAX59 is a finite state machine driven by a main loop. Once values and data structures have been initialized, program control enters the main loop which scans, parses, and generates TI-59 code for one

WBASIC source line. At the end of line, the loop checks for end of source file. As long as the end of file is not detected, the main loop repeats its processing of each succeeding line of source code. When end of file is found, control exits the main loop and post-processing begins on the generated TI-59 code. This includes insertion of suspended code data, optimization, address resolution, and final output of code and associated data.

The entire cross-compilation process is broken down into 15 functional areas. These are outlined by the contour diagram in Figure 4.1 Note that solid lines indicate actual procedures (P/) or functions (F/), while dotted lines only indicate logical association. Although these areas often depend upon one another, the particular services performed by each differ enough to allow independent analysis. Were it not for limitations imposed by the Pascal language, many more procedures and functions would have been tightly packaged in order to hide implementation details between functional areas. What follows is a survey of the major functional areas of EAX59 and the design details for each.

1. Initialization

In the interest of execution time efficiency all run-time actions which required completion prior to the start of parsing are incorporated into procedure INITIALIZE. As a result, some variables and databases which would otherwise have been local to their respective procedures, required globalization. One particular example is the reserved word table. This data structure is built of data supplied from outside the program in file RWTBLF. Since the reserved words contained in this file need only be loaded once (at the start of execution), there was no good reason to place them into procedure SCAN, which is called often by parse routines. While it would have been possible to use a

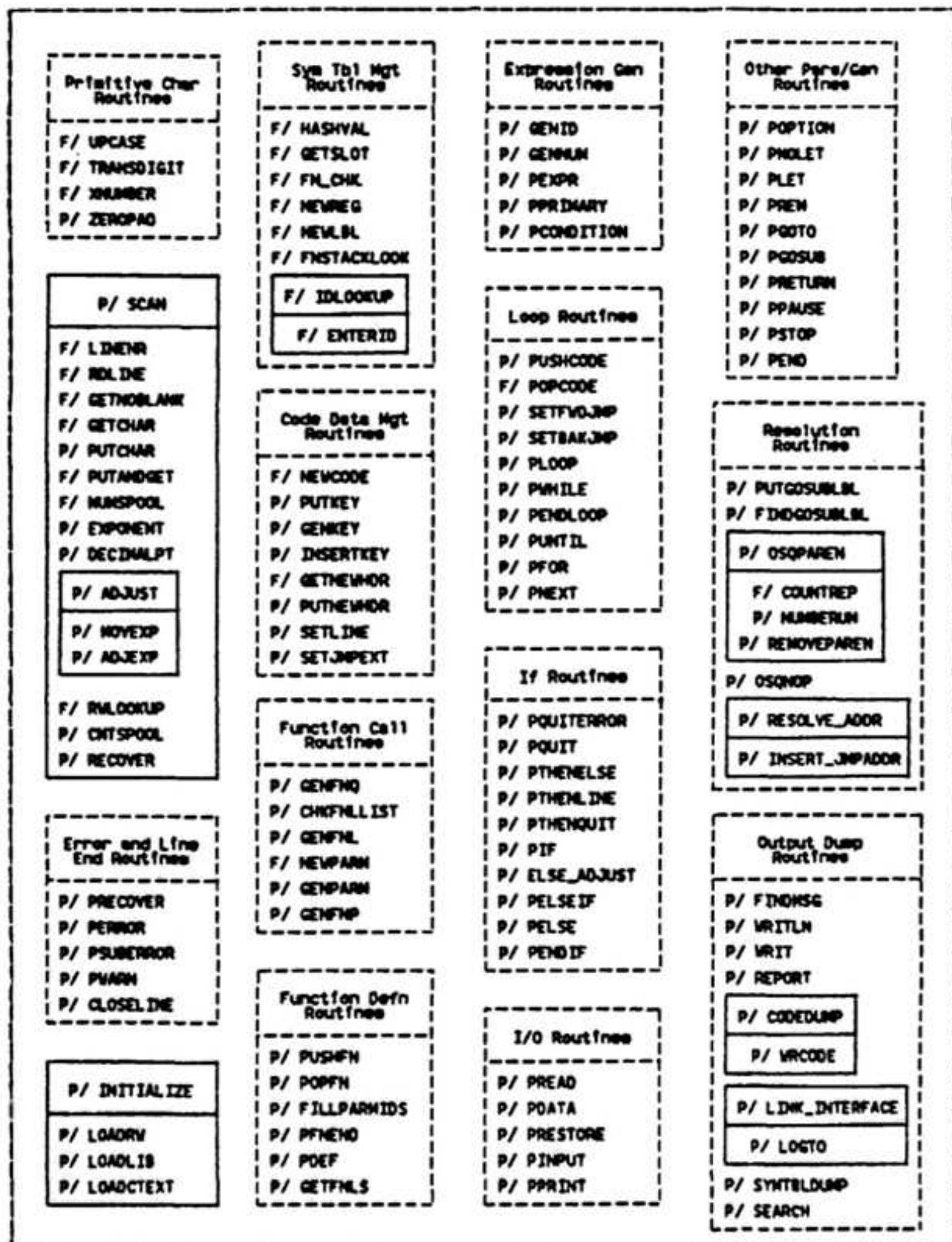


Figure 4.1 Cross-Compiler Contour (PROGRAM BAX59).

boolean switch to detect the initial call to SCAN and subsequently generate the reserved word table at this first call, this would have increased the coupling between modules. Loading the reserved word table from outside the program offers several advantages. First, changes to the reserved word table can be made easily (this process will be discussed when the scanner is considered). Second, the user can check one file (RWTBLF) in order to see the reserved words recognized by FAX59. Third, the loading routine does not need to know the words themselves, only the name and format of the file in which they reside.

Built-in functions are also loaded by procedure INITIALIZE. This operation actually requires two ordered steps. First, the symbol table must be created. Second, function recognition and generation data is read from outside files (BIFNCF and BIFNLF) and the symbol table management routines are used to put this data in the table. Loading the built-in functions by using the same routines which the parser will use to manage the symbol table, ensures table consistency and promotes readability. This approach has been taken as often as possible in designing the cross-compiler.

The complexity of the above initialization processes as well as the the TI-59 keycode text (CTEXTF) loading process required that these operations be abstracted into individual subprograms. However, procedure INITIALIZE performs many other pre-compilation activities which could be performed sequentially. Probably the most important of these activities is the simple initialization of variable values. The importance of this task is elevated by a serious hole in the Pascal VS implementation. The Pascal VS compiler will not detect the failure to initialize a variable value prior to its use. What is worse, random values which exist in pointer references or other variable storage

areas will be used as is, whether they were put there by the user or not. As a result, failure to initialize values was a major source of error during development of BAX59. These types of errors tended to be extremely difficult to debug, since they often surfaced late and usually in modules long thought to be robust.

In summary, procedure INITIALIZE loads all information which will be needed by the translation and resolution stages, constructs conveniences data such as character sets, assigns starting values to all variables and pointers.

2. Scanner

At the lowest level of abstraction within the translation component of BAX59 is the scanner, procedure SCAN. This single, self-contained subprogram is basically designed on three important concepts. First, the scanner is itself a finite state machine. Second, with the exception of procedure INITIALIZE and some system constants, its implementation is transparent to the rest of the cross-compiler. Third, the database which it uses for token recognition is simple, time efficient, and general.

The state machine logic of procedure SCAN provides knowledge of token streams in free format and nothing more. Its primary job is to read the source file character by character in order to isolate and recognize single tokens. However, procedure SCAN is designed to do much more. First, it reads and converts line numbers. It also fills as necessary the line buffer and accumulator, the data structures which store the line and token currently being scanned respectively. The scanner also detects the end of a source file which has no explicit WBASIC "END" statement. This allows a more graceful conclusion to what might otherwise be an abrupt exit.

Two other functions performed by the scanner will illustrate its transparency to the rest of the cross-compiler. These are recognition of the end of a line and of continuation lines. Procedure SCAN reads and loads the line buffer (LINEBUF) with a new line of source text each time the end of line character ("@") or continuation line character ("&") is found. The only difference is that the end of line token number must be passed up to the parsing routines so that the main loop will know when it has parsed one entire line. On the other hand, the continuation character can remain invisible to the parser, which views only whole source lines.

As mentioned before, the database used by procedure SCAN is the reserved word table. Although referred to as a table, the internal representation of this database is actually three coordinated arrays constructed from the RWTBLF file by procedure INITIALIZE. These arrays are used to compare the characters of a token in the accumulator to the characters of reserved words. The simplicity and efficiency of this comparison is illustrated in Figure 4.2, which depicts a condensed schematic of the arrays. Note that the characters in the RWCHAR array are arranged in order of increasing word length. A reserved word look up is based upon the length of the word in the accumulator. Comparison begins at the first character of the first word in the RWCHAR array which matches the length of the token in the accumulator. Comparison ends when either all characters in the accumulator match a string in the RWCHAR array, or when the characters of all words of a given length have been compared to the accumulator without success.

The RWORD array references the start index of each word in the RWCHAR array, while the RWLENG array references the start index for the first word in the RWORD array for that length index. The indexes of the RWORD array are used

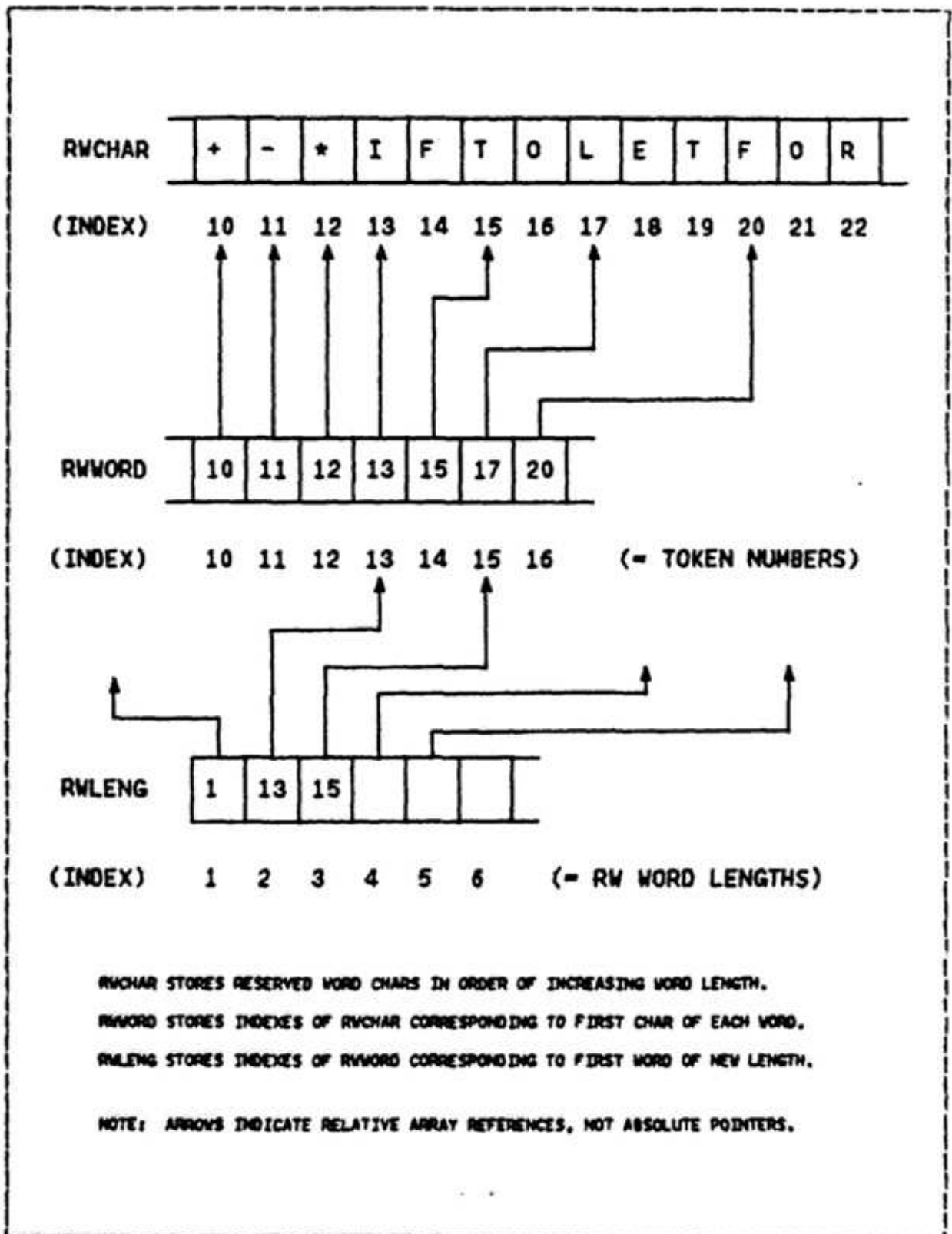


Figure 4.2 Reserved Word Table Arrays.

as the token numbers returned by the scanner after a successful look up. Should a look up operation fail to recognize the accumulator token, then the token is assumed to be a variable identifier. Other token types, such as numerics, are recognized prior to table look up. This mechanism and the fact that the scanner is independent from the parser require that WBASIC keywords be reserved. If keywords were overloaded as variable identifiers, the parser would have to communicate its token type expectation to the scanner. Token overloading would greatly hamper readability.

One scanner related problem which required a relatively complex solution concerned the conversion of WBASIC to TI-59 numeric values. The calculator display window restricts the number of significant figures which can be entered from the keyboard. For numbers without exponents, a maximum of ten digits (with decimal point) can be entered. Numbers with exponents are allowed a maximum of eight digits in the mantissa and two in the exponent. Because of this restriction, a rather complicated procedure was designed to convert WBASIC numeric values to TI-59 compatible values without losing equivalence. Procedure ADJUST performs decimal point shifting and exponent modification on WBASIC numerics which contain too many significant figures for the TI-59. The operation can, of course, reduce significance by truncation of excess digits. Except for this loss of digits, equivalence is maintained.

3. ERROR HANDLING

At this point it is appropriate to discuss error recognition and recovery. As implied earlier, the error detection capability in BAX59 is relatively weak and incomplete as compared to full language compilers. The reason is that the system requirements specified error-free input

source files. The primary use for this system is as a supplement to an existing BASIC language interpreter or compiler. Debugging of TI-59 programs is a hard enough task without adding the complexity imposed by absence of BASIC language run-time diagnostics. Therefore, users of BAX59 are strongly advised to ensure that a WBASIC program is correct syntactically, semantically, and logically by running it in the WBASIC environment, before translation to TI-59 code. Error handling in BAX59 is restricted to detection of subset related exceptions, calculator capacity limits, and errors of opportunity.

The cross-compiler is designed for recognition of two major types of errors: fatal errors and warning messages. Fatal errors are further categorized as scanner or parser detected.

Warning messages are generally unrelated to WBASIC syntactic or semantic problems. They refer to potential difficulties with the TI-59 run-time environment, most commonly (but not always) calculator capacity. Such conditions as too many registers in use, too many labels in use, or excessively nested subroutine calls, will trigger warnings. Each message is explicit and cautions the user of a situation which is considered abnormal to the calculator. Since these errors are unrelated to the WBASIC source code, warning messages do not halt the parsing or code generation processes. However, TI-59 code generated from a WBASIC source file that produced warnings is not guaranteed to execute properly, if at all.

The warning message is similar to non-fatal errors in full language compilers. The reason for continuation of code generation is slightly different. A user of BAX59 will most likely need to modify and tailor his WBASIC program to fit calculator constraints and capacities. Warnings are a non-fatal means of providing near equivalent code data for

use in comparison of efficiencies, capacities, or consistency. Even though the code may not successfully execute on the calculator, it still represents a direct translation from WEASIC and is a fairly accurate indication of program size, register/label use, etc.

Unlike a warning, one fatal error will flag the main loop against further parsing and code generation. However, scanning for tokens continues until the end of the source file is reached. Thus, only a single fatal syntax error can ever be detected in one BAX59 execution, although the scanner will continue to detect any number of lexical errors. Fatal errors are also categorized as subset or non-subset related. Non-subset errors are those previously referred to as errors of opportunity. During the coding phase of development, simple syntax checks were often inserted into the logic of the parsing routines. These were usually one-line IF-THEN-ELSE constructs which cost very little but were highly protective. For example, the main loop calls procedure PGOTO whenever the GOTO command is recognized. Since error-free input is required, this procedure could have been written to assume that the next token must be a numeric line reference. Instead, it was a simple matter to check the next token's type and call a syntax error (FERROR) if it were not numeric. Note, however, that the logic of PGOTO will not call an error for a numeric token which contains a fractional part, clearly a syntax error. In fact, the cross-compiler is not likely to detect an error at all. Execution may result in a Pascal VS runtime error. The reason is that the numeric string will be converted to an integer value based on ordinal values of the characters. The decimal point will appear to BAX59 as any other character. However, its ordinal value will be added during conversion resulting in an inconsistent integer value for the line reference. The routine used to set jump

pointers will probably not be able to find the line number since it is already in error.

Although often incomplete, these error traps provided much assistance in tracking system bugs. The technique used was to translate simple source programs known to be correct. Errors, tripped at these check points by system bugs, usually indicated the likely trouble spots. The faulty routine had helped in parsing either the statement which caused the error or the statement which immediately preceded it. Since the token which tripped the error was also known, the exact routine and the specific bug were easily found.

Subset related errors were defined in the software requirements. The user must be told where and how he has misused the system. BAX59 incorporates all WBASIC keywords (Version 2.0) in its reserved word table. The main loop logic contains the information to distinguish between implemented keywords and unimplemented keywords. This technique allows the reserved word table and implemented subset to be easily expanded (or contracted). Such a technique strongly supports the requirement for maintainable source code.

There is more room for improvement in the area of error detection and handling than in any other aspect of the cross-compiler. The capability could certainly be extended to protect against all possible syntactic and semantic errors so that prior compilation or interpretation would be unnecessary. However, the benefits to be gained are questionable, since run-time and logical debugging of TI-59 programs is no easy task. A special file to hold error message text might help to reduce some of the awkwardness in portions of the code which issue these messages. Within this file the messages could be indexed by number, thereby allowing more verbose and possibly clearer explanations of errors. Generally speaking, the critical resource of time forced the design of error handling to be barely adequate.

4. Symbol Table Management

One of the more important duties of the parser is to manage the symbol table. The BAX59 symbol table is a variable bucket hash table similar to one described by Aho and Ullman [Ref. 6]. The data structure used is an array of pointers. The indices to this array of base pointers are hash values computed by taking the modulo 99 sum of the ordinal values of identifier characters. This operation is performed by procedure HASHVAL. Figure 4.3 depicts the structure of the table itself and its four types of identifier entries. Three of the four types of identifiers are functions. These will be discussed later in this chapter. The important structural feature to notice now is that each node has a SLOTP field regardless of variant tag. The SLOTP field is each entry's link in the variable length chain which forms a bucket. In order to insure that no uninitialized pointer references or variables occur, new nodes are created as needed by the separate function GETSLOT. The job of this function is to create the node and to insure that all of its fields have been initialized to default values. This same approach to data structure construction is used throughout BAX59 in order to protect against random initialization by the Pascal VS compiler.

The look up operation of procedure IDLOOKUP is simply to hash the characters of the identifier token in the accumulator to the correct base pointer bucket in the symbol table. The IDENT field of each slot node is compared to the accumulator token until either a match or a nil pointer is found. If a match is found, a pointer to that slot is returned. If no match is found, then a slot for the accumulator identifier token is automatically added to the symbol table in the bucket just searched. A pointer to this new slot is then returned. This is in accordance with the

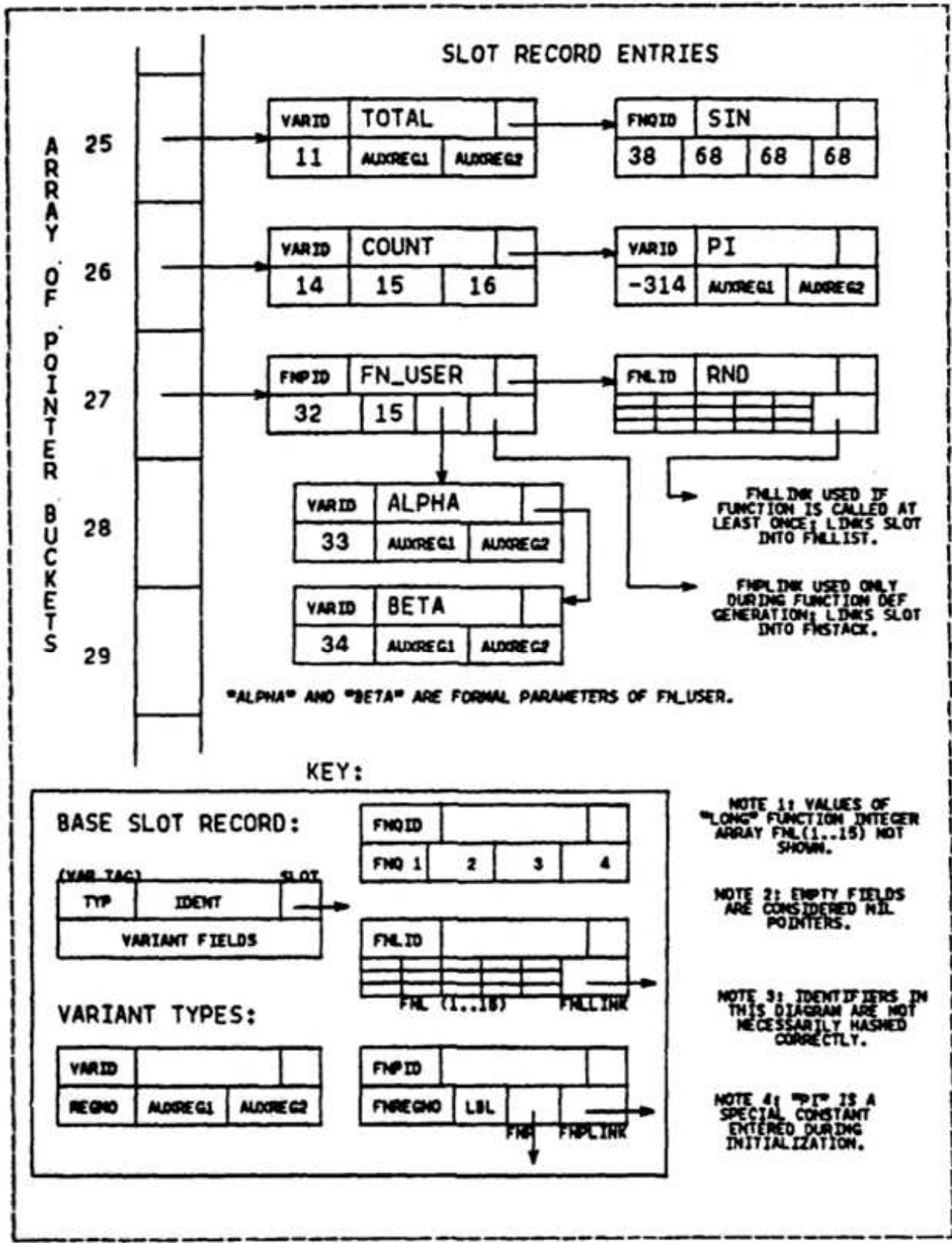


Figure 4.3 Symbol Table.

semantic rules of the BASIC language which allows implicit declaration of variable names by using them in statements.

The insertion of new identifiers into the symbol table is performed by procedure ENTERID. This procedure creates the new slot, fills all fields which are known, and links the slot into the symbol table. It is also during this process that identifiers are assigned TI-59 registers for code generation purposes. Function NEWREG handles the register pool, which is actually nothing more than an implicit stack of integers. An important feature regarding the assignment of registers to variable names is that the user has some control over these assignments from outside the program. Included in the LABELP file is a place to list register numbers which the user wants to reserve for his own use. Function NEWREG will not assign these numbers to WBASIC variable names. The significance and power of this control feature becomes more apparent during the discussion of functions. The user is cautioned against reserving the last assignable register number (system parameter in constant declaration block: REGBASE). Reserving this register will short circuit the logic which reports a TI-59 memory overflow warning message, the situation in which too many registers are in use.

As a final note, there are two forms of output which are closely associated with the symbol table. One is the WBASIC variable name to TI-59 register mapping which correlates variables to register assignments. The other is an optionally available symbol table image, which lists each table entry in bucket order with type and register assignment. Both outputs are discussed in the last section of this chapter.

5. Expressions

The most fundamental and most common construct seen by the parse routines is the arithmetic expression. The many similarities between BASIC language expressions and TI-59 expressions make them relatively easy to parse and generate. However, a few subtle differences cause abnormal situations requiring careful design. If there is one lesson to learn from this discussion, it is this: in compiler design, when in doubt revert to the grammar specification.

TABLE I
Production Rules for Expressions

```
<EXPRESSION> ::= <PRIMARY> {<BINARYOP> <PRIMARY>}
<PRIMARY> ::= {+|-} <PRIMARY> |
               <UNSIGNED NUMBER> |
               <IDENTIFIER> |
               ( <EXPRESSION> )
```

Table I lists the grammatical specification for a WBASIC expression. The two production rules in Table I are abstracted by the two BAX59 parse procedures PEXPR and PPRIMARY. They are designed to parse and generate code through mutual recursion. Careful examination of the case statements within these procedures will reveal the differences between WBASIC and TI-59 expressions. While both use infix notation for binary operators, unlike WBASIC, TI-59 unary operators and function applications are postfix. This minor twist in notation adds a little complexity to the logic of the expression parsing routines. However, once designed, the code for translation of expressions became the

fundamental base upon which assignment statements, conditional expressions, functions, and many other constructs could be built.

6. Unstructured Jumps

Some of the easiest constructs to understand and implement were the unstructured control statements GOTO and GOSUB. To realize their simplicity it is necessary at this point to introduce the code data structure which is constructed by the generation routines.

Illustrated by Figure 4.4, the code data structure is, perhaps, the most unique design concept of this cross-compiler. There are two types of nodes: WBASIC line number nodes and TI-59 keycode nodes. Since unstructured constructs in WBASIC are dependent upon source line numbers, there had to be a method of associating the TI-59 code with those same line numbers. Figure 4.4 shows how line nodes and code nodes are linked to duplicate this association. It is important to note that the TI-59 code chain is completely independent (and may be traversed as such) of the WBASIC line number chain. The line nodes merely provide a frame of reference for the TI-59 code.

Procedure SETLINE, called at the beginning of the main driving loop, is responsible for insuring that new line nodes are created and inserted into proper order. As each line is parsed, special holding pointers (FIRSTLP, LASTLP, BEGINCP, ENDCP, LPLEAD, LPTRAIL, LPCUR, CPCUR) keep track of all key locations in the structure. As Figure 4.4 indicates, it is possible to have line nodes created and linked prior to their encounter in the source code. This occurs whenever a forward jump (GOTO) is parsed. Since the line reference of a forward jump has not been parsed, its line node would not exist. However, the jump pointer (JMPP) must be anchored to a node. So the line node and an anchoring

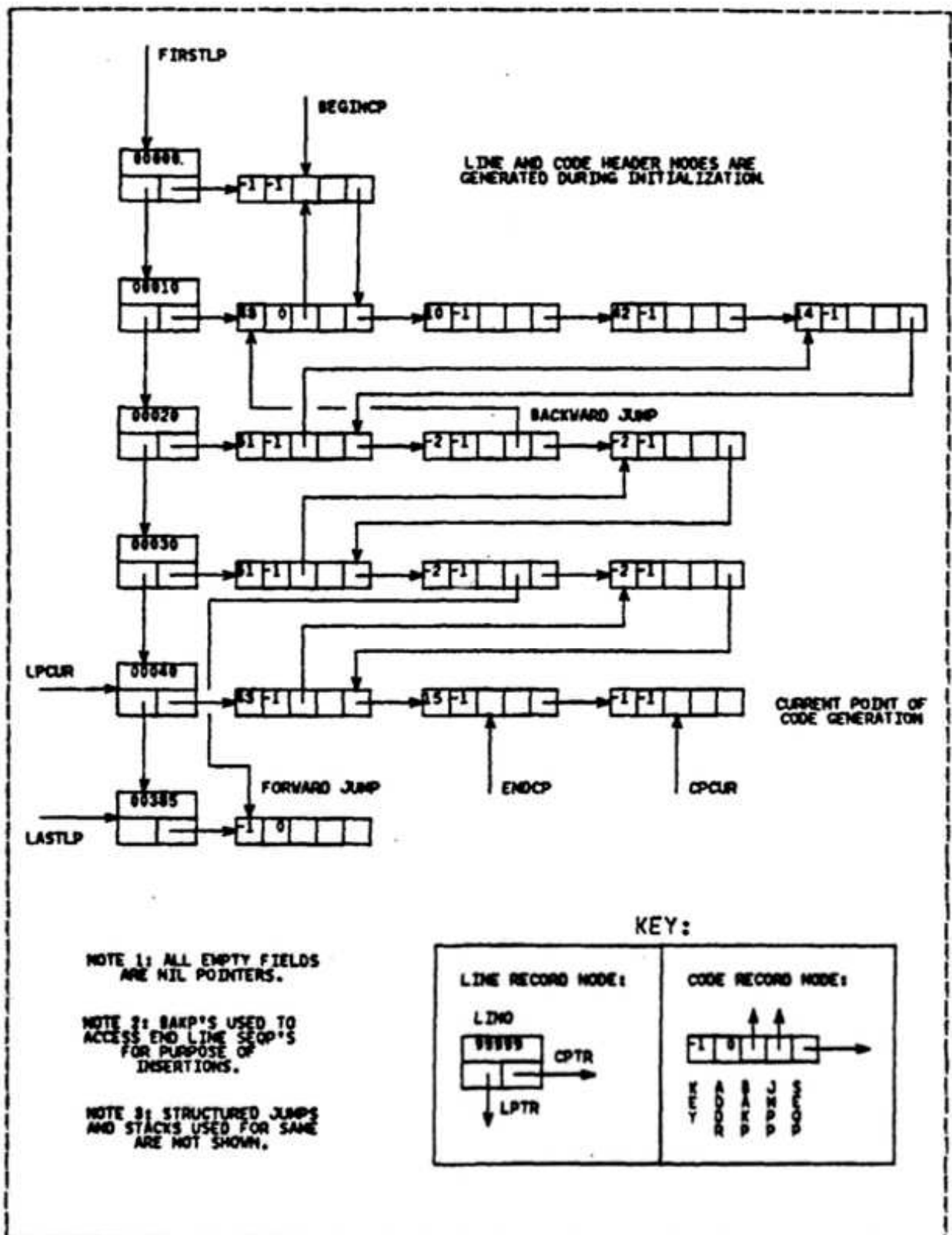


Figure 4.4 Code Data Structure.

code nodes are inserted in correct order ahead of the current line. Procedure SETJMPNEXT sets forward as well as backward jump references. Of course, setting backward jumps is easier because the line number node has already been created and is in place. It should be noted that procedure SETLINE always checks its forward line number chain before creating a new line node. If a line node exists whose line number field (LINO) is equal to the next WBASIC line number, then it will be used instead.

The technique for handling GOSUB statements is similar to but slightly more involved than the GOTO. Since the GOSUB is actually an unstructured subroutine call, it was necessary to maintain consistency in code generation so that the linker could recognize the difference between subroutines and unconditional jumps. All TI-59 subroutines are prefaced with and called by a label name. Therefore, while initially the GOSUB can be treated as a GOTO, at some time later a label must be inserted at the head of the subroutine body, which is the node referenced by the jump pointer. This is done during the resolution phase of compilation by procedure FINDGOSUBLBL. This fairly tricky insertion is one reason for the existence of the back pointer field (BAKP) in code nodes. This operation will be explored in the sub-section on resolution.

7. Looping and Branching

Users of BAX59 are strongly urged to practice structured programming when writing WBASIC code. Translation to and execution on the TI-59 are far more regular and predictable when the input source code is structured and readable. Much of the design of the entire system is based upon the assumption that the source program will be structured. You will understand why more thoroughly in the section describing linker design.

We begin discussion of structured looping and branching by introducing procedure PCONDITION. This procedure is fundamental to the parsing and code generation for simple boolean expressions (compound boolean expressions have not been implemented). While WBASIC has a fairly common set of boolean operators, the TI-59 does not. There was a need to construct efficient sequences of TI-59 code which are equivalent to the WBASIC boolean operators. These equivalences, shown in Table II, are implemented in

TABLE II					
TI-59 Keycode Sequences Equivalent to Boolean Operators					
A = B	A <> B	A >= B	A <= B	A > B	A < B
RCL A	RCL A	RCL A	RCL A	RCL A	RCL A
X->T	X->T	X->T	X->T	X->T	X->T
RCL E	RCL E	RCL B	RCL B	RCL B	RCL E
INV	X=T	X->T	INV	X>=T	X->T
X=T		INV	X>=T		X>=T
		X>=T			

procedure PCONDITION. While it would have been possible to implement compound boolean expressions (AND and OR), the lack of time and the fact that their logic could be duplicated using IF statements prevented this enhancement. It was, however, a very simple extension of logic to recognize and translate a negation (NOT).

In the implementation of a block structured language which allows nesting, the use of stacks is an important technique. And so it was with looping in BAX59. By nature loops involve backward jumps. As with unstructured jumps, there existed a need to anchor pointers on code nodes whose

source code had yet to be translated. In the case of loops, the inverse of this concept is also true. There was a need to create pointers from code nodes whose source code had yet to be translated. Since structured nesting of loops is checked by the WEASIC interpreter, it was possible to pre-create these nodes and push them onto a stack until their place of insertion is encountered. This is exactly how loops are translated using LOOPSTACK and ENDLOOPSTACK. When the LCOP statement is encountered, a NOP keycode node is created and pushed onto the ENDLOOPSTACK. In the case of the WHILE statement, a boolean expression is processed, a forward reference node is created, a jump pointer is set to the node (for the false branch to end of loop), and this node is pushed onto the LOOPSTACK. When the ENDLOOP or UNTIL is found, the stacks are popped and the NOP nodes are inserted. The nature of correct nesting guarantees that NOP code nodes popped from the stack will have jump pointers referencing or will be referenced by the appropriate code nodes.

Iterative loops are written by using the FOR-NEXT construct. The stack implementation is similar to that described above. The main difference is in the additional calculator resources required for such a loop. Unlike ordinary variable names, the FOR loop variable requires from two to three register assignments. The fields AUXREG1 and AUXREG2 in the VARID tagged slot record are used for this purpose. AUXREG1 holds the TI-59 register number which will store the upper (lower) limit of the FOR index variable. The FOR index variable increment will always default to +1 unless the STEP option is used. If STEP is used, then AUXREG2 will hold the register number which stores the increment value. The user should understand that use of a FOR loop carries a fairly heavy overhead in terms of both register and program step use. A single simple FOR

statement in WBASIC translates to use of two registers and over 20 program steps. Most of this overhead is caused by the run-time checking of the FOR index variable value against its limit for each iteration through the loop.

Branching is another construct implemented with stacks. There are actually two forms of branching: the unstructured or line-oriented IF statement and the block structured IF statement. The unstructured IF is actually only partially implemented. WBASIC allows either a jump to a line number or execution of any single statement within each of the IF branches. Because the structured IF can be written to perform the same way, it was decided to restrict the unstructured IF to allow line jumps or the QUIT statement. The implementation of line-oriented jumps has already been discussed. An IF-THEN-QUIT or IF-THEN-ELSE-QUIT is handled by setting a jump pointer to the ENDLOOPSTACK. The effect is to force program control to exit the current loop. If control is not within a loop (i.e. ENDLOOPSTACK is nil), then an error condition is raised. FOR loops are not considered loops in this context.

The more powerful of the two IF statements is the block structured IF-ELSEIF-ELSE-ENDIF. This form disallows the use of keyword THEN, since it is implied. Once again, stacks are used to implement the structured IF. The logic and its correspondence to the manipulation of stacks is roughly similar to that of looping. Instead of directing jump pointers to the end of loops, they are directed to the next ELSEIF or ELSE. An unusual situation occurs, however, in the case of the IF statement. Stack manipulation for the IF-ENDIF is slightly different from that for the IF-ELSE-ENDIF or the IF-ELSEIF-ELSE-ENDIF. To understand the problem, assume the viewpoint of a parser which has just evaluated the condition of a structured IF. At this point you do not know whether or not an ELSE/ELSEIF or an

immediate ENDIF will follow the true branch. To which stack will the jump pointer of the false branch be set? In order to cover both possibilities, a pointer to a node pushed onto each stack is required. However, there is only one jump pointer field (JMPP) for the code node which represents the jump address to the false branch. Our solution uses the back pointer field (BAKP) to reference a node in the IPSTACK, while the jump pointer field (JMPP) references a node in the ENDIFSTACK. Procedure ELSEADJUST performs the resetting of pointers required when an ELSE/ELSEIF is encountered. When the ENDIF is encountered, the BAKP is tested for a nil pointer. A nil BAKP at the top of the ENDIFSTACK indicates that an ELSE/ELSEIF has been seen. This is because procedure ELSEADJUST is the only routine which can clear the BAKP reference before the ENDIF is encountered.

The cause for all the foregoing complexity is the fact that IF-ENDIF has a single false branch which must be a jump past the ENDIF. The tail of the true branch merely falls through the ENDIF. On the other hand, IF-ELSEIF-ELSE-ENDIF can have several false branches, only one of which may jump to the ENDIF. The tails of the all true branches must be jumps to the ENDIF. The logic of BAX59 is designed to recognize and generate equivalent TI-59 code for any of these possibilities.

8. Functions

The most powerful feature of the BAX59 cross-compiler is the translation of functions. Both built-in and user-defined functions are handled. In order to take full advantage of the calculator's capabilities, it was necessary to design three distinctly different types of functions. The first type, referred to as "quick" functions, are the common arithmetic/trigonometric functions such as LOG, SIN, COS. The second type of function harnesses the power of the

Solid State Software module. These are referred to as "long" functions. Both of these first two types are built-in functions. The third type is user-defined "parameter" functions, which are translated from #BASIC source code specifications.

The difference between "quick" and "long" functions is basically the number of TI-59 program steps generated for each. "Quick" functions generally translate to a one or two step TI-59 keycode sequence. However, they may have as many as four steps. Because they are short, "quick" functions are inserted as in-line macros. On the other hand, "long" functions may translate to as many as 15 steps. Therefore, their length requires that they be called as subroutines rather than translated in-line.

Read from the BIFNQF and BIFNLF files respectively, the code for both "quick" and "long" functions is entered into the symbol table during initialization. BIFNQF and BIFNLF may be revised by the user from outside the cross-compiler. By knowing the TI-59 key stroke sequence, a user may add his own functions to either file. As a special user note, the format for additions to these files is critical. The number of key strokes in a function sequence may not exceed the maximum limit for the type of function. If less than that limit, then the end of the sequence must be padded with NOP (68) key strokes to the maximum limit. These limits may be altered by adjusting the system parameters FNQLEN and FNLEN in the constant declaration block of the BAX59 source code.

Most all functions that could be implemented as "quick" have been and are listed in the BIFNQF file. However, only the RND(X) (random number generator) function has been implemented as "long." To illustrate the concept of "long" function, we will walk through the design of RND(X).

Suppose you desire to write a TI-59 program which uses a random number generator. You might write your own pseudo-random number generator subroutine, but the TI-59 has such a routine built into its read-only Solid State Software module. Use of this built-in facility would clearly be more space efficient. WEASIC also has such a function, RND(X). If it had not, it would be possible to write one in WEASIC using the DEF FN_RND(X) statement. Before translation it would be necessary to remove the function definition block and replace the FN_RND(X) calls with RND(X). However, this is not required in our example. You must ensure that the TI-59 registers used by the Solid State Software module to run the RND(X) function are reserved in the LABELF file. This information can be found in the Master Library Manual [Ref. 7], which is the Master Solid State Software module reference guide. RND(X) uses registers 01, 02, 03, 04, 05, 06, and 09. "Long" functions always take a single parameter (even if it is a dummy). Register number 10 has been designated to store this parameter and should also appear on the reserved list. This parameter register assignment may be changed in system parameters of the BAX59 source code if desired. Each time it is encountered within the WEASIC source program, RND(X) will translate as a call to a subroutine whose single parameter is stored in register 10. The first time seen, the RND symbol table node will be linked to a special list (FNLLIST). At the conclusion of code generation, FNLLIST will be traversed and the key sequence which executes RND(X) as well as any other "long" functions on the list, will be added to the code data structure as subroutine bodies.

The real power of this facility lies in its user-controlled flexibility. The user may convert almost any program function in the Solid State Software module into a single parameter "long" function. All he must do is reserve

the correct registers in the LABELP file, list the key sequence in the BIFNLF file, and, if necessary, fix the values of all but one of the function input parameters (or create a dummy). If the function does not exist in WBASIC, then he must write the DEF block for it in order to check program correctness prior to translation.

Having strayed from implementation design toward system utility, we now return to implementation discussion of so-called "parameter" functions. The name given these user-defined functions applies more to how they are implemented rather than to their nature. The parsing routines always expect parameters but do not require them. Parameterless functions are, indeed, recognized.

Although the cross-compiler will correctly translate a function definition (DEF statement or block) whether it occurs before or after its respective call, the linker requires that all subroutine/function bodies be placed after the main program.

When a new function identifier is recognized (by the "FN_" prefix), a new FNPID tagged slot is created for the symbol table. Procedure GENPARM is then called upon to parse actual parameter expressions, generate code which performs their run-time evaluation, and construct the formal parameter list. Parameters, if found, are linked in order to the FNP field of the symbol table slot for the function. While registers are assigned to these parameters, the corresponding formal parameter names cannot yet be entered since they are not known until the function DEF statement is found. Note that formals are assumed to match actual parameters by both order and quantity. There are no checks in BAX59 to insure this correctness. Only a run through the WBASIC interpreter will verify parameter correspondence.

When the function definition is found in the DEF statement, a process similar to parsing the call takes place. The formal parameter names are now inserted into the parameter list attached to the function slot in the symbol table. Before the function body is processed, the slot is pushed onto the FNSTACK. This stack simulates an activation record stack. Each identifier look up that is performed by procedure IDLOOKUP requires that the FNSTACK be examined for active functions. If a formal parameter name is found in an active function parameter list which matches the identifier being sought, then its register assignment is used for code generation. As a result, standard rules of variable visibility and scoping apply. When the end of a function body (FNEND statement) is encountered, the function slot is popped from the FNSTACK and its formal parameters are no longer visible to the run-time environment.

As a final note, the user should know that "parameter" function names receive their own register assignment. This register is the place in which the final value of the function is returned. This register is zeroed during run-time just prior to the execution of the function call. However, after execution the value in this register persists until the next call on the function. This corresponds to an identical situation in the WBASIC run-time environment.

9. Code Resolution

If the physical end of the WBASIC source program is reached, or if a WBASIC END statement is found, parsing is stopped, the bodies for any "long" functions used are generated, and the code data structure is closed out with nil pointers. At this point, the code resolution phase of compilation begins.

The first step in resolution is to locate and insert labels at the destinations of all unstructured subroutine (SBR) calls. These, of course, were generated by GOSUB statements. Since GOSUB is a line-oriented jump, then there is a pointer in the code data structure referencing the destination of that jump. Procedure FINDGOSUBLBL traverses the code data nodes searching for SBR keycodes which are followed by a node with a non-nil jump pointer (JMPP). A very complicated check is made to ensure that the SBR label has not already been inserted by an identical SBR call. If not, then the back pointer (BAKP) is used by procedure PUTGOSUBLBL to assist in the insertion of the label at the jump destination. Once the insertion has been completed, the address field (ADDR) of the JMPP target is set from zero to negative one and the jump pointer (JMPP) is set to nil. This signifies to other routines that this jump has been resolved. The process continues until the end of the code data is reached.

The next step in resolution is to perform a special brand of TI-59 "peephole" optimization. The most common forms of excess parentheses pairs are removed. Such forms as "(RCL nn)" and "(2.333E-12)" will have been generated as a result of parsing even simple assignment statements and expressions. Since the parentheses in these expressions are unnecessary and use up valuable program steps, they are removed, provided they are not referenced by a jump pointer. If referenced by a jump pointer, the node's address field (ADDR) value will be 0 instead of -1 or -2. Removal of these will cause dangling jump references.

Looping and branching generate many place holding NOP keycodes. These are also an unnecessary use of program steps. However, remember that almost all of these were generated to anchor or project jump pointers. Thus, before removal their jump pointers must be reset. Procedure OSQNOP

passes over the code data twice, once to reset all jumps to and from NCP's, and the second to locate and remove the NOP's. It is important to realize that there is a distinction between a useless NOP and one which is acting as a label identifier or a jump address place holder. Because the TI-59 requires that particular keycodes be followed by labels, register numbers, or addresses it is easy to check keyccode usage. This information is actually loaded during initialization into the UNIT field of the CODETEXT record. It is an integer 0..3 which indicates whether the TI-59 code node is a one, two, three, or four keystroke instruction. This information is used to pass over keycodes which are required parts of a larger instruction.

The final stage of resolution is to convert relative jump (pcinter referenced) addresses into absolute (numerical) addresses. This must be the last step because previous code insertion/deletion routines constantly change absolute addresses. At this point no code insertion or deletion occurs. Procedure RESOLVE_ADDR passes over the code data twice. The first pass fills the address fields (ADDR) of all code nodes in sequential order starting at 000. Now that each exact absolute address is known, all jump pcinters which are still marking address space and referencing a destination node can be resolved. A TI-59 coded address consists of two parts. During the second pass procedure INSERT_JMPADDR is called at non-nil jump pointers to read the destination address, split it into its two integer parts, and insert the parts into the address space nodes. Once all jumps have been resolved, the code data structure is ready for output and linking.

10. Input/Output

In this subsection we discuss two input/output related issues: Implementation of I/O constructs and OPTION messages to the compiler. The limited capabilities of the calculator required that file handling and string handling aspects of WBASIC be eliminated from our subset. For similar reasons the I/O constructs which could be translated from WEASIC required restrictions.

While the WEASIC I/O statements INPUT and PRINT normally provide for file management, the BAX59 implementation cannot. The cross-compiler recognizes PRINT followed by any number of simple expressions separated by commas. The TI-59 code generated will evaluate these expressions and print their values (to either the display register or the Texas Instruments PC-100 Printer Cradle). On the other hand, the INPUT statement takes any number of variable identifiers separated by commas. For each identifier in the INPUT list, the TI-59 program halts execution, displays the register assignment for that identifier, and stores the input value entered by the user in the register assigned.

Many programs require the reading of large amounts of data, often at the start of execution. In this situation the INPUT statement tends to generate an excessive amount of program step overhead. Unless the program is designed to be interactive, this overhead unnecessarily increases TI-59 program size. In order to provide a more space efficient means of data entry, a limited translation of the WEASIC DATA and READ statements was designed. In some sense, these statements provide a substitute for file handling. The DATA statements are placed at the beginning of the WBASIC source program. Each statement may be followed by numeric data items separated by commas. The total number of data items in one program is limited to the number of unreserved

registers available in the calculator (based upon the system parameter REGBASE). If this limit is exceeded, a warning message will be issued. READ statements take variable identifiers and may be written with the DATA statements, however, the number of variables input to READ statements should never exceed the number of data items provided by DATA statements. This condition will also cause a warning message and further DATA/READ statements will be ignored. The parse routines make register assignments to the variables in the READ statements, and concurrently build a list which maps the data items to their respective registers and variable names. This list is one form of compiler output. Using the list the user can pre-load TI-59 registers with numeric values and be assured that they will be in correspondence with the translation of variable names. More importantly, no TI-59 program steps are used for this initial input. In fact, the data could be read from a magnetic card into a memory bank prior to execution.

As we have previously implied, there are many forms of output which can be generated by the cross-compiler. Additionally, the user will probably have to do some debugging. We have chosen to provide a primitive set of tools and options which can be toggled on or off from outside the BAX59 source program. The toggles are set or reset by using the OPTICN statement in the WBASIC program. Caution! Do not confuse this statement, which is unique to the BAX59 cross-compiler, with the WBASIC OPTION statement. They are not the same. BAX59 does not recognize WBASIC OPTION parameters and WBASIC does not recognize BAX59 OPTION parameters. Table III lists the possible options available to BAX59 users. To toggle the options, simply include an OPTION statement as the first line of the program to be translated. Desired parameter settings should follow the OPTION reserved word separated by spaces. Positive parameters set the

TABLE III
BAX59 OPTION Statement Parameters

Parameter	Option	Default
+0	Generate linker interface file	false
+1	Generate code for PC-100 printer	true
+2	Optimize out unnecessary parentheses	true
+3	Optimize out unnecessary NOP's	true
+4	Translated TI-59 code to list file	true
+5	Image of symbol table to list file	false
+6	Contents of code structure to list file	false
+7	Each lexical token to terminal	false
+8	Each lexical token to list file	false

toggle true; negative parameters reset the toggle false. In the case of the zero parameter, the sign has no effect.

As a final note, an OPTION statement may not be placed in the WEASIC source program until it is ready for translation. Also, placing an OPTION statement in any line but the first may produce unpredictable results.

E. LINKER

The linker's purpose is to produce a segmented version of the compiled code and present the code in a format that is user friendly. The informal strategy used to accomplish this was discussed in the preliminary design phase of Chapter III. The detailed design that supported the solution strategy called for the linker to operate sequentially through three major phases. In Figure 4.5, the contour diagram for the linker is presented. The preprocessor phase of the linker includes actions from some of the SYSTEM UTILITY procedures and the BLD_SEGMENTBL procedure. The remaining two procedures, COALESCE and INSTRUCTIONS, accomplish the segmenting and postprocessing activities.

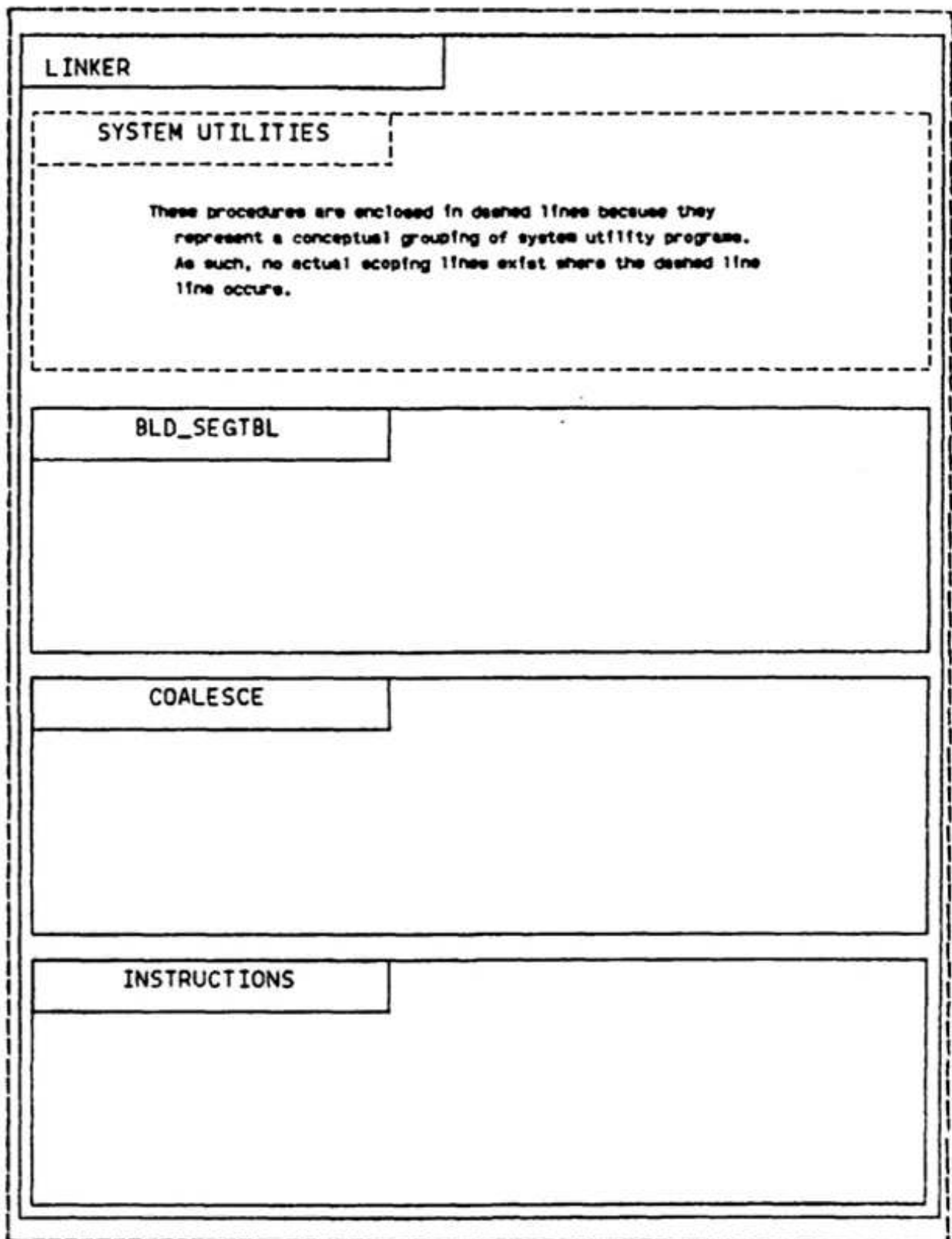


Figure 4.5 Linker Contour.

Each of these major actions were described in the preliminary design phase in Chapter III. The detailed design of these specific operations will be presented in the next sections. Only those major design considerations required for understanding the operation of the major operative phases will be presented.

1. PREPROCESSOR

As was mentioned in the preliminary design, the primary purpose of the preprocessor is to reproduce the compiled linked code list and generate a table that represents the sequential segments of the compiled code.

The informal strategy called for a two step operation. In the first step, textual integer pairs are read from an input file into a data record. Each record is linked to the preceding record forming a linked list which reproduces the linked list of code generated by the compiler. The next step evaluates the linked list to determine where the sequential segments are located. Information concerning each sequential segment is stored in another record and linked to the preceding sequential segment record, thus forming a linked list of sequential segment description records. Evaluation for sequential segments would occur by TI-59 labeled subroutines. Each list of sequential segment records would be pointed to by a header record which contains the subroutine name. Each of these subroutine name header records would be linked to other subroutine name header records in the same order in which they were detected in the generated code.

Two data structures were needed to support this strategy. The first structure comprises a linked list of records. Each record contains all the information that is contained in one program step in the TI-59 calculator, including the address of the instruction and the instruction

integer code. Each record is linked to the following record by a dynamic pointer, which captures the sequential nature of the compiled code. Another dynamic link is provided for those records containing keycodes that may cause the flow of control to change from a sequential flow. The generated linked list of records is a complete internal representation of the compiled code.

A second data structure is needed to represent a sequential segment of code. Vital program control flow information must be captured by the structure so that segmentation rules may be applied during linker processing. To accomplish this a sequential segment table was developed utilizing a record format to describe each segment. This table record holds data such as segment start address, stop address, whether the segment is covered by an iterative backloop, a list of forward jumps and a list of subroutine invocations that originate within the segment. Each one of these records is linked to the following sequential segment's record. In addition, the sequential segment records are grouped according to subroutine. That is to say, only those sequential segments residing within one TI-59 subroutine definition are connected together in sequential order.

The linked sequential segments are tied together by other records of the same basic type but different variants. Each subroutine grouping of sequential segment records is pointed to by a linked list of header nodes. These nodes contain the name of the subroutine and the subroutine definition address. Each header is linked to another subroutine header in the same order in which subroutine definitions occur within the generated TI-59 code.

To capture information relating to forward jumps, a variant of the sequential segment record is used to build a forward jump list. This list contains the originating

address of the forward jump and the address of the instruction to which control is transferred. Because the actual jump address is used, the link to the jump location is termed relative. Each jump node is dynamically linked to following jump nodes to form a jump list. This list is, in turn, dynamically pointed to by the sequential segment in which the jumps originate.

To capture information regarding subroutine invocations, the same type of structures is used as for the jump node lists. The only difference is that the subroutine lists point to the invoked subroutine in a dynamic manner. That is to say, that a dynamic pointer is set to the first sequential segment of the invoked routine in the sequential segment table. This is basically the only difference between the subroutine invoke list and the forward jump list.

Figure 4.6 is the contour diagram of a conceptual grouping of procedures referred to as the SYSTEM UTILITIES group. These procedures are not explicitly grouped together by code; rather, the grouping is to facilitate discussion and understanding. There are several operations within this group which manipulate the data objects.

In creating the linked compiled code list of records two separate procedures are used. The first procedure is called INPUT. This procedure builds the initial linked structure. It utilizes an input file containing the integer pairs representing TI-59 code steps. Essentially it creates one record for each pair and links the previous record to the new record. The only thing not done is the setting of pointers to represent an indirection in the flow of control.

This is the job of the procedure SET_JMPS. In this procedure, the major activity is the detection in the actual keycode portion of the compiled code of an instruction that represents a possible change in control flow. When one is

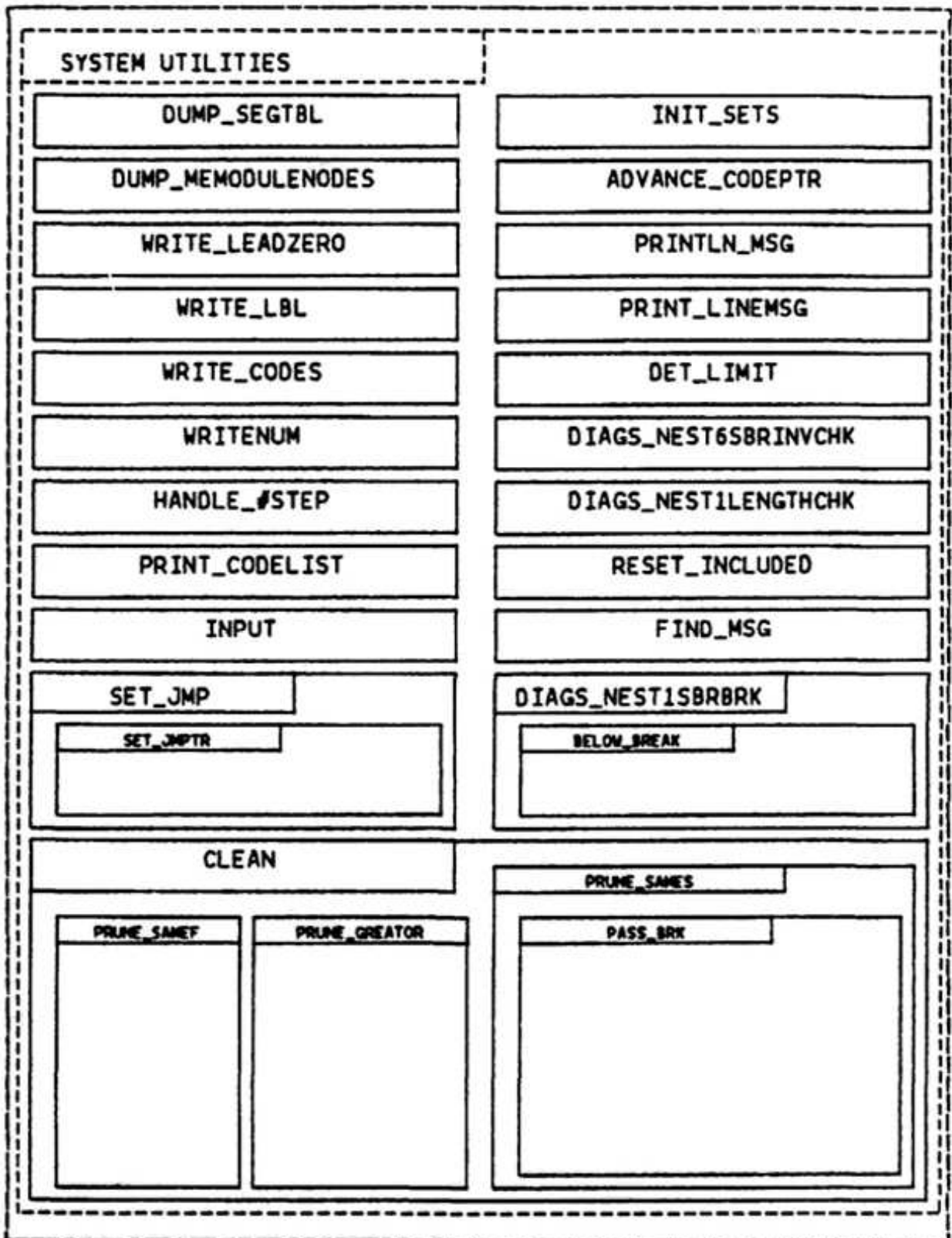


Figure 4.6 System Utilities Contour.

detected, a jump pointer (the indirection pointer) is set pointing to the record containing the next TI-59 program step to be executed.

The operations which create the second data object are a little more complicated and are contained in Figure 4.7. The action of building the sequential segment table data structure is broken down into three steps. The first step begins the formation of the table while the second step completes segment detection. The last step captures other information and ensures that internal interfacing requirements are met.

The first part of this procedure is accomplished through the BLD_PRIMSEG_TBL procedure. This operation passes over the compiled codelist structure and determines where a subroutine starts, stops, or issues a back jump command, and locates the terminal points of the back jump commands. Each of these points is called a critical point. When detected, each critical point is inserted in the segment table data structure under the header node containing the TI-59 subroutine code name which is being processed. In addition, each of the jump commands with their initiation and termination points are inserted into the structure. This completes the first major step.

The second operation is accomplished through the BLD_ADVSEG_TBL procedure. In this procedure the initialized data structure is fleshed out. Up to now only critical points have been inserted. As these are points and are not double ended, segments have not been delineated. This procedure examines the segment data structure and adds points to delineate where a segment starts and ends. It does this by subtracting one from the point following it and taking this to be its end point. This results in a series of records which are all covered by an iterative backloop, with the exception of the first record. This is noted in

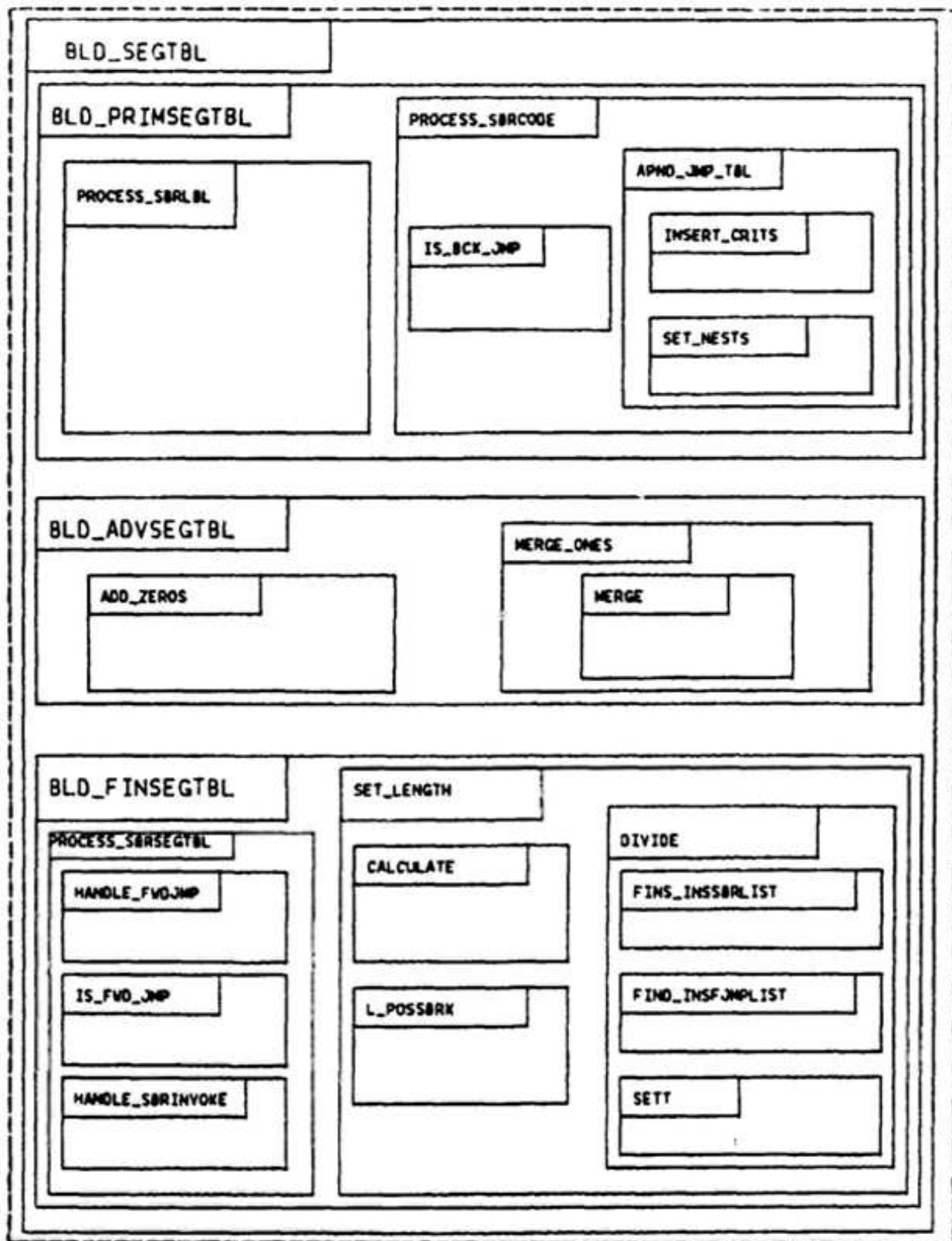


Figure 4.7 BLD_SEG_TBL Contour.

the record. Next, the procedure examines the modified structure and determines by examining the addresses where there are holes in the table. These holes correspond to sequential segments that are not covered by an iterative backloop. These records are then inserted into the structure. Lastly, adjacent segments that are covered are merged to form one record representing a sequential segment that is covered by the largest iterative backloop.

The segment table structure is completed in the last procedure, called EID_FINSEGTBL. In this operation two primary things happen. First, PROCESS_SBRSEGTBL evaluates the compiled code list and determines where forward jumps and subroutine invocations occur. It inserts these locations into the proper sequential segment that covers the area where the call or invocation occurs. Second, SETT_LENGTH checks that each sequential segment does not violate the memory size limit of the calculator. It does this by checking each sequential segment record and calculating a size. If the size is too great, then the segment is divided in half and a new segment is inserted into the table. This is not done for segments that are covered by an iterative loop as this would represent a break of an iterative loop. Other actions that must occur include readjusting forward jump lists and subroutine invocation lists if a division does occur. One interesting point worth noting is that when the length check is made additional steps must be allocated to the actual length to compensate for the possibility of prompting code being added for an invocation to a subroutine that does not currently reside in memory. This is the purpose of L_FCSSBRK in Figure 4.7.

The data structure operational procedures access the data structures through pointers which point to the structures. The pointer to the compiled code list is referred to as BUILT_CODE. The pointer to the sequential segment table

structure is called SEGTBL. These are the only data which are passed among procedures. One point to remember is that SEGTBL points to the header node list containing the names of the subroutines. The actual sequential segment lists reside underneath the header nodes.

Since understanding the data structure and its construction is essential to understanding the remainder of the linker, an example will be examined to demonstrate the preceding sections.

In Figure 4.8 a sample topology of a TI-59 program is given. It includes four subroutines of various sizes and with various control flow indirections. In looking at the diagram it is important to note the absolute address locations given, for these will be critical to understanding the development of the sequential tables.

As was mentioned, the first operation is the restructuring of the generated TI-59 code. Figure 4.8 represents approximately this structuring. The actual code line is rebuilt internally in the machine and is pointed to by pointer EUILT_CODE.

Figure 4.9 is the completed sequential segment table, without the linked header node list. To understand the concept of sequential segment a comparative look at Figures 4.8 and 4.9 must be made. In Figure 4.9 the first sequential segment is defined as being between addresses 000 and 049. This is reflected in Figure 4.8. When looking at the sequential record one sees that the forward jump information is captured in the forward jump list node which, in this case, is only one node long. When looking at the second sequential segment one notes that there is a nested back jump. The sequential segment is defined to be that segment which is covered completely by back jumps. In this case it extends from 050 to 199. If for some reason the back jumps shown in Figure 4.8 did not fully contain each

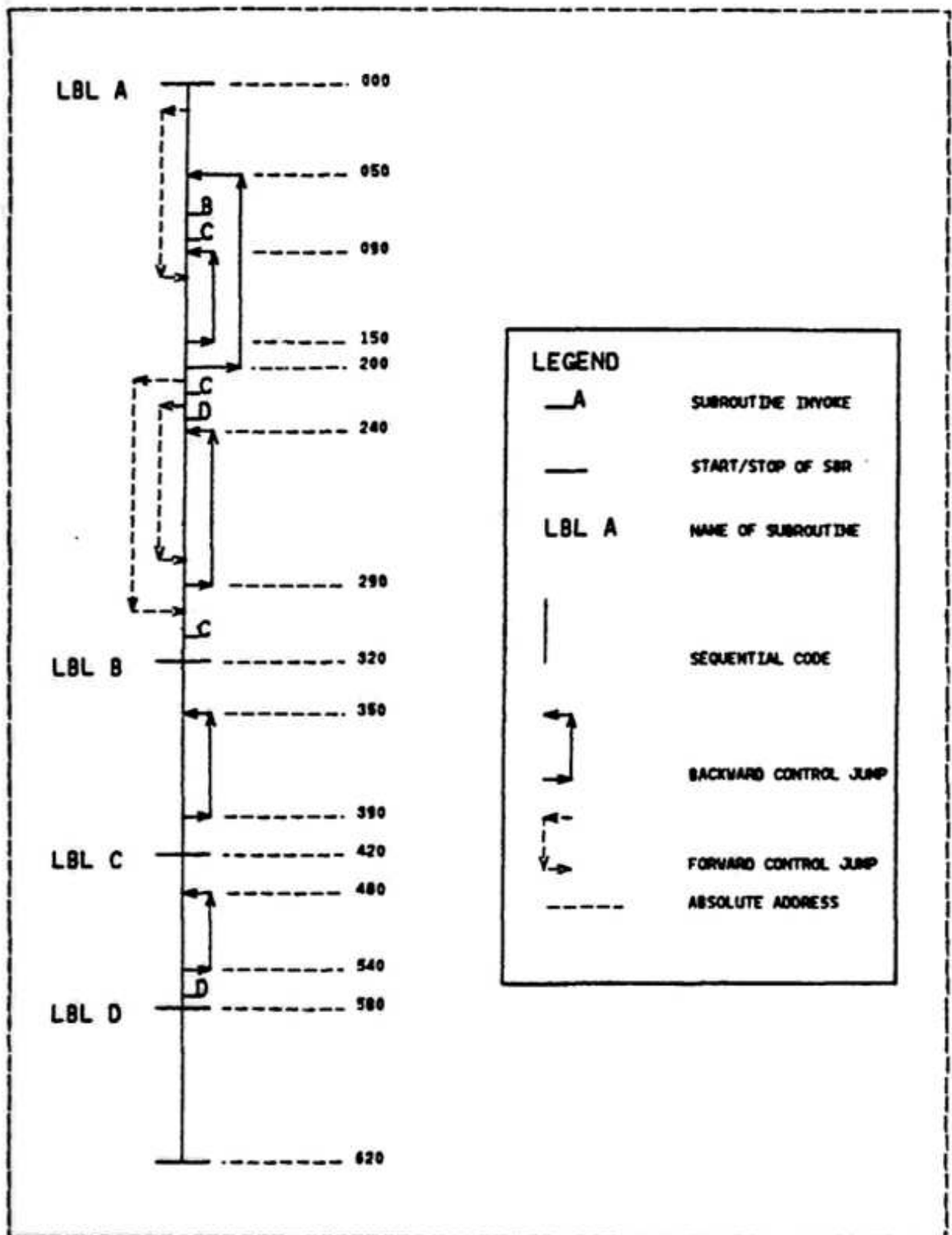


Figure 4.8 TI-59 Code.

other, that is to say, one jump started at 199 and stopped at 090 and the other started at 150 and stopped at 050, then the cover would still extend from 050 to 199. The reason is that this region of code is probably caught in an iterative loop and cannot under any circumstances sustain a break within this cover.

Another point to be made is the manner in which each subroutine's sequential segments are recorded together. In addition each invocation is recorded as is each forward jump. During the completion of the table, invocations to the same routines from different invocation locations are deleted, thus leaving only one link to the called routine for that sequential segment.

A final point concerns the recursive nature of the structure. By assuming that the first subroutine is the main routine and that all other lower level routines are below it (in the sense that they are pointed to from invocation nodes) one can see that any routine used to combine segments can be used on any subroutine's sequential segments. This opens the door for recursion to be used in a bottom up recombination scheme to be discussed later.

Many problems were encountered in the development of the preprocessor phase of the linker. Only the most difficult or annoying will be discussed.

One of the first problems concerned the multiple meaning of program steps in the generated TI-59 code. A separate TI-59 program step may be either a command, register number, flag number, or part of an address. The meaning is dependent on the last valid command. Commands can affect the interpretation of a program step as far as three step positions away (analogous to the concept of one-byte, two-byte, and three-byte instructions in assembly code). This had to be taken into account when doing any operation requiring an interpretation of the code. This

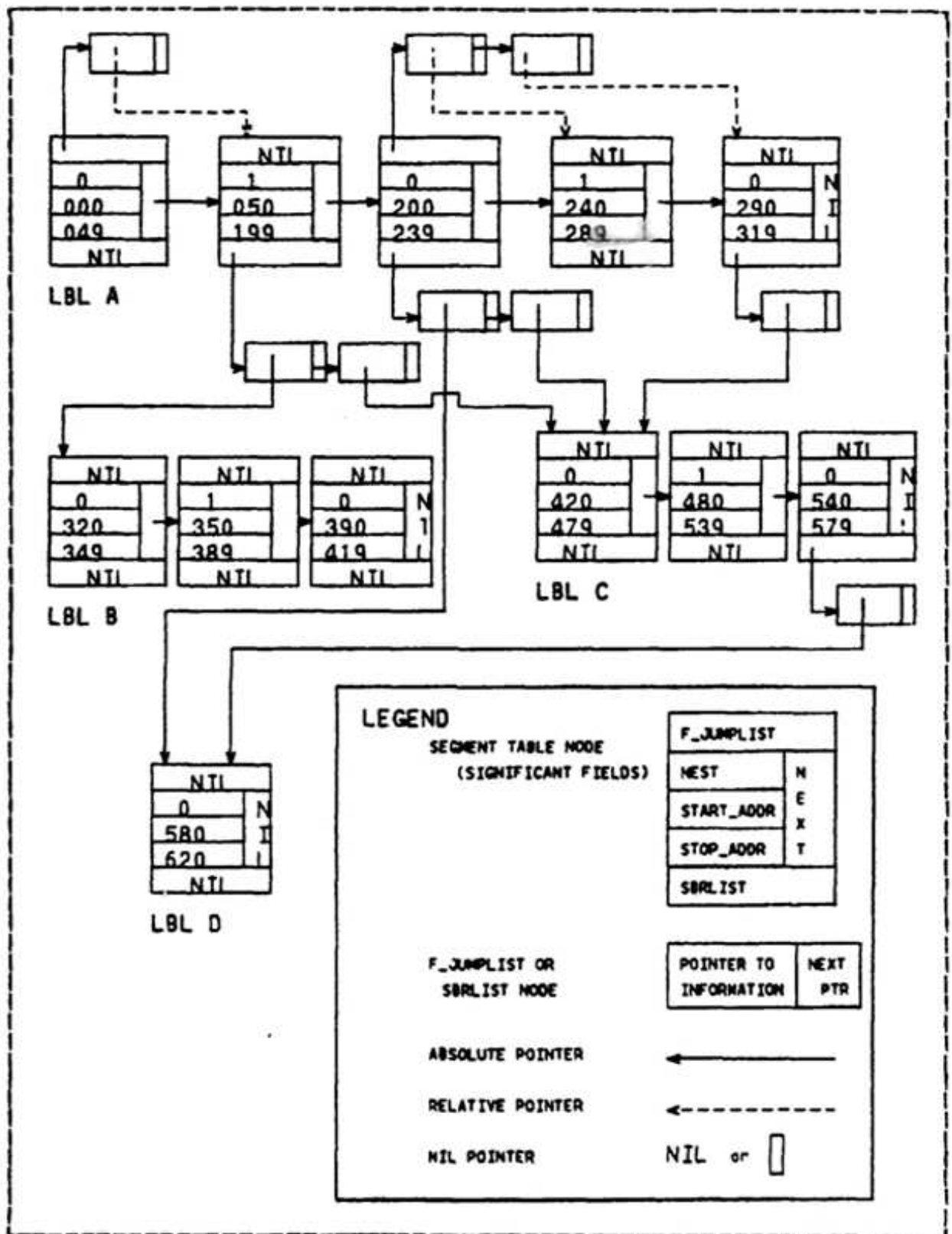


Figure 4.9 Sequential Segment Table.

resulted in special code sets being initialized and special routines being written to print out labels and move the compiled code list pointers. All of these are shown in Figure 4.6.

Improvements in this operative phase could be realized. In the early stages it was decided to separate the compiled code lists from the segment table lists. This was to avoid accidental tampering with the compiled code, since the integrity of the compiled code was the paramount consideration. It would be possible to make the compiled code lists a variant of the segment table. Then, instead of having relative pointer indexes to compiled code addresses, an absolute pointer could be used. This may reduce the size of the program significantly in that types would now be compatible and a reduction in the number of output routines due to the different types would be realized.

2. Segmentor

After the input file has been preprocessed then the linker passes into the segmenting phase of the operation. The routines that support this phase are built into the Pascal procedure called COALESCE depicted in Figure 4.10.

The informal strategy called for the sequential segments of a subroutine to be combined to form larger sequential segments. This recombination would be allowed as long as memory limits were not violated. This required that invoked subroutines be combined first before the caller so as to make room for the invoked routine's code. If the invoked routine could not reside then a break was placed in the dynamic link to the invoked routine and prompting code added to the caller's length for memory size checking purposes.

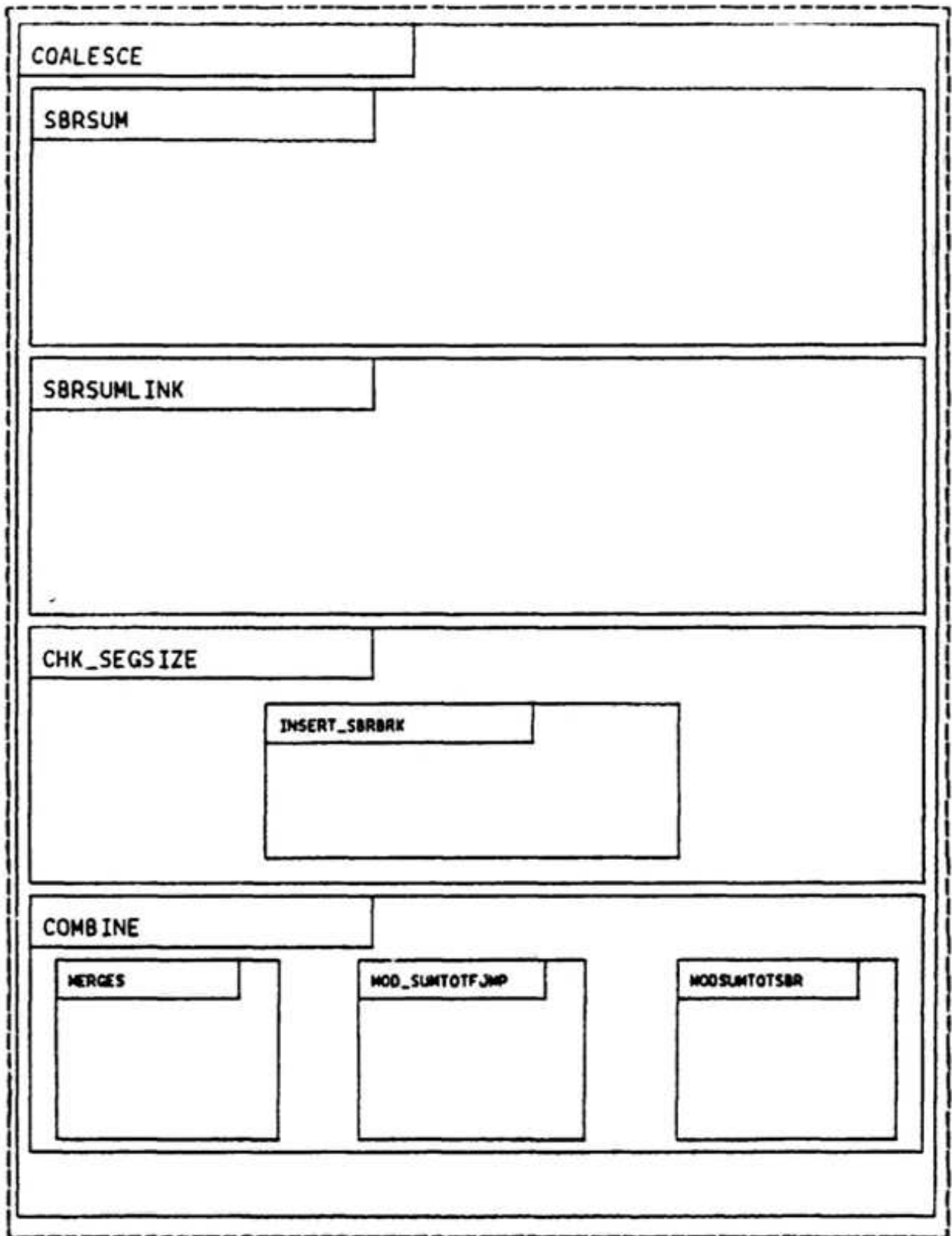


Figure 4.10 COALESCE Contour.

This strategy has one major requirement: the invoked routine must be recombined before the caller can be recombined. For this reason a recursive solution was adopted. In this solution, the main subroutine is recombined. The first part of the recombination process is to ensure that invoked subroutines will reside with the calling sequential segment. If a subroutine is encountered that is not combined or coalesced, then the program will recurse on the new routine. Recursion will close out upon completion of coalescing of a particular routine. When all the sequential segments have been checked then adjacent segments are combined.

Another part of the strategy calls for the combination process to stop when a size limit is encountered. When this happens then some sort of break notation must be used to mark where the limit was exceeded so as to prevent production of code segments that exceed the memory capabilities of the calculator. After the break has been set then the process of recombination begins on the other side of the break with the non-combined segments, starting again with a memory limit of zero.

This process of breaking and checking limits results in the sequential segment table containing break points. These break points delineate the exact locations where program segmentation will occur. These points will mark those portions of code which can fit in the calculator memory according to the rules of segmentation outlined in the preliminary design.

The data structure that supports this strategy is the sequential segment table. No other structure is used. The only addition to the structure is the node referred to as a subroutine invocation break node. This node is inserted between a subroutine invocation list node and the invoked subroutine sequential segment table. All other

changes to the structure involve removing nodes and combining adjacent information into one node.

There are two major activities that support the above strategy. The first activity is the checking of a segment and the second activity is the combining of adjacent segments. Overseeing these activities is a single driver. This topology was suggested by the recursive nature of the solution. The procedures which support these activities are shown in Figure 4.10.

The driver is represented by the Pascal procedure COALESCE. This routine is called whenever a new subroutine is encountered that has not been combined or coalesced. The interior Pascal procedure CHK_SEGSIZE verifies that the specific segment it is looking at, together with all called subroutines on the subroutine invocation list for that segment, will reside in calculator memory. This routine uses SBRSUM and SBRSUMLINK to determine the lengths of invoked routines. It recurses mutually by calling COALESCE in the event that an invoked routine has not yet been coalesced. It determines this by looking at a boolean field in the segment table. This field is set true if the subroutine has been coalesced. The other procedure, COMBINE, accomplishes the actual combination of adjacent sequential segments. It uses the length predictor routines MOD_SUMTCTF_JMP and MOD_SUMTOTSBR to predict a combined length which takes into account any changes that might occur in the subroutine invocation or forward jump lists. If the combined length is within limits then a recombination occurs; if not, pointers are advanced. This means that any sequential segment records which follow the initial sequential segment records are part of a new memory calculation. In other words, any sequential segment links that are not nil represent a break between the linked sequential segments. Upon exiting COALESCE, the subroutine that has

just been coalesced is marked as such in the sequential record's bcclean field reserved for that information.

To visualize the result of the segmenting phase another look at the example is provided. Figure 4.11 represents a segmentation based on a memory limit of 550 steps. Note that each of the invoked subroutines has been coalesced into a single sequential segment. Also note that a break was made in the main subroutine. This is shown by the fact that the main routine is not a single segment record. By examining the table it can be seen that the routine labelled "C" will be copied twice when the two memory sized segments are produced.

To interface between modules in this recursive environment several things were assumed or used. The first was that the data structure would serve as the repository of global data. In addition, a variable would be used to keep track of the current size of the combining memory program steps. This variable was passed as a parameter in order to preserve its value throughout the recursion. All pointers were passed as local parameters. This preserved locations in the data structure as the algorithm progressed through the different levels of recursion.

These operations did require some other work in order to obtain valid data that would correctly calculate code lengths to include multiple copies of subroutines. The problem occurred when there were multiple invocations to the same subroutine from different segments (or even different subroutines) that up to now were all included in the same memory limit calculation. To solve this another field was placed in the segment record to indicate whether or not the particular routine had been included or not in length calculations. Whenever a sum was calculated and a routine included then the field was set true. Whenever a new memory limit was reset back to zero following the implantation of a

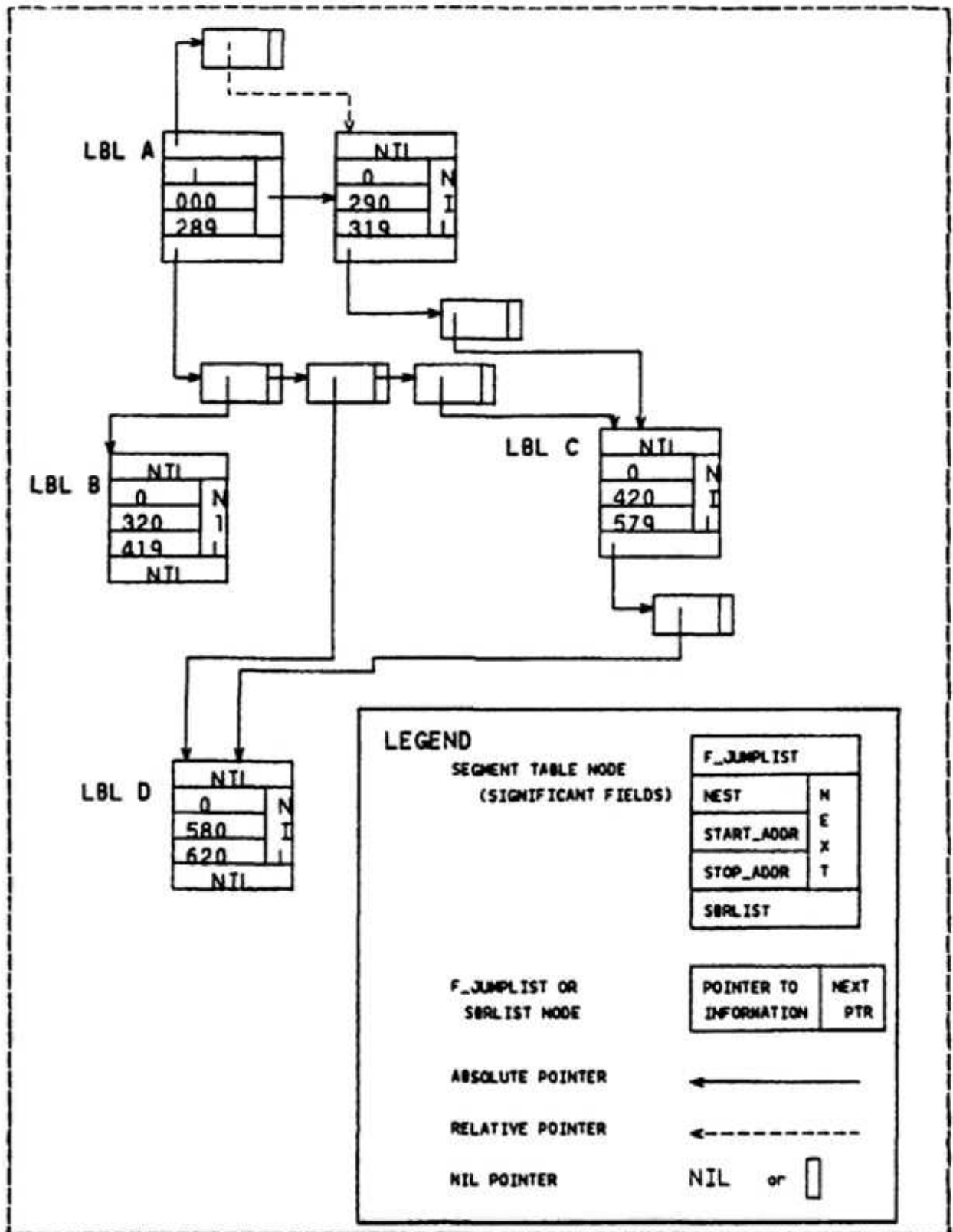


Figure 4.11 Coalesced Segment Table.

break point, a SYSTEM UTILITY procedure was used to reset all the included fields back to false. This means that only one copy of a subroutine would be considered for each calculator sized memory computation.

Future implementations should develop a better method for recording whether a segment is coalesced or included. The inclusion of this field in the segment table record was a "quick fix." This fix results in wasted storage as it is only used in the first record for each subroutine. An improvement would be to use another variant record to record all current data, with the exception of coalesced and included information, for all other sequential segment nodes other than the first sequential segment node. This would save memory.

3. Ecst Processor

After the segmentation phase, the linker passes into the `ecstprocessor` phase. It is this phase that provides the required output for the user.

The informal strategy divides this phase into three distinct operations. The first operation designates the start of each calculator sized segment of code. These segments, which meet the memory requirements of the calculator, are referred to as memory modules. The second operation copies the required code into each segment and inserts any segmentation prompting instructions that are needed for successful code execution. The last operation consists of outputting the segmented code in a user friendly instruction sheet format. This completes all linker actions.

In order to support the informal strategy, several data structures are used, two of which were described earlier. They are the segment table and the compiled code list. At this point the segment table has been coalesced

and contains the locations of the segmentation breaks that will minimize card reads. The compiled code will be copied by segment as delineated in the segment table. A third structure is built in this phase and a fourth structure is provided with the program.

The third structure is referred to as the memory module data structure. It is a Pascal variant of the segment table, which allows compatible pointer references between the two objects. The structure consists of a linked list of head nodes, which are named by respective memory module number. They represent one calculator's worth of available memory programming steps as determined by the calculator partition. Each node of this linked list points to two locations. One location is to the first sequential segment table record node following a segmentation break. The second link is to the copied code that will make up the programming steps of the memory module. In Figure 4.11, there would be two memory module header nodes. The first header node always points to the first sequential segment of the main subroutine. In the example this first memory module would be pointing to the node beginning with address 000. The second memory module node would point to a record that follows a break. This would be to the sequential segment node beginning with address 290. Just as there are no other breaks, there are no other memory module nodes. The other pointers would point to a linked list of code.

The copied compiled code list is a part of the memory module data structure. It is another Pascal variant of the same record type. This list is similar in structure to the compiled code list reproduced during the preprocessor phase. The only difference is in type. Another difference is that there are no jump pointers or dynamic pointers indicating a change in flow of control. The structure is just a linked list of sequential code. This structure is pointed

to by the memory module header node. Another point to be made is that the linked code list, when completed, does contain other code that is needed for prompting. As such it is not a one for one copy of the compiled code list. Lastly, a look at Figure 4.11 will show exactly the segments of code that can be expected to form the two memory module structures. By looking at the sequential segment nodes and following their dynamic pointers of the subroutine lists all required code start and stop addresses are given. It is this "look down" facility of the sequential segment table that make it so useful.

The fourth data object is provided with the linker program. It is a textual file which contains text messages which are used by the linker. Each message is delineated by a \$XXX where XXX is an integer. The linker, when provided the number portion of a message, can easily locate the message. Once located it can either extract values or copy the message verbatim to an output file. This is what occurs during the formatting of the instruction messages.

The operations that build and manipulate the data structures function in three phases. Figure 4.12 is the contour diagram of the subroutine that supports these operations.

The first operation is the construction of the header memory module nodes. This is accomplished by the Pascal procedure BLD_MEMMODULENODES depicted in Figure 4.12. This procedure traverses the segment table and looks for break points. When it finds one, it checks to see if the break has already been detected. If it has not been detected then it builds the header node and assigns it a memory module number. The reason for the check is that the traversing mechanism is based on recursion. In this strategy, traversing is begun with the main subroutine. In Figure 4.11 this would correspond to subroutine LBL A, node

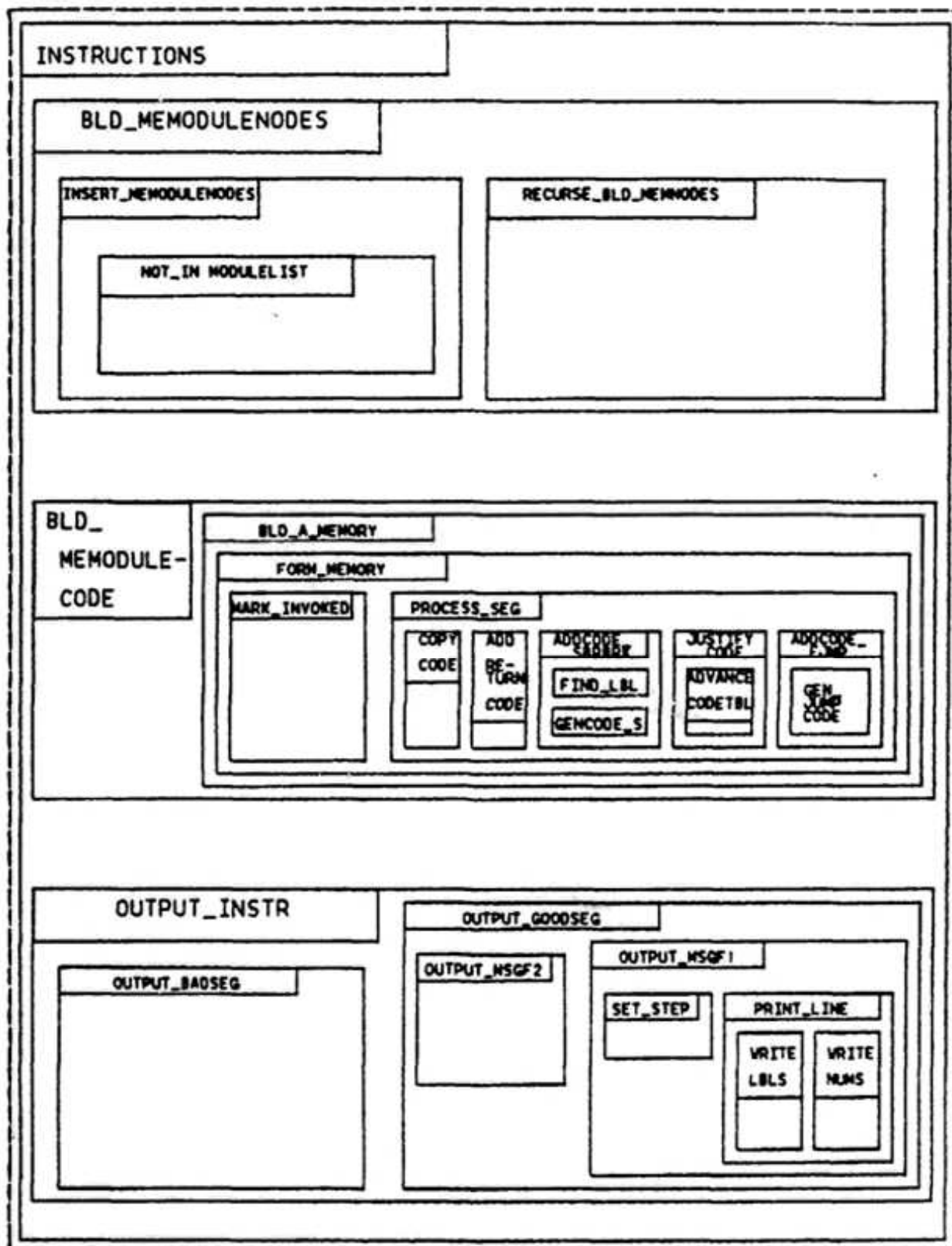


Figure 4.12 INSTRUCTIONS Contour.

start address 000. It then searches right along the same subroutine to detect all breaks of that subroutine. Then it resets back to the start of LBL A and begins to check the subroutine list of each node that comprises LBL A. Recursion is implemented at this point when the subroutine link is traversed and another subroutine is discovered. If a break is discovered in the subroutine list then another memory header node is built and the program bypasses the break and recurses on the next subroutine. This traversal mechanism leads to multiple discoveries of the same breaks. Consequently, the check is made to ensure multiple copies are not placed in the memory module header list.

The next operation consists of copying code from the compiled list, resetting address key codes for jumps and adding prompting code to each specific memory module code list. Figure 4.12 contains the Pascal procedure BLD_MEMODULECODE which accomplishes the above tasks. This is done by moving down the memory module header list in a sequential fashion. At each header, the link to the segment table is traversed to determine exactly what segments of the compiled code are to be copied. This is the duty of BLD_A_MEMORY in Figure 4.12. Once the start and stop points are determined and copied then recursion is utilized to traverse the subroutine links of the sequential segment table to obtain the required copies of resident subroutines to support the functioning of higher level segments. During this operation the segment table is used as a "check pad," that is, copies are marked included after being copied and are reset upon completion of copying. Another function accomplished during the processing of a memory module segment is the addition of prompting code. Lastly, addresses are reset and justified to include the resetting of jump address key codes to reflect new jumps to internal prompting messages and the absolute address of the originally compiled code list.

Once the memory module data structure is completed then the structure is presented to the user in an instruction type of format. This is the purpose of the Pascal procedure OUTPUT_INSTR. This procedure utilizes two data structures. It uses the provided message file structure and the memory module structure. The first action is to output the instruction introduction. This the procedure does through the use of the message file and the SYSTEM UTILITY programs FIND_MSG, PRINTLN_MSG and PRINT_LINEMSG. These are depicted in Figure 4.6. These procedures allow the linker to copy verbatim messages in preformatted form. Once this is done the procedure copies the codelist from each of the memory module lists of code. Once a specific module is copied the driver routine, OUTPUT_MSGP1, prints out specific information to delineate each memory module. After this action is accomplished, the procedure traverses the segment table and prints out additional user information that will aid the user in the execution of his program.

Interfaces between modules are accomplished as usual with pointers. These pointers point to their respective data structures. Global information is recorded in the data structures or in special global variables which are passed as parameters during recursive operations.

One major problem that was encountered and solved in an interesting manner concerns the formatting of the output. The vast amount of instructional information that was required to be output made inclusion in the source code ridiculous. To solve this, the message file system was developed. This system consists of a text file containing preformatted messages and a several procedures located in the SYSTEM UTILITIES contour in Figure 4.6. Each message is delineated by a "\$" and a number. Two types of messages can be processed. One kind of message results in a complete copy from the first line following the message code (\$XXX)

down to the line preceding the next "\$" encountered. The other messages are one-line messages which copied until the "\$" at the end of the message. This gives the programmer the capability to write out blocks of text and to write out text and computer generated information on the same line.

Another capability provided by the package is the ability to search out messages from other files. The procedure FIND_MSG takes as parameters a file as well as a message number. This facility allowed the linker to be loosely coupled with the cross-compiler by interfacing with a message number coded file produced by the cross-compiler. All that the linker needed to know was under which number a required piece of information was coded and the interface file name to affect an interface.

An improvement might be realized in the output of the generated code list. Currently there are two separate sets of procedures used to output code lists. This was primarily due to typing differences. However, the second set of print procedures located in INSTRUCTIONS (see Figure 4.12) is probably more efficient. Furthermore, if the reconstructed code were changed to be a variant of the segment table structure then a reduction in Pascal code lines would be realized through the elimination of a code list group of printing procedures. A further increase in efficiency may be realized in any operations requiring use of the reconstructed compiled code list.

C. INTERFACE ENGINEERING

In any detailed design, careful consideration must be given to interfacing criteria. Interfacing criteria should be as explicit as possible, however this is not always possible. Sometimes, design decisions or engineering interpretations have implications that affect other modules or

submodules. These are generally of an indirect nature in that the interface is implied in the system structure and is not explicitly passed from module to module.

These types of interfaces surface in the detailed design phase. Decisions regarding TI-59 address labelling and structure of the TI-59 subroutine greatly affected the design. In addition, assumptions about the use of structured programming and the prohibition of recursive WEASIC programs facilitated system design. Simple redefinition of the use of WEASIC commands READ/DATA provided an easy form of I/C, but again was an implied interface. These types of implied interfaces will be examined in the following sections since they are critical to understanding the system operations and to future maintenance.

1. Addressing TI-59 SBR

One implied interface resulted from a decision on the mechanism of subroutine invocation which would be used by the system. This decision arose from the fact that the TI-59 calculator may invoke subroutine code in several different ways.

To understand why a decision was required a look at the subroutine naming conventions and procedure for invocation are in order. A subroutine name is composed of two program steps. The first step is the keycode 76. This is the LBL code. It tells the calculator that the next key is a subroutine name. The next program step is the actual subroutine name. The keys which may serve as actual subroutine names may be one of two types. The first type comes from keys which are undefined. That is to say the keys are not used by the calculator to perform calculator functions; they are strictly reserved for naming subroutines. The other type of name comes from defined keys. By this we mean that the calculator uses the keys in some fashion in

addition to naming subroutines. These keys are overloaded: they can have two meanings.

To define which meaning is to be interpreted the calculator requires that a subroutine definition be preceded with a special key called the label key. This alerts the calculator to the fact that the next program step is a subroutine name. To invoke a subroutine, the calculator requires double meaning labels to be preceded with a special key called the invocation key. This alerts the calculator to the double meaning much as the label key. For undefined keys, this alert key is not required though its use will not alter any transfers. There exists another method of invoking subroutines. This calls for the invocation alert key to be followed by an absolute address. It is the duality of meanings which presented problems.

To overcome these problems two decisions were made. The first decision required that all undefined keys be treated as if they were defined keys. This resulted in only one case to be developed to handle subroutine naming. The penalty for this decision is that an added step was generated whenever an undefined key was used as a label and a call to that label was made. The other decision disallowed the use of absolute addresses in the subroutine invocation. All invocations would use labels. This decision carried a penalty in terms of execution speed. The calculator must search program step memory to locate named subroutines whereas address references can be reached directly. On the other hand there are several benefits. One program step was saved since a label requires only one step whereas an absolute address requires three. The other benefit permitted the definition of the invocation alert key as a two step instruction and not a two or three step instruction. This was the primary reason for this implementation. All of the system was designed with this decision in mind.

2. Structured Subroutines

A problem arose with the defining of subroutines. The calculator permits subroutines to be defined within other subroutines. Though this in itself is not extraordinary, the fact that nested definitions may be closed out with the same subroutine return key is not usually permitted. This type of structure made subroutine detection for segmentation purposes very difficult.

To solve this problem it was decided to allow only structured subroutine definitions. That is to say, only one return was permitted for each label. In addition, it was decided to disallow nested definitions. In fact, the decision was made to require that the programmer position his main program first in WBASIC source code and to have all of his subroutines follow in the manner described above.

This interface design decision forced a specific program structure. This structure was easy to detect, easy to compile and easy to segment. These benefits were realized at a cost of not being able to generate efficient or "tricky" code and of reducing the programmer's leeway in developing his WBASIC source code program structure.

3. Recursion

Although TI-59 calculator instructions tend to be evasive on the subject, recursive programs can be executed in a few special situations. Some BASIC languages support recursion, others do not. Although the WBASIC language supports recursion, limitations imposed by the calculator forced us to disallow the translation of recursive source programs. Other reasons for this decision involved complexities which such programs would present to the linker.

4. Input/Output

The development of the input/output structure was designed around the three major limitations of the calculator: programming steps available, storage registers available, and the calculator numeric display. In order to develop an efficient system to allow for input and output, restrictions were placed on some WBASIC commands.

Since the calculator display allows only numeric input and output, then any information passed between a human operator and the calculator must be in numeric form. Prompts used to communicate with the human must be imbedded in the compiled code. In the case of input and output, these prompts have an overhead in that they use up valuable programming steps. For example, with each INPUT command in WBASIC, a total of seven steps are generated to produce a prompt and store a value in the calculator. The problem with this occurs when large numbers of variables need to be input. If 60 variables are required for input then the INPUT command will generate 420 steps. This is clearly unacceptable.

To solve this problem the semantics of the WBASIC commands READ and DATA were modified. Use of these commands within a source program does not increase the number of program steps in the translation. A table of WBASIC variable names assigned to TI-59 registers is produced and placed in the interface file. This allows the human to input all of his data prior to execution without having to pay the penalty of generating extra prompting code.

Another decision concerns the location of the DATA and READ statements: All DATA/READ pairs must occur together at the beginning of the program. The reason for this should be clear. It would be impossible to place the commands in the middle of the program because they would

have no effect. With the exception of a single NOP (which is eliminated during peephole optimization), these commands generate no code. As a result, no run-time modification of variables can occur. Furthermore, DATA/READ statements placed within loops would invalidate the DATA to READ mapping, a static table.

The whole purpose of the redefinition of these commands was to give the user an optional form of I/O. This was expected to increase the efficiency of the generated code. These decisions influenced much of the design of the system.

V. TESTING

The test program and its generated output are provided in Appendices K through P of this report. The test program was chosen for several reasons. The first reason was to test the actual ability of the system to produce a usable TI-59 coded program. The second reason was to attempt to obtain an idea as to the efficiency of the system-generated code. The following sections present a description of the test program and comment on the efficiency of the generated code.

A. TEST PROGRAM DESCRIPTION

In developing a test program several considerations had to be taken into account. The first consideration required that the generated code be verifiable. To verify the generated code, it was decided to program a solution to a problem for which verifiable solutions existed. Verification would be achieved if the system-generated solutions matched those of the existing solutions for the same inputs. The second consideration was to attempt to arrive at some sort of efficiency comparison between the system generated TI-59 program and an optimized hand coded TI-59 program.

The test problem which was selected fit the above criteria. First of all, solutions to known input values existed. This ensured proper verification of the generated code. Secondly, a highly optimized hand coded solution to the problem existed for comparison. The problem selected is called the "Gunnery Problem."

The gunnery problem is to determine firing data for a howitzer cannon. It consists of inputting the following data: piece location, target location, piece corrections and howitzer ballistic coefficients. The output which is generated consists of time to target, elevation to achieve target hit and the lateral deflection (an angular measurement) to align with the target.

The solution involves calculating a range to target as well as an azimuth to the target. The azimuth is then converted to a lateral deflection, which is understood by the piece to be the correct azimuth on which to align. Next, three quadratic equations are solved to determine elevation, time of flight and shell drift. These are applied to the lateral deflection and to the piece elevation. A decision based on the calculated range to target is needed to ensure correct ballistic coefficients are used in the computation. These coefficients are based on the charge which is to be fired to achieve the range.

The hand coded version solution has been in use since 1976. The accuracy of this version was checked by artillery ballistic tables and by the Field Artillery Digital Computer (FADAC), the recognized source of all correct firing data. By using this solution, vast quantities of test input and output were available for program verification runs.

The hand developed version is highly optimized. For example, the hand version stores eight numbers in four storage registers. This is accomplished by storing the numbers on either side of the decimal point in a real number. This real number is then decomposed in a subroutine to obtain the correct number. In addition, this version makes use of calculator commands that were not implemented by the FAX59. For example, there is great use of the indirect store, decrement and skip on zero and polar to rectangular commands. All of these features made the hand coded version a highly optimized program.

B. TEST COMMENTS

The WEASIC solution was coded using basically the same program structure as the hand coded solution. This source code is presented in Appendix K. The interface file generated by the cross-compiler is presented in Appendix O. The final output generated by the linker is presented in Appendix F.

In verifying the accuracy of the generated code, many runs were made using input data for which known solutions existed. There were no deviations from the known solutions in any of the test runs. The conclusion is that the generated code was accurate.

In comparing efficiency a common unit of measurement was needed for comparison. This common unit was chosen to be the TI-59 program step. The reason for this is simple. Both registers and program steps reside in the same memory. Registers occupy eight programming steps. To measure how much memory a program uses it was only necessary to multiply the number of storage registers times eight and add it to the number of program steps to arrive at a figure which measured memory usage within the calculator.

This was done with the BAX59 generated code and the hand coded programs. The hand coded solution used 441 steps and 60 registers for a total step count of 921. The BAX59 generated code used 652 steps and 86 registers for a total step count of 1340. This represented an increase over the hand coded solution of 419 steps or a 45% increase.

A time to solution comparison was made next since the primary purpose of the BAX59 system is to allow an individual to quickly obtain a program capable of running on the TI-59 calculator. To begin with, the hand coded version was developed by three individuals over a period of several weeks. The BASIC program used as input for the BAX59 system

took only one person one day to write and debug. The utility of a higher level language greatly simplified the programming process. It is this savings in programming development which makes the desirability of the BAX59 system readily apparent.

In looking at the system-generated program some more comments can be made about where the relative overhead occurs. Of the total 652 programming steps it was noted that 84 steps were due to prompting code. Six storage registers were used as manual return registers while one register was used as a temporary display storage register. Another register was used in the manual subroutine return prompting scheme. This totals for the Linker an overhead of 148 programming steps. This Linker overhead represents 11% of the total generated steps indicating that the compiler generated at least 34% more code than the hand optimized coded program.

One last comment concerning the actual running of the BAX59 generated program needs to be made. In running the BAX59 code it was determined that time of execution became totally dependent on the amount of cards required to be read in and out. If the head reader in the calculator functioned properly, then this required small amounts of time to accomplish. If however, the head reader malfunctioned, then the reading of cards became an ordeal. In addition, the user had to pay close attention to the program prompting scheme or he would become lost between card reads.

VI. CONCLUSION

In the following discussions, the test program is evaluated together with the actual design. Based on this evaluation several conclusions are drawn and recommendations presented.

A. EVALUATION OF TEST RESULTS

In examining the test data, it is important to realize that this is only a single test. As such, it does not represent the whole set of BASIC programs which may be executed by the BAX59 software system. However, the test does give an insight into the actual efficiencies which might be expected. While actual numerical data is given these data should not be viewed as a statistical analysis of the system. Rather, the data is meant to provide some frame of reference for the discussion of system efficiency.

In examining the test program, it is noted that excess code generated amounts to roughly 45%. Of this, approximately 11% can be attributed to the linker, while 34% can be attributed to the cross-compiler. Although the total overhead seems rather large, the reason for building the system must be recalled. The primary reason is to facilitate the rapid design and implementation of programs on the TI-59 calculator. In view of this, it becomes clear that the overhead is secondary to the problem. Our real yardstick for success is whether or not TI-59 programs can be developed more rapidly than hand coded programs. The test program provides an insight into this side of the problem. Development time was about one order of magnitude faster compared to the hand coded solution.

This rapid development time more than justifies the high overhead. This is true in most academic applications, as most program executions are limited in nature. If on the other hand, a program is to be executed many times, then an optimized hand coded program might be better than the machine generated version. The final decision lies with the user. His program execution requirements, and the amount of time he has available to design and build his solution, will drive his selection.

E. CONCLUSIONS

In view of the target machine limitations, it is probably safe to conclude that the system is a valid first-cut prototype. The prototype proved the desirability and feasibility of the concept, that quick calculator programming can be realized with the minimum of effort. The following paragraphs discuss the prototype's limitations and suggest reasons why the current system is not yet useful as a good production system.

The calculator is severely limited in memory capacity. The TI-59 calculator has only 959 program steps for program usage. The overhead in code generation and segmentation prompts use up 45% of these steps. Together with the fact that only three memory partitions between program steps and storage registers are available, the 45% becomes a significant driving figure in calculator use.

The memory problem restricts the varieties of programs which may be written for translation. Programs may not be written which require a main routine having a long back jump that covers a vast portion of memory. This is because the linker segmentation rules will be violated underneath the covered iterative segment. The segmenter will fail to segment. Thus, only programs which are sequential in nature are suitable for the system.

Another related problem is that the smaller the memory, the more segmentation breaks will occur in the code. The more segmentation breaks, the more card reads will be required to achieve a successful program execution. Often, a problem will arise with the card reader of the calculator. Like any piece of equipment with a motor and magnetic tape head, it is fairly sensitive and prone to failure. If the card reader fails just once in the execution sequence, then there is a high probability that the program will fail to terminate successfully. The minimization of magnetic card reads is desirable for this reason as well as for reduction of user thrashing.

Another restriction occurs in the language subset. Arrays were not implemented in the first prototype design. This limitation impacts directly on the types of programs which can be developed for and translated by BAX59.

Arrays are used primarily for iterative work. Without arrays, iterative work, while still possible, is very limited. Much computer power lies in the ability to execute iteration rapidly. As noted in preceding paragraphs, this limitation occurs as a result of the small memory capacity of the machine. Because of this, it is felt that the implementation of arrays should occur when the prototype is matched with a more capable calculator.

We have suggested major limitations of the system as it currently stands. For these reasons it is felt that the system should be viewed only as a first working prototype. However, we feel that this prototype successfully demonstrates the concept that the power and efficiency of calculator programming can be greatly extended through higher level language programming.

C. RECOMMENDATIONS

If the concept of BAX59 is useful, then the next logical step is to develop the second prototype. The second prototype should not be hindered by the current TI-59's limitations. Otherwise, the next machine should be a more practical one, allowing easy hand held calculator programming. Many firms now market machines which have built-in BASIC language interpreters.

1. Hardware Related Suggestions

In order to avoid the major restriction, namely memory, the BAX59 system must target a larger capacity calculator. This calculator should have about 10,000 program steps and approximately 400 to 500 storage registers. A memory partitioning capability should be available to maximize memory usage.

As a follow on to increased memory capacity, the next prototype should have a hardware device available which will enable the host computer to download the generated calculator program into the target machine. This would eliminate hand punching program steps, which would be prohibitive on a calculator program of 10,000 or more steps.

The linker algorithm examines the dynamic structure of the program to facilitate segmenting the program. This algorithm segments sequential code that is not covered by a back loop. It may be possible to use this algorithm in the development of a single page swapping mechanism/system for a small microcomputer. The purpose of the algorithm in such a system would be to segment a program too large for the microcomputer, in such a manner so as to minimize the number of single page swaps with the system's disk storage unit. Such a system might be desirable for a small microcomputer in which memory is a problem.

2. Array Implementation

A very useful extension of the language subset would be the inclusion of arrays. Coupled with memory expansion, the capability to process arrays would make it possible to do more iterative programming.

A simple but costly implementation of single dimensional arrays is, perhaps, the most feasible approach. For each array declared, three registers will be required to store indexing and access information. Call the WEASIC record fields which will store the assigned register numbers BASE, HOLD, and CALL. BASE stores the number of the array base register (array index 0). During assignment statement translation, HOLD stores the index for an array identifier on the left hand side of the assignment statement equal sign. CALL stores the index of any array identifier currently being evaluated in simple expression. Of course, registers will be necessary for storage of the array itself. The simplest technique is to require that all array index ranges start at zero. Additionally, there will be no runtime checking of range limits. With the foregoing restrictions, estimated TI-59 program step requirements for translation of even simple assignments involving single dimensional arrays are very high. Evaluation of one array reference such as "A(X)" translates to 13 program steps. The assignment statement "A(X)=A(Y)" translates to 26 steps. A more complex reference such as "A(A(X))" requires about 24 steps. Together with the number of registers needed to store values and access data, usage of calculator memory may rapidly approach capacity levels with array manipulation.

Particular array values are accessed with the IND (indirect) instruction. Our basic strategy is to add the evaluated array subscript to the register number in BASE, and store it in HOLD (for left side of assignment

statements) or CALL (for all others). When a value is to be assigned to an array element on the left of an assignment statement, then the right side is evaluated and "STO IND" HOLD stores the value at the correct location. If it is only necessary to evaluate an array item within an expression, then "RCL IND" CALL recalls the value of the appropriate item. More efficient translation schemes might be possible; however, our technique has been tested on the TI-59 calculator, and works well.

The most difficult aspect of implementation is the the task of installing the translation scheme and a parsing scheme. Fortunately, the BASIC language requires explicit array declarations using the DIM statement. Procedure PDIM must be written to parse these declarations, create symbol table entries, and make register assignments. The SLOTRCD record would need an additional variant tag type for array types, call it ARRID. A slot with this tag would carry fields EASE, HOLD, and CALL in its variant part. Finally, we would have to adjust the simple expression parsing and code generating procedures PEXPR and PPRIMARY so that they could recognize array references and act accordingly. Of course, there are other source code adjustments that would be necessary to fine tune the system. However, this discussion has suggested our outline of major steps involved in array implementation.

APPENDIX A

WBASIC SUBSET RECOGNIZED BY BAX59

Command Reserved Words:

DATA	ENDIF	GOTC	NEXT	REM
DEF	ENDLOOP	IF	PAUSE	RESTORE
ELSE	FNEND	INPUT	PRINT	RETURN
ELSEIF	FOR	LET	QUIT	STOP
END	GOSUB	LOOP	READ	UNTIL
				WHILE

OFTICN (special--does not follow WBASIC syntax)

Supplemental Reserved Words:

NCT	<>	<	+	(
STEP	<=	>	-)
THEN	>=	=	*	!
IC	**	!	/	:

ε (special--recognized by scanner directly)
- (special--recognized by scanner directly)

Unimplemented Reserved Words:

CHAIN	LINECT	RANDOMIZE	SLEEP
CLOSE	LOCK	REMOVE	SORT
DIM	MAT	RENAME	TAGSORT
ENDGUESS	ON	RESUME	UNLOCK
GUESS	OPEN	SCRATCH	USE

OFTICN (special--syntax of WBASIC not implemented)

AND	OR	.
#	\$	

Description of Command Reserved Words:

DATA: Create an internal data list (see READ, RESTORE)

DATA <integer|real> [, <integer|real>]

DEF: Define a single or multi-line function (see FNEND).

single line: DEF <fn_name> [{ <parameter-list> }] = <expr>
multi-line: DEF <fn_name> [{ <parameter-list> }]
 ...body of definition
 FNEND

ELSE: Indicate instructions to execute if no IF/ELSEIF conditions were satisfied (see IF, ELSEIF, ENDIF).

ELSE

ELSEIF: Cause execution of a number of statements depending on the given condition (see IF, ELSE, ENDIF).

ELSEIF <boolean-expr>

END: Mark the end-of-source in the program (last line).

END

ENDIF: Indicate the end of an IF-ELSEIF-ELSE structure (see IF, ELSEIF, ELSE).

ENDIF

ENDLOOP: Mark the end of a loop (see UNTIL, WHILE, LOOP).

ENDLOOP

FNEND: Mark the end of a function definition (see DEF).

FNEND

FOR: Mark the start of a loop (see NEXT).

FOR <for-var> = <expr> TO <expr> [STEP <expr>]
 ...statements to execute in loop
NEXT <for-var>

GOSUB: Transfer control to the line specified, until a RETURN is reached (see RETURN).

GOSUB <line#> (Note: GO SUB is not recognized)

GOTO: Transfer control to the line specified (see ON).

GOTO <line#> (Note: GO TO is not recognized)

IF: (1) Cause transfer of control to either of two statements or QUIT a loop depending on a condition.

```
IF <boolean-expr> THEN <line#>|QUIT [ELSE <line#>|QUIT]
NOTE: This is an exception to WBASIC which
      allows any single statement after THEN
      and/or ELSE.
```

(2) Cause execution of either of a group of statements depending on a condition (see ELSE, ELSEIF, ENDIF).

```
IF <boolean-expr>
...statements to execute if expression TRUE
[ELSEIF <boolean-expr>
...statements to execute if 2nd expression TRUE]
[ELSE
...statements to execute if none are TRUE]
ENDIF
```

INPUT: Transmit data from the terminal to a number of variables (see PRINT). No variables stops execution.

```
INPUT [<expr> {, <expr>}]
```

LET: Assign the value of an expression to a variable.

```
[LET] <var> = <expr>
```

LOOP: Mark the beginning of a loop (see WHILE, ENDOOP, UNTIL).

```
LOOP
...statements to execute in loop
ENDLOOP
```

NEXT: Mark the end of a FOR loop (see FOR).

```
FOR <for-var> = ...
...statements to execute in loop
NEXT <for-var>
```

PAUSE: Suspend execution of the program.

```
PAUSE
```

PRINT: Transfer a series of values to printer or display. If no expression is found, a line space will result. (see OPTION)

```
PRINT [<expr> {, <expr>}]
```

QUIT: Leave the current block (WHILE, UNTIL, LOOP).

```
QUIT
```

READ: Transfer data from the list of items specified in DATA statements (see DATA, RESTORE).

```
READ <variable> {, <variable>}
```

REM: Indicate that the line is a comment. The exclamation character (!) may also be used to indicate a comment.

REM [<comment>]

RESTORE: Cause the next READ statement to get data values starting at the first item in the DATA list (see READ, DATA).

RESTORE

RETURN: Transfer control to the statement following the last GOSUB executed (see GOSUB).

RETURN

STOP: Terminate program execution.

STOP

UNTIL: Mark the end of a loop to be executed until the given condition is true (see WHILE, LOOP, ENDLOOP).

LOOP

...statements to execute in loop
UNTIL <boolean-expr>

WHILE: Mark the beginning of a loop to be executed until the given condition is no longer true (see LOOP, ENDLOOP, UNTIL).

WHILE <boolean-expr>

...statements to execute in loop
ENDLOOP

OPTICN: Set/reset boolean toggles within B&X59 to control generation of output files.

CAUTION: This is not the #BASIC OPTION!
This OPTICN should be used only after a correct #BASIC program has been constructed and is ready for translation.

OPTICN <opt-param> [<opt-param>]

where <opt-param> (option parameter) is
an integer range -8..+8;
sign indicates the direction of toggle:
positive = true/on, negative = false/off;
sign is assumed positive if omitted.

	Default
0 = generate linker interface file.....	false
1 = generate code for PC-100 printer.....	true
2 = optimize out unnecessary parentheses.....	true
3 = optimize out unnecessary NOP's.....	true
4 = translated TI-59 code to list file.....	true
5 = contents of symbol table to list file.....	false
6 = contents of code structure to list file.....	false
7 = each lexical token to terminal.....	false
8 = each lexical token to list file.....	false

Description of Supplemental Reserved Words:

NOT: Negate a boolean expression (see IF, WHILE, UNTIL).

NOT <boolean-expr>

STEP: Designate the increment (decrement) value of a FOR variable (see FOR). (default = +1)

FOR <for-var> = <expr> TO <expr> STEP <expr>

THEN: Mark the beginning of the true branch of a line-oriented IF statement (see IF).

IF <boolean-expr> THEN <line#>|QUIT [ELSE <line#>|QUIT]

TO: Mark the expression which represents the limiting value of a for-variable (see FOR)

FOR <for-var> = <expr> TO <expr> [STEP <expr>]

Symbols and Operators:

<>	not equal	+	addition
<=	less than or equal	-	subtraction
>=	greater than or equal	*	multiplication
**	raise to the power	/	division
<	less than	(open expression
>	greater than)	close expression
=	equal	:	list item separator
!	end of line cmt	.	decimal point

Special Characters:

(not reserved words--recognized directly within the scanner without reference to the RW table)

& Signifies that the current line is continued on the next line or is a continuation of the last line.

CO120 REM This comment is too long for one line, so it &
CO130 & must be continued on the next line.

- Underscore; used within variable identifier names to assist in readability; also used to designate a user defined function identifier.

LET FINAL_SUM_VALUE = FIRST_VALUE + SECOND_VALUE
DEF FN_FACTORIAL...

Built-in Functions:

AES	ATN	CSC	IP	PI	SIN
ACCS	COS	EXP	LOG	RND	SQR
ASIN	COT	FP	LOG10	SEC	TAN

ABS: Returns the absolute value (magnitude) of parameter
AES (x)

ACOS: Returns the arccosine (in radians) of parameter
ACCS (x)

ASIN: Returns the arcsine (in radians) of parameter
ASIN (x)

ATN: Returns the arctangent (in radians) of parameter
ATN (x)

COT: Returns the cotangent of parameter angle in radians
CCT (x)

CSC: Returns the cosecant of parameter angle in radians
CSC (x)

EXP: Returns the value of e raised to the power x
EXP (x)

FP: Returns the fractional part and sign of parameter
FP (x)

IP: Returns the integer part and sign of real parameter
IP (x)

LOG: Returns the natural logarithm (base e) of parameter
LOG (x)

LOG10: Returns the logarithm (base 10) of parameter
LOG10 (x)

PI: Returns the value of the constant pi
(pi = 3.141593 WBASIC ==> pi = 3.14159265359 TI-59)
PI

RND: Returns a pseudo-random number in the range (0,1)
RND (x)

SEC: Returns the secant of parameter angle in radians
SEC (x)

SIN: Returns the sine of parameter angle in radians
SIN (x)

SQR: Returns the square root of parameter
SQR (x)

TAN: Returns the tangent of parameter angle in radians
TAN (x)

APPENDIX B
CONDENSED BAX59 USER'S GUIDE

This guide is intended to be a useful compendium of important points the user should consider when preparing, cross-compiling, and linking a WBASIC source program with the BAX59 system. Included are suggested programming techniques which will optimize and improve resulting TI-59 code. Some of the information contained in the design document is repeated here for the sake of consistency. There are a few previously unmentioned items, many of which are essential to successful use of the system.

1. Whether you are translating a prewritten WBASIC program or one which you are writing yourself, review it for constructs and functions which are not implemented in BAX59. Use Appendix A as a quick reference for this purpose. Finding and eliminating unimplemented functions is more important than constructs. BAX59 will detect and report construct subset errors, however, unimplemented functions are assumed to be variable identifiers and will be entered in the symbol table as such. An error may or may not be reported as a result of faulty syntax; this depends upon the context of the unimplemented function.
2. Every line of the source program must have a line number, including blank lines. Every line number must be in chronological order between 00000 and 99999 exclusive.
3. The end of the source file does not require the END statement; however, whenever an END statement is encountered, the end of the source program is assumed

and translation halts. The END statement will not generate a TI-59 program stop. If you desire that a TI-59 program halt gracefully, you must use the STOP statement(s) in the WBASIC source code.

4. There is no overhead involved in using blank lines or comments. Although a NOP is generated for each, it is subsequently optimized out during resolution. This also allows unconditional jumps to such constructs without cause for concern. However, such practice is not recommended and will hamper debugging.
5. You are strongly urged to practice structured programming. While an unstructured program will be translated, its physical correspondence to the original WBASIC source code is likely to be less recognizable. Also, such a translation will very probably confuse the linker. Ensure that your subroutines have only one entry point and one RETURN, otherwise you will definitely confuse the linker! In order to improve the physical correspondence between source and translation, you are encouraged to use blank lines and comments. This will usually assist in debugging, if required.
6. The structured order of WBASIC program parts should be as follows:
 - A. OPTION statement (for BAX59 only)
 - B. DATA/READ statements
 - C. Main Body (including a at least one STOP)
 - D. Function DEF's and Subroutines
 - E. END statement

Note: The linker expects to find the label representing the main body (LBL A) as the first two TI-59 ksyccodes.

7. There is only one type of numeric data: real. Integers and reals are both considered reals in the TI-59 run-time environment. Numeric entries without

exponent will always be truncated to ten significant digits. If the entry contains an exponent or if the entry must be converted to exponent notation in order to maintain equivalence, then only eight significant digits will be saved (plus two in the exponent). These rules can have a profound effect upon precision errors in numerical computations.

8. Avoid proliferation of variable names. Variable names use registers, your most precious resource! Whenever possible reuse variable names to prevent new register assignments.
9. Do not forget to reserve registers for your own requirements. The "long" function facility always requires one in the range 10-99. The linker always requires two in the range 00-09. An additional two per segment in the range 10-99 will be taken dynamically by the linker after the cross-compiler has made all its assignments. The interface information is passed through the SCRATCH file. Never reserve the last available register, otherwise a memory overflow will never be reported in a warning message.
10. Optimize expressions using the standard rules of operator precedence. Failure to do so may result in unnecessary generation of parentheses.
11. Avoid use of the STEP option in FOR loops. Rely on the default increment (+1) whenever possible.
12. If a user-defined function is to be applied for its side effects only, then use one variable name to invoke all such function calls. For example:

```
00120 INVOKE = FN_ALPHA
00130 INVOKE = FN_BETA (X,Y)
00140 INVOKE = FN_GAMMA
```
13. Although contrary to principles of good structured programming, do not pass any parameters unless

absolutely necessary. Parameter passing uses a great deal of program steps. Furthermore, since actual parameters may be simple expressions, nesting of parentheses can become arbitrarily deep very quickly in a function call. The TI-59 places a limit of nine on the depth of parentheses nests.

14. Remember that the TI-59 allows subroutine (SBR) nesting to a maximum of six levels. This restricts the depth of recursion. However, if the recursive call always returns to the same address, the recursion will probably work. This is because the subroutine return stack will always maintain the correct return address, even if it overflows.
15. BAX59 distinguishes between upper and lower case characters in variable identifiers. Exceptions are the "E" in an exponent, the "FN_" preface of a function, and reserved words. Built-in functions must be written as they appear in the BIPNQF or BIFNLF files (currently, all upper case).
16. WBASIC trigonometric functions compute in radians. By default, TI-59 trigonometric functions compute in degrees. If trigonometric functions are translated, then prior to execution the TI-59 must be reset to the radian mode by entering "2ND RAD."
17. If you plan to modify TI-59 code that has been generated by BAX59, remember that only subroutines have relative addresses. All other addresses are absolute justified. Code insertions or deletions do not rejustify absolute addresses! Unless you are familiar enough with the calculator to know what you are doing, you will create more problems than you fix.
18. The special construct, PAUSE, is provided for assistance in debugging. This is somewhat like a message to the compiler. It translates to the TI-59

keyccdes 82 and 31, which are a void key and the "LRN" key. These keycodes cannot be entered directly into the calculator. Instead, enter "STO 31" followed by the editing sequence "BST BST NOP SST." The original "STC 31" will have been changed to "NOP 31." When encountered during run-time, these keys will interrupt execution and cause the calculator to enter its Learn Mode. Other than stopping execution, no other harmful effects result. To resume processing, simply hit "LRN" to show the contents of the current display register, and "R/S" to continue execution. This facility provides a convenient method of process suspension which corresponds physically and logically to the WBASIC source code.

19. The BAX59 CPTICN statement may provide other useful facilities for debugging. However, most of these were designed for debugging during installation of enhancements to the BAX59 Cross-Compiler Beware of CPTICN parameters 6, 7, and 8. These tend to produce a great deal of output!
20. The BAX59 system will not execute properly unless all associated data files are available to the host operating system on which BAX59 will run. You may modify the information contained in them, however, you should not change the formats. (These files are Appendices D, E, F, G, H, and J.)
21. Do not design excessively long iterative loops or back jumps. The linker cannot handle them. If iteration is required, design loops which translate to back jumps that are well within the TI-59 memory constraints.
22. The key to successful use of the linker is to break very large programs into smaller parts which can be processed sequentially without much repetition.

APPENDIX C

CROSS-COMPILER SOURCE CODE

```

*****
*
*                               BAX59 CROSS-COMPILER
*
* IMPLEMENTS A RECURSIVE DESCENT PARSER AND GENERATES
* A LINKED RECORDED DATA STRUCTURE OF TI-59 CODE
* TRANSLATED FROM WEASIC LANGUAGE SOURCE CODE.
* THE DATA STRUCTURE IS USED TO RESOLVE RELATIVE TI-59
* ADDRESSES. A LISTING OF THE ORIGINAL WBASIC PROGRAM
* INCLUDING DETECTED ERRORS IS GENERATED DURING
* TRANSLATION. VARIOUS FORMS OF OUTPUT BESIDES THE
* LISTING FILE CAN BE TOGGLED ON/OFF FROM WITHIN THE
* INPUT FILE USING THE "OPTION" STATEMENT. THIS VERSION
* OF THE PROGRAM IS DESIGNED TO SUPPLEMENT A WATERLCO
* BASIC (WEASIC) LANGUAGE INTERPRETER OR COMPILER. AS
* SUCH, IT DOES NOT DETECT ALL WBASIC SYNTAX OR SEMANTIC
* ERRORS. WBASIC PROGRAMS SHOULD BE SUCCESSFULLY RUN
* IN THE WEASIC SYSTEM ENVIRONMENT PRIOR TO TRANSLATION
* WITH THIS CROSS-COMPILER. THE BAX59 SYSTEM INCLUDES
* AN INDEPENDENT LINKER (TSDRIVER) WHICH WILL PROPERLY
* SEGMENT AND ISSUE INSTRUCTIONS FOR MANUALLY LINKING
* AND EXECUTING A TI-59 PROGRAM GENERATED BY THE CROSS-
* COMPILER BUT WHICH IS TOO LARGE FOR THE CALCULATOR
* MEMORY CAPACITY.
*
*****

```

```

PROGRAM BAX59      (INPUT,   OUTPUT, BASICF, MSGF,
                  RWIENLF, LABELF, CTEXT?, BIFNOF, BIFNLF,
                  OUTFILE, LISTF, NAMEF, READF, SCRATCH);

```

```

*****
*
*                               SYSTEM PARAMETERS
*
*****

```

```

-----
*
*                               CONST DECLARATIONS (MAIN)
*
-----

```

```

CONST  RWCHARCT   = 270; (* TOTAL # OF CHARS IN RW ARRAY *)
      RWWORDCT    = 72; (* TOTAL # OF WORDS IN RW ARRAY *)
      RWIENGCT    = 9;  (* # OF CHARS IN LONGEST RW *)
      MAXTOKLEN   = 20; (* MAX ACCUM LENGTH => MAX TOKEN *)
      MAXLINLEN   = 66; (* MAX LENGTH OF BASIC TEXT LINE *)
      MAXEASLIN   = 99999; (* MAX BASIC PROGRAM LINE NUMBER *)
      HASHBASE    = 99; (* INDEX OF LAST BUCKET (0-99) *)
      STARTREG    = 00; (* 1ST REGISTER (LOWEST NUMBER) *)
      REGEASE     = 90; (* MAX # AVAILABLE REGISTERS *)
      LBLEASE     = 72; (* MAX # OF AVAILABLE LABELS *)
      FNCLN      = 4;  (* MAX # STEPS IN QUICK FUNCTION *)
      FNLIEN     = 15; (* MAX # STEPS IN LONG FUNCTION *)
      FNIREG     = 10; (* REG USED FOR PARM OF LONG FNS *)
      FNSTACKLIM = 6;  (* MAX SBR/FN NESTING LEVEL *)
      TEXTLEN    = 20; (* MAX # CHARS IN A CODE TEXT LN *)

```

 *
 * GICBAL DECLARATIONS *
 *

 * TI-59 KEY CODES: OTHER SYMBOLS: *

K_ZERO	=	0:	K_EE	=	52:	BLANK	=	' '
K_1	=	1:	K_OPAREN	=	53:	ENDLIN	=	'@'
K_2	=	2:	K_CPAREN	=	54:	ENDFIL	=	'#'
K_3	=	3:	K_DIVOP	=	55:	PERIOD	=	'.'
K_4	=	4:	K_ENG	=	57:	COMMA	=	','
K_5	=	5:	K_FIX	=	58:	USCORE	=	'_'
K_6	=	6:	K_INT	=	59:	EXCLAM	=	'!'
K_7	=	7:	K_DEG	=	60:	QUOTE	=	'\"'
K_8	=	8:	K_GTO	=	61:			
K_9	=	9:	K_PG MIND	=	62:			
K_10	=	10:	K_EXCIND	=	63:			
K_A	=	11:	K_PRDIND	=	64:			
K_EE	=	12:	K_MULTOP	=	65:			
K_EC	=	13:	K_PAUSE	=	66:			
K_DE	=	14:	K_IPXEQT	=	67:			
K_E	=	15:	K_NOF	=	68:			
K_M	=	16:	K_OP	=	69:			
K_BPR	=	17:	K_RAD	=	70:			
K_CPR	=	18:	K_SBR	=	71:			
K_DPR	=	19:	K_STCIND	=	72:			
K_2CLR	=	20:	K_RCLIND	=	73:			
K_INV	=	22:	K_SUMIND	=	74:			
K_LNX	=	23:	K_SUBOP	=	75:			
K_CE	=	24:	K_LBL	=	76:			
K_CIR	=	25:	K_IPXGET	=	77:			
K_2INV	=	27:	K_SIGMA	=	78:			
K_ICG	=	28:	K_XBAR	=	79:			
K_CP	=	29:	K_GRAD	=	80:			
K_TAN	=	30:	K_RST	=	81:			
K_XT	=	32:	K_GTOIND	=	83:			
K_XSQR	=	33:	K_OPIND	=	84:			
K_XSRT	=	34:	K_ADDOP	=	85:			
K_XINV	=	35:	K_STFLG	=	86:			
K_FGM	=	36:	K_IPFLG	=	87:			
K_FR	=	37:	K_DMS	=	88:			
K_SIN	=	38:	K_PI	=	89:			
K_CC	=	39:	K_LIST	=	90:			
K_IND	=	40:	K_RS	=	91:			
K_STO	=	42:	K_INV SBR	=	92:			
K_RCL	=	43:	K_DE CPT	=	93:			
K_SUM	=	44:	K_NEG	=	94:			
K_ECWR	=	45:	K_EQUALS	=	95:			
K_CMS	=	47:	K_WRITE	=	96:			
K_EXC	=	48:	K_DSZ	=	97:			
K_FED	=	49:	K_ADV	=	98:			
K_ABSX	=	50:	K_PRT	=	99:			

{*-----*}
* TOKEN NUMBERS USED MOST OFTEN OUTSIDE OF MAIN DRIVER. *

ERROR TOK	=	0:	IDENTOK	=	RWORDCT	+ 1:
CMTCKEXC	=	1:	NUMBERTOK	=	RWORDCT	+ 2:
EQUAL TOK	=	2:	ENDLINTOK	=	RWORDCT	+ 3:
PLUSTCK	=	3:	ENDFIL TOK	=	RWORDCT	+ 4:
MINUSTOK	=	4:				
MULTCK	=	5:	TOTOK	=	20:	
DIVTCK	=	6:	OBTOK	=	21:	
OPARENTOK	=	7:	CMTOKREM	=	24:	
CPARENTOK	=	8:	NOT TOK	=	28:	
GITCK	=	9:	AND TOK	=	30:	
LITCK	=	10:	THE NTOK	=	33:	
COMMATOK	=	11:	ELSETOK	=	34:	
EXPTOK	=	15:	QUIT TOK	=	38:	
NOTEQTOK	=	16:	STEPTOK	=	42:	
GTECTCK	=	17:	ENDLOOPTOK	=	67:	
LIECTCK	=	18:				

```

{ *-----*
*
*           TYPE DECLARATIONS (MAIN)
*
*-----*

```

```

TYPE  BASLNRNG  = 0..MAXBASLIN;           (* SUBRANGES *)
      TCKENRNG  = 0..FWWORDCT + 4;
      HASHRNG   = 0..HASHBASE;
      ACCRNG    = 0..MAXTOKLEN;
      LBLRNG    = 0..LBLBASE + 1;
      REGRNG    = 0..99;
      KEYRNG    = 0..99;
      CTEXTRNG  = -2..99;

      LNSTRING  = PACKED ARRAY (.0..MAXLINLEN + 1.) OF CHAR;
      TKSTRING  = PACKED ARRAY (.1..MAXTOKLEN.) OF CHAR;

      LEVEL1    = ARRAY (.1..RWCHARCT + 1.) OF CHAR;
      LEVEL2    = ARRAY (.1..RWWORDCT + 1.) OF INTEGER;
      LEVEL3    = ARRAY (.1..RWLENGCT + 1.) OF INTEGER;

      LBLSTACK  = ARRAY (.LBLRNG.) OF INTEGER;

```

```

{ *-----*
* DATA STRUCTURE USED FOR WBASIC READ/DATA STATEMENTS;
* ONE DATA ENTRY CONSISTS OF AN OPTIONAL SIGN (DEFAULT
* IS POSITIVE) AND AN INTEGER OR REAL NUMBER.
*-----*

```

```

      DATAITEM = RECORD
          NUMB : TKSTRING;
          SIGN : CHAR
      END;
      DATASTORE = ARRAY (.1..REGBASE.) OF DATAITEM; (* DATAITEM *)

```

```

{ *-----*
* DATA STRUCTURE WHICH HOLDS THE TEXT TRANSLATION OF
* ALL TI-59 KEY CODES READ FROM THE CTEXTF FILE;
* UNIT FIELD INDICATES INSTRUCTION TYPE:
* 0 = SINGLE STEP INSTR (INDEPENDENT)
* 1 = 2-STEP INSTR (FOLLOWED BY REG NUMBER)
* 2 = 3-STEP INSTR (FOLLOWED BY ABSOLUTE ADDR)
* 3 = 4-STEP INSTR (FOLLOWED BY REG OR FLAG NUMBER
* AND AN ABSOLUTE ADDR)
*-----*

```

```

      CODETEXT = RECORD
          UNIT : 0..3;
          CCLECHAR : PACKED ARRAY
                      (.1..TEXTLEN.) OF CHAR
      END;
      CTEXTSTORE = ARRAY (.CTEXTRNG.) OF CODETEXT; (* CODETEXT *)

```

```

*-----*
* CODERCD IS A SINGLE NODE IN THE TI-59 CODE DATA *
* STRUCTURE; LINERCD IS A SINGLE NODE IN THE CHAIN OF *
* WEASIC LINE NUMBERS TO WHICH THE TI-59 CODE STRUCTURE *
* IS ATTACHED; THIS PART OF THE CODE DATA STRUCTURE IS *
* USED TO LOCATE PORTIONS OF TI-59 CODE WHICH DIRECTLY *
* CORRESPOND TO WEASIC LINE NUMBERS. *
*-----*

```

```

CODEPTR = @CODERCD;
CODERCD = RECORD
    ADDR : INTEGER;
    KEY : INTEGER;
    JMPF : CODEPTR;
    SECP : CODEPTR;
    BAKF : CODEPTR;
END; (* CODERCD *)

```

```

LINEPTR = @LINERCD;
LINERCD = RECORD
    LINC : BASLINRNG;
    LPTR : LINEPTR;
    CPTR : CODEPTR;
END; (* LINERCD *)

```

```

*-----*
* SLCTRCDs ARE SYMBOL TABLE SLOTS; SLOTS ARE ATTACHED TO *
* HASH; HASH IS THE SYMBOL TABLE REPRESENTED AS A STATIC *
* ARRAY OF SLOT POINTERS; EACH SLOT POINTER IN HASH *
* REPRESENTS A SINGLE HASH BUCKET; EACH BUCKET MAY HAVE *
* ANY NUMBER OF SLOTS ATTACHED AT THE HEAD POINTER OF *
* EACH BUCKET (LIMITED ONLY BY MACHINE CAPACITY). *
* SLOTS MAY BE OF 4 DIFFERENT TYPES: VARIABLE IDS, *
* LONG FUNCTION IDS, QUICK FUNCTION IDS, AND PARAMETER *
* (OR PARAMETERLESS) FUNCTION IDS. *
*-----*

```

```

IDTYP = (VARID, FNQID, FNLID, FNPID);
SLCTPTR = @SLCTRCD;
HASH = ARRAY (.HASHRNG.) OF SLCTPTR;
OKEYSEQ = ARRAY (.1..FNQLEN.) OF KEYRNG;
LKEYSEQ = ARRAY (.1..FNLLEN.) OF KEYRNG;
SLCTRCD = RECORD
    IDENT : TKSTBING;
    SLCT : SLCTPTR;
    CASE TYP : IDTYP OF
        VARID : (REGNO : INTEGER;
                AUXREG1 : INTEGER;
                AUXREG2 : INTEGER);
        FNQID : (FNQ : OKEYSEQ);
        FNLID : (FNL : LKEYSEQ);
        FNPID : (FNLINK : SLCTPTR);
        FNP : (PNREGNO : INTEGER;
              FNP : SLCTPTR;
              LEL : LBLNG;
              FNPLINK : SLCTPTR);
END; (* SLCTRCD *)

```

```

(*)-----(*)
(*)
(*)          VAR DECLARATIONS (MAIN)
(*)
(*)-----(*)

```

```

(*)-----(*)
(*) CPTICN TOGGLES ARE BOOLEANS WHICH CAN BE SWITCHED FROM *
(*) WITHIN THE WBASIC SOURCE PROGRAM USING THE "OPTION" *
(*) STATEMENT; EXCEPT FOR ZERO (LINKER INTERFACE TOGGLE), *
(*) THE RULE IS THAT A '+' SETS/RESETS TOGGLE TRUE, WHILE *
(*) A '-' SETS/RESETS TOGGLE FALSE; DEFAULT VALUES ARE *
(*) INDICATED IN THE CCOMMENT FOLLOWING EACH DECLARATION. *
(*)-----(*)

```

```

VAR   LINK59 : BCOLEAN;          (* OPTION 0 = FALSE *)
      PC100  : BCOLEAN;          (* OPTION 1 = TRUE  *)
      OPTPAR : BCOLEAN;          (* OPTION 2 = TRUE  *)
      CPTNOF : BCOLEAN;          (* OPTION 3 = TRUE  *)
      CODUMF : BCOLEAN;          (* OPTION 4 = TRUE  *)
      SYDUMF : BCOLEAN;          (* OPTION 5 = FALSE *)
      DSDUMF : BCOLEAN;          (* OPTION 6 = FALSE *)
      TOKCUT : BCOLEAN;          (* OPTION 7 = FALSE *)
      TOKIIS : BCOLEAN;          (* OPTION 8 = FALSE *)

```

```

(*)-----(*)
(*) SETS USED IN VARIOUS TESTS FOR CHARS, TOKEN NUMBERS, *
(*) TI-59 KEY CODES, AND REGISTERS.
(*)-----(*)

```

```

LETTERS, DIGITS, ALPHANUM : SET OF CHAR;
DOUBLE1, DOUBLE2         : SET OF CHAR;
SPECIALS, SIGNS         : SET OF CHAR;
SUBERROR, CRITICAL      : SET OF CHAR;
BINCF TOKS, RELOFTOKS   : SET OF TOKENRNG;
TRAILTOKS, SIGNTOKS     : SET OF TOKENRNG;
BEGIN EXPRTOKS          : SET OF TOKENRNG;
NUMERICKEY              : SET OF KEYSRNG;
RESERVE_REG             : SET OF REGRNG;

```

```

(*)-----(*)
(*) RESERVED WORD TABLE CHARACTER AND INDEX ARRAYS.
(*)-----(*)

```

```

RWCHAR : LEVEL1;
FWWCRC : LEVEL2;
RWLZNG : LEVEL3;

```

```

(*)-----(*)
(*) SCANNER ASSOCIATED GLOBAL VARIABLES.
(*)-----(*)

```

```

ACCUM          : TKSTRING;
ACCINX         : ACCRNG;
LINEUF         : LNSTRING;
LNEINX        : 0..MAXLINLEN + 1;
LINUM, LLINUM, CLINUM : BASLINRNG;
TCKNUM, LTCKNUM : TOKENRNG;
LCCUNT, RCOUNT, ECOUNT : INTEGER;
FLAGCMT       : BOOLEAN;

```

```
{*-----*}
{* PARSE ASSOCIATED GLOBAL VARIABLES. *}
{*-----*}
```

```
ERRCRCT, WARNCT : INTEGER;
NEXTREG          : INTEGER;
LELCT           : LBLRNG;
RESERVECT       : REGRNG;
```

```
CTEXT           : CTEXTSTORE;
DATALIST        : DATASTORE;
IATAIX, READIX : 1..REGBASE;
INDEXERROR      : BOOLEAN;
FIRSTREAD       : BOOLEAN;
```

```
CLABEL          : LBLSTACK;
```

```
BUCKET          : HASH;
IDSLOT          : SLOTPTR;
LF, LFCUR       : LINEPTR;
LPLEAD, LPTRAIL : LINEPTR;
CF, CFCUR       : CODEPTR;
FIRSTLP, LASTLP : LINEPTR;
BEGINCP, ENDCP  : CODEPTR; (* MARKERS *)
```

```
FNSTACK, FNLLIST : SLOTPTR;
FNSTACKCT        : INTEGER;
```

```
IFSTACK, ENDIFSTACK : CODEPTR;
LOOPSTACK, ENDCOPSTACK : CODEPTR;
PCRSTACK, NEXTSTACK : CODEPTR;
```

```
{*-----*}
{* FILES *}
{*-----*}
```

```
RWTELF, LABELF, CTEXTF : TEXT; (* INITIAL FILES *)
BIFNCF, BIFNLF         : TEXT; (* BUILT-IN FNS *)
BASICF, MSGF           : TEXT; (* INPUT FILES *)
LISTF, NAMEF, READF    : TEXT; (* OUTPUT FILES *)
OUTFILE                 : TEXT; (* TERMINAL FILE *)
SCRATCH                 : TEXT; (* LINKER INTERFACE *)
```

```

{*****}
{*
  PRIMITIVE CHAR ROUTINES
*}
{*****}

{-01-----*}
{* UPCASE CCNVERTS ANY LOWER CASE EBCDIC (OR ASCII) CHAR *}
{* TC UPPER CASE EQUIVALENT. *}
{------*}

FUNCTION UPCASE (VAR CH : CHAR) : CHAR;
BEGIN
  IF CH IN ('a'..'i') THEN
    UPCASE := CHR(ORD(CH) - ORD('a') + ORD('A'))
  ELSE IF CH IN ('j'..'r') THEN
    UPCASE := CHR(ORD(CH) - ORD('j') + ORD('J'))
  ELSE IF CH IN ('s'..'z') THEN
    UPCASE := CHR(ORD(CH) - ORD('s') + ORD('S'))
  ELSE
    UPCASE := CH
END;
(* UPPERCASE *)

{*****}

{-02-----*}
{* TRANSDIGIT RETURNS THE INTEGER VALUE FOR NUMERIC CHARS *}
{------*}

FUNCTION TRANSDIGIT (CH : CHAR) : INTEGER;
BEGIN
  TRANSDIGIT := ORD (CH) - ORD ('0')
END;
(* TRANSDIGIT *)

{*****}

{-03-----*}
{* XNUMBER RETURNS INTEGER VALUE OF A NUMERIC CHAR STRING *}
{------*}

FUNCTION XNUMBER (ACCUM:TKSTRING; ACCINX:ACCRNG) : INTEGER;
VAR I, TEMPNR : INTEGER;
BEGIN
  TEMPNR := 0;
  FOR I := 1 TO ACCINX DO
    TEMPNR := TEMPNR * 10 + TRANSDIGIT (ACCUM(.I.));
  XNUMBER := TEMPNR
END;
(* XNUMBER *)

{*****}

```

```
(* -04 ----- *)
{* ZEROPAD WRITES INTEGERS TO AN OUTFILE WITH LEADING 0'S *}
{* ----- *
```

```
PROCEDURE ZEROPAD (VAR WFILE : TEXT; N, ZCT : INTEGER);
```

```
VAR I, TN : INTEGER;
```

```
BEGIN
```

```
  TN := N;
```

```
  REPEAT
```

```
    TN := TN DIV 10;
```

```
    ZCT := ZCT - 1
```

```
  UNTIL TN = 0;
```

```
  FOR I := 1 TO ZCT DO
```

```
    WRITE (WFILE, '0');
```

```
  IF N >= 0 THEN
```

```
    WRITE (WFILE, N:1)
```

```
  ELSE
```

```
    WRITE (WFILE, -N:1, '-')
```

```
END;
```

```
(* ZEROPAD *)
```

```

(*****
*
*                               SCANNER
*
*****

```

```

*-05-----*
* SCAN USES THE RESERVED WORD ARRAYS TO ISOLATE AND
* RETURN SINGLE TOKENS FROM THE WBASIC SOURCE FILE;
* IT ALSO REWRITES THE SOURCE CODE TO THE LISTF FILE
* ALONG WITH ANY SCANNER DETECTED ERRORS; SCAN INSERTS
* ITS OWN END-OF-LINE AND END-OF-FILE CONTROL CHARACTERS
* INTO ITS LINE BUFFER (LINBUF) AS THESE CHARS ARE
* DETECTED IN THE SOURCE FILE.
*-05-----*

```

```
PROCEDURE SCAN (VAR TCKNUM : TCKENRNG);
```

```
VAR TCHAR : CHAR;
    I : INTEGER;
```

```

(=====*)
*-05-01-----*
* LINENR RETURNS NUMBER OF THE CURRENTLY SCANNED WBASIC
* LINE; WRITES THE LINE NUMBER TO THE LISTF FILE
* PADDED WITH ZEROS.
*-05-01-----*

```

```
FUNCTION LINENR : BASLINRNG;
```

```
VAR I : BASLINRNG;
```

```

BEGIN
  READ (BASICF, I);
  ZERCPAD (LISTF, I, 5);
  LINENR := I
END;
(* LINENR *)

```

```

(*-----*)
*-05-02-----*
* RDLINE READS TEXT IMMEDIATELY FOLLOWING LINE NUMBER;
* RETURNS THE TEXT IN A MAXLINLEN CHAR BUFFER; UNUSED
* PORTION OF BUFFER IS FILLED WITH BLANKS; WRITES EACH
* CHAR OF TEXT TO THE LISTF FILE AS IT IS READ; REPORTS
* AN ERROR IF NUMBER OF CHARS EXCEEDS MAXLINLEN.
*-05-02-----*

```

```
FUNCTION RDLINE : LNSTRING;
```

```
VAR I, J, LINLENGTH : INTEGER;
    CH : CHAR;
```

```

BEGIN
  I := 0;
  WHILE (NOT (EOLN(BASICF))) AND (I < MAXLINLEN) DO
    BEGIN
      I := I + 1;
      READ (BASICF, CH);
      RDLINE(I) := CH;
      WRITE (LISTF, CH)
    END;
  WRITELN (LISTF);
  LINLENGTH := I;
  LNBINX := 0;
  IF LINLENGTH < MAXLINLEN THEN (* FILL UNUSED W/ BLANKS *)
    FOR J := LINLENGTH + 1 TO MAXLINLEN DO
      RDLINE(J) := BLANK;

```



```

IF (LINLENGTH = MAXLINLEN) AND (NOT EOLN(BASICF)) THEN
  BEGIN
    WRITE (LISTF, '**F** SCAN ERROR...LENGTH OF TEXT ');
    WRITEIN (LISTF, 'AFTER LINUM > ', MAXLINLEN:2, ' CHARS');
    ERRORCT := ERRORCT + 1;
    REPEAT (* LOCATE THE EOLN CHAR TO RECOVER *)
      GET (BASICF)
    UNTIL EOLN(BASICF)
  END;
  RDLINE(.LINLENGTH + 1.) := ENDLIN; (* INSERT EOLN CHAR *)
  GET (BASICF); (* MOVE FILE PTR PAST PASCAL EOLN CHAR *)
  IF EOF(BASICF) THEN (* OVERWRITE EOLN CHAR IF EOF *)
    RDLINE(.LINLENGTH + 1.) := ENDFIL
END; (* RDLINE *)

(*-----*)
(*-05-03-----*)
{* GETNOELANK RETURNS FIRST NON-BLANK CHAR STARTING WITH *}
{* THE CURRENT CHAR REFERENCED BY THE LINBUF INDEX. *}
(*-----*)

FUNCTION GETNOBLANK : CHAR;
BEGIN
  WHILE LINEUF(.LNBINX.) = BLANK DO
    LNBINX := LNBINX + 1;
  GETNOELANK := LINBUF(.LNBINX.)
END; (* GETNOELANK *)

(*-----*)
(*-05-04-----*)
{* GETCHAR RETURNS THE CHAR FOLLOWING LINE BUFFER INDEX *}
{* AND INCREMENTS THE LINE BUFFER INDEX. *}
(*-----*)

FUNCTION GETCHAR : CHAR;
BEGIN
  LNBINX := LNBINX + 1;
  GETCHAR := LINEUF(.LNBINX.)
END; (* GETCHAR *)

(*-----*)
(*-05-05-----*)
{* PUTCHAR INCREMENTS THE ACCUM INDEX AND PLACES A CHAR *}
{* INTO THAT POSITION IN THE ACCUM ARRAY. *}
(*-----*)

PROCEDURE PUTCHAR (CH : CHAR);
BEGIN
  IF NOT (ACCINX >= MAXTOKLEN) THEN
    BEGIN
      ACCINX := ACCINX + 1;
      ACCUM(.ACCINX.) := CH
    END
END; (* PUTCHAR *)

(*-----*)

```

```

(*-05-06-----*)
{* PUTANDGET PUTS A CURRENT CHAR INTO ACCUM AND GETS THE *}
{* NEXT CHAR FROM THE LINE BUFFER. *}
(*-----*)

```

```
FUNCTION PUTANDGET (CH : CHAR) : CHAR;
```

```

BEGIN
  PUTCHAR (CH);
  PUTANDGET := GETCHAR
END;
(* PUTANDGET *)

```

```
(*-----*)
```

```

(*-05-07-----*)
{* NUMSFCCL READS ALL DIGITS WHICH COMPRISE ONE UNSIGNED *}
{* INTEGER AND PUTS THEM INTO THE ACCUM IN SEQUENCE; *}
{* THE NUMBER OF DIGITS SPOOLED IS RETURNED. *}
(*-----*)

```

```
FUNCTION NUMSPOOL : INTEGER;
```

```

VAR I : INTEGER;
BEGIN
  IF TCHAR IN DIGITS THEN
    BEGIN
      I := 0;
      WHILE TCHAR IN DIGITS DO
        BEGIN
          TCHAR := PUTANDGET (TCHAR);
          I := I + 1
        END;
      NUMSFCCL := I
    END
  ELSE
    NUMSFCOL := 0
END;
(* NUMSPOOL *)

```

```
(*-----*)
```

```

(*-05-08-----*)
{* EXPONENT READS THE CHARS WHICH COMPRISE AN EXPONENT *}
{* PART, COUNTS THEM, AND PUTS THEM INTO THE ACCUM. *}
(*-----*)

```

```
PROCEDURE EXPONENT;
```

```

BEGIN
  IF (UPCASE (TCHAR) = 'E') THEN
    BEGIN
      TCHAR := PUTANDGET (TCHAR);
      IF TCHAR IN SIGNS THEN
        TCHAR := PUTANDGET (TCHAR)
      ELSE
        PUTCHAR ('+');
      ECCUNT := NUMSFCOL
    END
  END
END;
(* EXPONENT *)

```

```
(*-----*)
```

```
(*--05-09-----*)
{* DECIMALPT READS AND COUNTS THE DIGITS IN THE *
{* PRACTICAL PART OF A NUMBER AND PUTS THEM INTO ACCUM. *}
(*-----*)
```

PROCEDURE DECIMALPT;

```
BEGIN
  IF TCHAR = PERIOD THEN
    BEGIN
      TCHAR := PUTANIGET(TCHAR);
      RCCUNT := NUMSECOL
    END
  END;
(* DECIMALPT *)
```

```
(*-----*)
```

```
(*--05-10-----*)
{* ADJUST IS HIGHLY CALCULATOR-DEPENDENT; ADJUSTS ALL *
{* LITERAL NUMERICS TO TI-59 FORMAT; SETS VALUES FOR *
{* LCCUNT, RCOUNT, AND ECOUNT (LEFT, RIGHT, EXPONENT). *
{* TI-59 WILL ACCEPT FROM ITS KEYBOARD A MAX OF 10 DIGITS *
{* (PLUS DECIMAL POINT AND SIGN) FOR INTEGERS OR PEALS *
{* WITHOUT EXPONENT, OR 8 DIGITS (PLUS DECIMAL POINT AND *
{* SIGN) AND 2 DIGIT SIGNED EXPONENT. SINCE NUMBERS ARE *
{* SCANNED AND PUT INTO THE ACCUM AS THEY ARE READ IN *
{* WPASIC SOURCE CODE, THOSE WHICH EXCEED THE ABOVE *
{* MAXIMUMS MUST BE CONVERTED WHILE IN THE ACCUM. *
{* THIS ROUTINE IS DIFFICULT TO UNDERSTAND, MUCH LESS *
{* VISUALIZE, WITHOUT USING A SPECIFIC EXAMPLE TO WALK- *
{* THROUGH ON PAPER. *}
(*-----*)
```

PROCEDURE ADJUST;

```
(*--05-10-01-----*)
{* MCVEXP SHIFTS THE POSITION OF THE EXPONENT PART SO *
{* THAT THE NUMBER HAS MAX OF 8 SIGNIFICANT DIGITS, SIGN, *
{* AND DECIMAL POINT (NOTE: SIGNIFICANCE IS LOST). *}
(*-----*)
```

PROCEDURE MCVEXP;

```
VAR I : INTEGER;
BEGIN
  FOR I := 1 TO (2 + ECOUNT) DO
    ACCUM(.9 + I.) := ACCUM(.LCOUNT + RCOUNT + 1 + I.);
  ACCINX := 11 + ECOUNT
END;
(* MCVEXP *)
(*-----*)
```

```

(*-05-10-02-----*)
{* ADJEXF CCNVERTS NUMEERS TO EQUIVALENT TI-59 COMPATIBLE *}
{* FORM BY COORDINATING EXPONENT VALUE ADJUSTMENT WITH *}
{* DECIMAL FOINT MOVEMENT AND DIGIT TRUNCATION. *}
*-----*

```

```
PROCEDURE ADJEXP (DIFF : INTEGER);
```

```
VAR E1, E2, EXP, NEWEXP : INTEGER;
```

```

BEGIN
  E1 := CRD(ACCUM(.12.)) - CRD('0');
  E2 := CRD(ACCUM(.13.)) - CRD('0');
  EXP := 10 * E1 + E2;
  IF ACCUM(.11.) = '-' THEN
    NEWEXP := EXP - DIFF;
  ELSE
    NEWEXP := EXP + DIFF;
  E1 := TRUNC(NEWEXP/10);
  E2 := NEWEXP - (E1 * 10);
  ACCUM(.12.) := CHR(E1 + ORD('0'));
  ACCUM(.13.) := CHR(E2 + ORD('0'));

```

```
END; (* ADJEXP *)
```

```
(*-----*)
```

```

BEGIN (* ADJUST MAIN *)
  IF (LCCUNT > 10) OR ((LCOUNT > 8) AND (ECOUNT <> 0)) THEN

```

```

    BEGIN
      ACCUM(.9.) := PERIOD;
      IF ECOUNT = 0 THEN

```

```

        BEGIN
          ACCUM(.10.) := 'E';
          ACCUM(.11.) := '+';
          ACCUM(.12.) := '0';
          ACCUM(.13.) := '0';
          ACCINX := 13;

```

```
        END
```

```
      ELSE
```

```
        BEGIN
```

```
          MCVEXP;
```

```
          IF ECOUNT = 1 THEN
```

```
            BEGIN
```

```
              ACCUM(.13.) := ACCUM(.12.);
```

```
              ACCUM(.12.) := 'C';
```

```
              ACCINX := 13;

```

```
            END
```

```
          END;
```

```
          ADJEXP(LCOUNT - 8);
```

```
          LCCUNT := 8; RCOUNT := 0; ECOUNT := 2
```

```
        END
```

```
      ELSE IF (LCOUNT + RCOUNT > 10) OR ((LCOUNT + RCOUNT > 8) AND (ECOUNT <> 0)) THEN
```

```

        BEGIN
          IF ECOUNT = 0 THEN

```

```
          BEGIN
```

```
            ACCINX := 11;
```

```
            RCOUNT := 10 - LCCUNT
```

```
          END
```

```
        ELSE
```

```
          BEGIN
```

```
            MCVEXP;
```

```
            RCOUNT := 8 - LCCUNT
```

```
          END
```

```
        END
```

```
END;
```

```
(* ADJUST *)
```

```
(*-----*)
```

```

(*-05-11-----*)
{* RWLOOKUP LOOKS UP TOKEN IN RESERVE WORD TBL BASED UPON *}
{* TCKEN LENGTH; RETURNS TOKEN NUMBER; IF NOT FOUND, *}
{* TCKNUM = IDENTOK (IE. TOKEN IS ASSUMED IDENTIFIER). *}
(*-----*)

```

```
FUNCTION FWLOOKUP : INTEGER;
```

```
VAR MATCH : BOOLEAN;
    CHINDEX, WDINDEX, LGINDEX, ACINDEX : INTEGER;
```

```

BEGIN
  LGINDEX := ACINDEX;
  WDINDEX := RWLENG(.LGINDEX.);
  REPEAT
    MATCH := TRUE;
    ACINDEX := 1;
    CHINDEX := RWORD(.WDINDEX.);
    WHILE (MATCH) AND (ACINDEX <= LGINDEX) DO
      BEGIN
        IF UPCASE(ACCUM(.ACINDEX.)) <> RWCHAR(.CHINDEX.) THEN
          MATCH := FALSE
        ELSE
          BEGIN
            ACINDEX := ACINDEX + 1;
            CHINDEX := CHINDEX + 1
          END
        END;
      WDINDEX := WDINDEX + 1
    UNTIL (MATCH) OR (CHINDEX = RWORD(.RWLENG(.LGINDEX+1.)));
    IF MATCH THEN
      RWLOOKUP := WDINDEX - 1      (* BACK-UP THE WORD INDEX *)
    ELSE
      RWLOOKUP := IDENTOK
  END;
(* RWLOOKUP *)

```

```

(*-----*)
(*-05-12-----*)
{* CMTSFCCL READS AND DISREGARDS THE TEXT OF COMMENTS. *}
{* FLAGCMT IS USED FOR COMMENTS CONTINUED ON NEW LINE. *}
(*-----*)

```

```
PROCEDURE CMTSPOOL;
```

```

BEGIN
  FLAGCMT := FALSE;      (* RESET COMMENT CONTINUATION FLAG *)
  WHILE NOT (TCHAR IN CRITICAL) DO
    TCHAR := GETCHAR;
    IF TCHAR = 'E' THEN
      FLAGCMT := TRUE
  END;
(* CMTSPOOL *)

```

```

(*-----*)
(*-05-13-----*)
{* RECOVER SCANS AND DISREGARDS THE REMAINDER OF THE *}
{* CURRENT TOKEN AND STOPS AT START OF NEXT TOKEN. *}
(*-----*)

```

```
PROCEDURE RECOVER;
```

```

BEGIN
  WHILE NOT (TCHAR IN (CRITICAL + (.BLANK.))) DO
    TCHAR := GETCHAR      (* SKIP TO NEXT TOKEN AND RETURN *)
  END;
(* RECOVER *)

```

```
(*=====*)
```

```

BEGIN (* SCAN MAIN *)

LTCKNUM := TOKNUM; (* SAVE LAST TOKNUM IN LTOKNUM *)
TOKNUM := ERRORTOK; (* INITIALIZE NEW TOKEN NUMBER *)
ACCINX := 0; (* INDICATES TOKEN LENGTH IN ACCUM *)
LCCUNT := 0; (* NO. OF DIGITS LEFT OF DECIMAL *)
RCCUNT := 0; (* NO. OF DIGITS RIGHT OF DECIMAL *)
ECOUNT := 0; (* NO. OF DIGITS IN EXPONENT *)
TCHAR := GETNOBLANK; (* GET NEXT NON-BLANK CHAR *)

IF (TCHAR = ENDFIL) THEN
  TCKNUM := ENDFILICK

ELSE IF (TCHAR = '&') THEN
  BEGIN
    CLINUM := LINENR; (* READ LINE NO. OF CONT LINE *)
    LINBUF := RDLINE; (* READ TEXT OF CONT LINE *)
    TCHAR := GETCHAR; (* MOVE LNBINX PAST TRAIL "&" *)
    TCHAR := GETNOBLANK; (* FIND LEADING "&" ON NEW LINE *)
    TCHAR := GETCHAR; (* MOVE LNBINX PAST CONT "&" *)
    IF FLAGCMT THEN
      CMTSFOOL (* COMMENT CONTINUATION *)
    ELSE
      SCAN (TOKNUM) (* SCAN NEXT TOKEN AFTER "&" *)
    END
  END

ELSE IF (TCHAR = ENDLIN) THEN
  BEGIN
    LINUM := LINUM; (* PASS LINUM TO LAST LINUM *)
    LINENR := LINENR; (* READ LINE NO. OF NEW LINE *)
    LINBUF := RDLINE; (* READ TEXT OF NEW LINE *)
    TCHAR := GETCHAR; (* MOVE LNBINX PAST ENDLIN CHAR *)
    TCKNUM := ENDLINTOK (* ASSIGN TOKEN NO. FOR ENDLIN *)
  END

ELSE IF (TCHAR IN LETTERS) THEN
  BEGIN
    WHILE (TCHAR IN ALFANUM) DO
      BEGIN
        TCHAR := PUTANDGET(TCHAR);
        IF TCHAR = USCORE THEN (* ASSUMES USCORE WILL *)
          TCHAR := PUTANDGET(TCHAR) (* NOT OCCUR AT END *)
        END;
        IF ACCINX <= RWLENGCT THEN
          BEGIN
            TCKNUM := HWLOOKUP;
            IF TOKNUM = CMTOKREM THEN (* LOOK FOR REM CMT *)
              CMTSFOOL
            END
          ELSE
            TCKNUM := IDENTOK
          END
        END
      END

ELSE IF (TCHAR IN DIGITS) THEN
  BEGIN
    LCCUNT := NUMSFICOL;
    IF TCHAR = PERIOD THEN
      DECIMALPT
    ELSE IF (UPCASE(TCHAR) = 'E') THEN
      EXECNENT;
      ADJGST;
      TCKNUM := NUMBERTOK
    END
  END

```

```

ELSE IF (TCHAR = PERIOD) THEN
  BEGIN
    DECIMALPT;
    EXFCNENT;
    ADJUST;
    TCKNUM := NUMBERTOK
  END

ELSE IF (TCHAR IN DCUBLE1) THEN
  BEGIN
    TCHAR := PUTANDGET(TCHAR);
    IF (TCHAR IN DCUBLE2) THEN
      TCHAR := PUTANDGET(TCHAR);
    TOKNUM := RWLOCKUP
  END

ELSE IF (TCHAR IN SPECIALS) THEN
  BEGIN
    TCHAR := PUTANDGET(TCHAR);
    TCKNUM := RWLOCKUP;
    IF TOKNUM = CMTICXC THEN      (* LOOK FOR EXCLAM CMT *)
      CMTSPOOL
  END

ELSE IF (TCHAR IN SUBERROR) THEN
  BEGIN
    WRITE (LISTF, '**P** SCAN ERROR FOUND AT "', TCHAR);
    WRITEIN (LISTF, '"...CHAR NOT IN THIS SUBSET');
    ERRRCCT := ERRRCCT + 1;
    RECOVER
  END

ELSE
  BEGIN
    WRITE (LISTF, '**P** SCAN ERROR FOUND AT "', TCHAR);
    WRITEIN (LISTF, '"...UNRECOGNIZABLE CHAR');
    ERRRCCT := ERRRCCT + 1;
    RECOVER
  END;

FOR I := (ACCINX + 1) TO MAXTOKLEN DO
  ACCUM(.I.) := BLANK; (* BLANK OUT REMAINDER OF ACCUM *)

(*-----*)
(* DEBUGGING TOOL:  LISTS TOKNUM AND TOKEN AS IT IS READ. *)
IF TOKCUT THEN (* OPTION 7 *)
  WRITEIN (OUTFILE, ':6, TCKNUM:2, ' |', ACCUM, '|');
IF TCKLIS THEN (* OPTION 8 *)
  WRITEIN (LISTF, ':6, TOKNUM:2, ' |', ACCUM, '|')
(*-----*)

END; (* SCAN *)

```

```

{*****}
{*
{*          ERROR/LINE END HANDLING ROUTINES
{*
{*****}

{*--06-----}
{* PRECOVER SCANS A LINE AND DISREGARDS TOKENS UNTIL IT
{* FINDS A COMMENT, AN END OF LINE, OR AN END OF FILE.
{*-----}

PROCEDURE PPRECOVER;

BEGIN
  WHILE NOT (TOKNUM IN TRAILTOKS) DO
    SCAN (TCKNUM)
  END;
  (* PRECOVER *)

{*****}

{*--20--*}
PROCEDURE GENKEY (OPCCDE : INTEGER);          FORWARD;

{*****}

{*--07-----}
{* PERROR IS THE GENERAL PURPOSE ERROR HANDLER WHICH:
{* GENERATES SPACE FOR REGISTER OR ADDRESS INSERTION IN
{* ORDER TO PREVENT THE CODEDUMP ROUTINE FROM CAUSING A
{* SYSTEM ERROR DUE TO INVALID TI-59 CODE GENERATION;
{* ANNOTATES THE LISTING FILE WITH THE ERROR LOCATION;
{* INCREMENTS THE ERRCR COUNT; RECOVERS TO END OF LINE.
{*-----}

PROCEDURE PERROR;

BEGIN
  GENKEY (-1); (* GENERATE REG/ADDR SPACE TO PROTECT CODE *)
  GENKEY (-1); (* DUMP ROUTINE FROM OPERATING SYS ERROR. *)
  WRITELN (LISTF, '**F** FATAL ERROR FOUND AT ', ACCUM, '');
  ERRCRCT := ERRCRCT + 1;
  PRECOVER
END;
  (* PERRCR *)

{*****}

{*--08-----}
{* PSUBERROR DISTINGUISHES AN ERROR WHICH IS A RESULT OF
{* USING A WEASIC COMMAND NOT IN THIS IMPLEMENTATION.
{*-----}

PROCEDURE PSUBERROR;

BEGIN
  WRITELN (LISTF, '**F** SUBSET ERROR FOUND AT ', ACCUM, '');
  ERRCRCT := ERRCRCT + 1;
  PRECOVER
END;
  (* PSUBERRCR *)

{*****}

```



```

(*-09-----*)
{* PWARN IS USED TO LOCATE THE CAUSE OF A WARNING FOR THE *
{* USER AND TO INCREMENT THE WARNING MSG COUNT; NOTE THAT *
{* NORMAL COMPILE CONTINUES. *
(*-----*)

```

PROCEDURE PWARN;

```

BEGIN
  WRITELN (LISTF, '**W** WARNING TRIGGERED AT ', ACCUM, '');
  WARNCT := WARNCT + 1
END;
(* PWARN *)

```

(*****)

```

(*-10-----*)
{* CIOSELINE IS A GENERAL PURPOSE PROCEDURE USED WHEN THE *
{* END OF A LINE IS EXPECTED BUT IT IS NOT KNOWN WHETHER *
{* CR NOT THE LINE BUFFER INDEX IS IN FRONT OF OR AT THE *
{* END OF LINE/FILE CHAR; MAY ALSO BE A COMMENT PRIOR TO. *
(*-----*)

```

PROCEDURE CIOSELINE;

```

BEGIN
  IF NOT (TCKNUM IN TRAILTOKS) THEN
    SCAN (TCKNUM);
  IF NOT (TCKNUM IN TRAILTOKS) THEN
    PERRCH
END;
(* CIOSELINE *)

```

```

{*****}
{*
{*          SYMBOL TABLE MANAGEMENT ROUTINES
{*
{*****}

{*--11-----}
{* HASHVAL RETURNS THE HASH VALUE OF THE STRING CONTAINED *
{* IN THE ACCUM (SUM OF ORD VALUES OF CHARS MOD HASHEASE) *
{*-----}

FUNCTION HASHVAL (ACCUM:TKSTRING; ACCINX:INTEGER) : HASHENG;
VAR  HASHSUM, I : INTEGER;
BEGIN
  HASHSUM := 0;
  FOR I := 1 TO ACCINX DO
    HASHSUM := HASHSUM + ORD (ACCUM(.I.));
  HASHVAL := HASHSUM MOD (HASHEASE + 1)
END;                                     (* HASHVAL *)

{*****}

{*--12-----}
{* GETSLOT HASHES ID STRING IN ACCUM INTO SYMBOL TABLE; *
{* INSERTS NEW SLOT INTO CORRECT HASH BUCKET AND ENTERS *
{* IDENTIFIER NAME INTO IDENT FIELD; RETURNS SLOT POINTER *
{* TO THE NEW SLOT BUT VARIANT TAG TYP IS UNDECLARED. *
{*-----}

FUNCTION GETSLOT (ACCUM:TKSTRING; ACCINX:ACCRNG) : SLOTPTR;
VAR  CURHASH : HASHENG;
     IDSLCT  : SLOTPTR;
BEGIN
  CURHASH := HASHVAL (ACCUM, ACCINX);
  NEW (IDSLCT);
  IDSLCT?.IDENT := ACCUM;
  IDSLCT?.SLOT := BUCKET(.CURHASH.);
  BUCKET(.CURHASH.) := IDSLCT;
  GETSLOT := IDSLCT
END;                                     (* GETSLOT *)

{*****}

{*--13-----}
{* FN_CHK RETURNS TRUE/FALSE AS TO WHETHER OR NOT AN *
{* ACCUM STRING DESIGNATES A USER DEFINED FUNCTION OR NOT *
{*-----}

FUNCTION FN_CHK (ACCUM : TKSTRING) : BOOLEAN;
BEGIN
  FN_CHK := FALSE;
  IF (UPCASE {ACCUM { .1. }} = 'F') AND
     (UPCASE {ACCUM { .2. }} = 'N') AND
     (ACCUM { .3. } = USCORE)
  THEN
    FN_CHK := TRUE
END;                                     (* FN_CHK *)

{*****}

```

```

(*-14-----*)
(* NEWREG RETURNS THE VALUE OF THE NEXT REGISTER WHICH *)
(* AVAILABLE FOR USE AS VARIABLE STORAGE; RESERVED REGS *)
(* ARE SKIPPED AND A COUNT IS MAINTAINED TO IDENTIFY THE *)
(* POINT AT WHICH TOO MANY REGISTERS HAVE BEEN USED. *)
(* NOTE THAT IF THE LAST REG IS RESERVED AN OVERFLOW WILL *)
(* NOT BE DETECTED AND PROCESSING WILL CONTINUE. NOTE *)
(* ALSO THAT A MEMORY OVERFLOW DOES NOT STOP THE PARSER *)
(* FROM ANALYSIS AND CODE GENERATION, HOWEVER THE REG *)
(* SUMMARY MAY NOT REFLECT ACCURATE REGISTER INFO. *)
(*-----*)

```

```

FUNCTION NEWREG : INTEGER;
BEGIN
  WHILE NEXTREG IN RESERVE_REG DO
    NEXTREG := NEXTREG + 1;
    NEWREG := NEXTREG;
    IF NEXTREG = (REGBASE + STARTREG) THEN
      BEGIN
        (* NOTE THAT IF LAST REG IS RESERVED, THEN *)
        PWARN; (* THIS WARNING WILL NOT BE TRIGGERED *)
        WRITE (LISTP, '***** MEMORY OVERFLOW...> ');
        WRITE (LISTP, REGBASE:1);
        WRITELN (LISTP, ' VARIABLE NAMES IN USE...REUSING. ');
        NEXTREG := STARTREG (* RESET THE REGISTER STACK *)
      END
    ELSE
      NEXTREG := NEXTREG + 1
  END;
  (* NEWREG *)

```

```

(*-15-----*)
(* NEWLBL RETURNS THE KEY CODE FOR THE NEXT LABEL ON THE *)
(* LABEL STACK; IF A LABEL STACK OVERFLOW OCCURS, THE *)
(* SUMMARY ITEM WHICH INDICATES NUMBER OF LABELS USED MAY *)
(* NOT REFLECT ACCURATE INFORMATION. *)
(*-----*)

```

```

FUNCTION NEWLBL : LBLNG;
BEGIN
  NEWLBL := CLABEL(.LBLCT.);
  LBLCT := LBLCT + 1;
  IF LBLCT = LBLBASE + 1 THEN
    BEGIN
      PWARN;
      WRITE (LISTP, '***** LABEL OVERFLOW...> ');
      WRITE (LISTP, LBLBASE:1);
      WRITELN (LISTP, ' IN USE...RESET TO 0. ');
      LBLCT := 1
    END;
  END;
  (* NEWLBL *)

```

```
(*-16-----*)
{* FNSTACKLOCK SEARCHES THE FNF ACTIVATION STACK FOR THE *}
{* IDENTIFIER IN ACCUM AND RETURNS A POINTER TO ITS SLOT *}
{* IF IT IS A FORMAL PARAMETER IN AN ACTIVE FNP, OTHER- *}
{* WISE THE POINTER RETURNED IS NIL. *}
*-----*)
```

```
FUNCTION FNSTACKLOOK (ACCUM : TKSTRING) : SLOTPTR;
```

```
VAR FLOCK, PARMPTR, TPARMPTR : SLOTPTR;
    FOUND : BOOLEAN;
```

```
BEGIN
```

```
  FLCCK := FNSTACK;
```

```
  POUND := FALSE;
```

```
  PARMPTR := NIL;
```

```
  TPARMPTR := NIL;
```

```
  WHILE (FLCCK <> NIL) AND (NOT FOUND) DO
```

```
    BEGIN
```

```
      PARMPTR := FLOCK@.FNP;
```

```
      IF PARMPTR <> NIL THEN
```

```
        REPEAT
```

```
          TPARMPTR := PARMPTR;
```

```
          FCUND := (PARMPTR@.IDENT = ACCUM);
```

```
          PARMPTR := PARMPTR@.SLOT
```

```
        UNTIL (FOUND) OR (PARMPTR = NIL);
```

```
      FLCCK := FLOCK@.FNPLINK
```

```
    END;
```

```
  FNSTACKLOCK := PARMPTR;
```

```
  IF FCUND THEN
```

```
    FNSTACKLOCK := TPARMPTR
```

```
END;
```

```
(* FNSTACKLOCK *)
```

```
(*****)
```

```

(*-17-----*)
(* IDICCKUP FIRST SEARCHES ACTIVE FNP STACK (FORMAL      *)
(* PARAMETERS), THEN THE SYM TBL UNTIL IT FINDS THE     *)
(* IDENTIFIER NAME CURRENTLY IN ACCUM; RETURNS POINTER TO *)
(* THE SICT FOR THAT IDENTIFIER; CREATES AND ENTERS A    *)
(* SLCT FOR THE IDENTIFIER IF IT DOES NOT YET EXIST.     *)
(*-----*)
FUNCTION IDLOOKUP (ACCUM:TKSTRING; ACCINX:ACCRNG) : SLOTPTR;
VAR  HICOK, TLOCK : SLOTPTR;

(*-----*)
(*-17-01-----*)
(* ENTERID INSERTS AN APPROPRIATE SLOT INTO THE SYM TBL  *)
(* FOR THE IDENTIFIER IN ACCUM AND RETURNS A POINTER TO  *)
(* THAT SICT.                                           *)
(*-----*)
FUNCTION ENTERID (ACCUM:TKSTRING; ACCINX:ACCRNG) : SLOTPTR;
VAR  IDSICT : SLCTPTR;
BEGIN
  IDSLOT := GETSLOT (ACCUM, ACCINX);
  IF NOT FN CHK (ACCUM) THEN
    WITH IDSLOT@ DO
      BEGIN
        TYP := VARID;
        REGNO := NEWREG;
        AUXREG1 := -1; (* USED FOR INDEX VARS TO *)
        AUXREG2 := -1; (* FOR-NEXT LOOPS ONLY. *)
        WRITE (NAMEF, ' ':8);
        ZERCFAD (NAMEF, REGNO, 2);
        WRITELN (NAMEF, ' ':3, IDENT)
      END
    ELSE
      WITH IDSLOT@ DO
        BEGIN
          TYP := FNPID;
          FNP := NIL;
          FNPLINK := NIL;
          FNREGNO := NEWREG;
          LEL := NEWLBI;
          WRITELN (NAMEF, ' ':5, FNREGNO:5, ' ':3, IDENT)
        END;
      ENTERID := IDSLOT
    END;
  (* ENTERID *)
  (*-----*)

```

```

BEGIN (* IDLOOKUP MAIN *)
  TLCK := FNSTACKLOCK (ACCUM);
  IF TLOCK = NIL THEN
    BEGIN
      HLOCK := BUCKET (.HASHVAL (ACCUM, ACCINX) .);
      IF HLOCK <> NIL THEN
        BEGIN
          TLOCK := HLOOK;
          HLOOK := HICOK@.SLOT;
          WHILE (HLOCK<>NIL) AND (TLOCK@.IDENT<>ACCUM) DO
            BEGIN
              HLOOK := HLOOK@.SLOT;
              TLOCK := TLOCK@.SLOT;
            END;
          IF TLOCK@.IDENT = ACCUM THEN
            IDLOCKUP := TLOCK
          ELSE
            IDLOCKUP := ENTERID (ACCUM, ACCINX)
          END
        ELSE
          IDLCOKUP := ENTERID (ACCUM, ACCINX)
        END
      ELSE
        IDLCCKUP := TLOCK
      END;
    END;
  END;

```

(* IDLOOKUP *)

```

{*****}
{*
{*      CCODE DATA STRUCTURE MANAGEMENT ROUTINES
{*
{*****}

{*--18-----*}
{* NEWCCODE RETURNS A POINTER TO A NEW CODE DATA NODE;
{* FIELDS ARE INITIALIZED TO A DEFAULT VALUE (EXCEPT KEY)
{*-----*}

FUNCTION NEWCODE (OPCCODE : INTEGER) : CODEPTR;
VAR CP : CODEPTR;

BEGIN
  NEW (CP);
  CP@.SECF := NIL;
  CP@.JMFP := NIL;
  CP@.EAKF := NIL;
  CP@.ADDR := -1;
  CP@.KEY  := OPCCODE;
  NEWCODE := CP;
END;                                     (* NEWCODE *)

{*****}

{*--19-----*}
{* PUTKEY INSERTS A CODE DATA NODE INTO THE CODE DATA
{* STRUCTURE IN FRONT OF CPCUR BUT AFTER ENDCP WHICH IS
{* TO SAY THAT THE NODE IS THE NEW END CODE DATA NODE.
{*-----*}

PROCEDURE PUTKEY (VAR HOLDF : CODEPTR);

BEGIN
  IF HOLDF = NIL THEN
    HOLDF := NEWCODE (-1);
  ENDCP@.SEQP := HOLDF;
  HOLDF@.SECF := CPCUR;
  ENDCP := ENDCP@.SEQP
END;                                     (* PUTKEY *)

{*****}

{*--20-----*}
{* GENKEY FILLS THE CPCUR REFERENCED NODE WITH THE OPCCODE
{* PASSED TO IT, AND ATTACHES ANOTHER EMPTY NODE TO THE
{* CODE DATA STRUCTURE AFTER THE ONE JUST FILLED.
{*-----*}

PROCEDURE GENKEY;                                     (* FORWARD DECLARATION WITH *)
VAR CP : CODEPTR;                                     (* ERROR HANDLING ROUTINES *)

BEGIN
  CPCUR@.KEY := OPCCODE;
  CPCUR@.SECF := NEWCODE (-1);
  CPCUR := CPCUR@.SECF;
  ENDCP := ENDCP@.SECF
END;                                     (* GENKEY *)

{*****}

```

```
(*--21-----*)
(* INSERTKEY IS USED AFTER ALL CODE HAS BEEN GENERATED TO *)
(* INSERT LABELS TO GCSUB REFERENCES; NOTE THAT ENDCP *)
(* MUST NOT BE MOVED OR USED BY THIS PROCEDURE SINCE IT *)
(* NOW RESIDES AT THE END OF THE CODE STRUCTURE AND WILL *)
(* BE USED TO FLAG THE END OF THE STRUCTURE TO TRAVERSING *)
(* PCINTER PROCEDURES. *)
(*-----*)
```

```
PROCEDURE INSERTKEY (CPCODE : INTEGER; VAR LOCUS : CODEPTR);
VAR CP : CCDEPTR;
```

```
BEGIN
  CP := NEWCODE (OPCODE);
  CP@.SECF := LOCUS;
  LOCUS := CP
END; (* INSERTKEY *)
```

```
(*****)
```

```
(*--22-----*)
(* GETNEWHDR CONSTRUCTS AND RETURNS A POINTER TO A LINE *)
(* AND CODE DATA NODE PAIR; ALL FIELDS OF BOTH NODES ARE *)
(* INITIALIZED; THIS PAIR OF NODES IS USED TO HEAD THE *)
(* CHAIN OF CODE GENERATED FOR A PARTICULAR BASIC SOURCE *)
(* LINE. (NOTE THAT NEWCODE INITIALIZES CODE NODE) *)
(*-----*)
```

```
FUNCTION GETNEWHDR (LINUM : BASLINRNG) : LINEPTR;
```

```
VAR LP : LINEPTR;
    CP : CODEPTR;
```

```
BEGIN
  NEW (LP);
  LP@.CPTR := NEWCODE (-1);
  LP@.LINC := LINUM;
  LP@.LPTR := NIL;
  GETNEWHDR := LP
END; (* GETNEWHDR *)
```

```
(*****)
```

```
(*--23-----*)
(* PUTNEWHDR INSERTS A HEADER (LINE/CODE DATA NODE PAIR) *)
(* INTO THE CODE DATA STRUCTURE AT THE POSITION OF THE *)
(* CURRENT LINE MARKING POINTER, LPCUR. NOTE THAT THE *)
(* BAKP IS USED HERE TO MARK THE LOCATION OF THE START OF *)
(* CODE CORRESPONDING TO A NEW SOURCE LINE. *)
(*-----*)
```

```
PROCEDURE PUTNEWHDR (VAR LPCUR, LP : LINEPTR);
```

```
BEGIN
  LP@.LPTR := LPCUR@.LPTR;
  LPCUR@.LPTR := LP;
  LPCUR := LP;
  ENDCP@.SECF := LP@.CPTR;
  CPCUR := LP@.CPTR;
  CPCUR@.EAKP := ENDCP
END; (* PUTNEWHDR *)
```

```
(*****)
```



```

(*-24-----*)
{* SETLINE COORDINATES THE SET UP OF THE CODE DATA *}
{* STRUCTURE FOR THE BEGINNING OF CODE GENERATED FOR A *}
{* NEW WEASIC SOURCE LINE; IF LINE NUMBER HEADER NODES *}
{* ALREADY EXIST BECAUSE FORWARD JUMP REFERENCES REQUIRED *}
{* THEIR EXISTENCE, THEN THESE NODES ARE CHECKED FIRST *}
{* FOR THE CURRENT LINE NUMBER. *}
-----*)

```

```
PROCEDURE SETLINE (VAR LPCUR, LP : LINEPTR);
```

```
VAR LNUMBER : BASLINRNG;
```

```
BEGIN
```

```
  LNUMBER := LINUM;
```

```
  IF TOKNUM = ENDLINTOK THEN
```

```
    LNUMBER := LIINUM;
```

```
  IF LPCUR@.LPTR <> NIL THEN
```

```
    BEGIN (* IF LINE DATA NODES *)
```

```
      IF (LPCUR@.LPTR@.LINO) = LNUMBER THEN (* EXIST FROM JUMP REFS. *)
```

```
        BEGIN (* CURRENT LINE NUMBER IS *)
```

```
          LPCUR := LPCUR@.LPTR; (* THE NEXT ONE IN CHAIN *)
```

```
          ENDCP@.SEOP := LPCUR@.CPTR;
```

```
          CPCUR := LPCUR@.CFTR;
```

```
          CPCUR@.BAKP := ENDCP
```

```
        END
```

```
      ELSE
```

```
        BEGIN
```

```
          LF := GETNEWHDR (LNUMBER);
```

```
          PUTNEWHDR (LPCUR, LP)
```

```
        END
```

```
      END
```

```
    ELSE
```

```
      BEGIN
```

```
        LP := GETNEWHDE (LNUMBER);
```

```
        PUTNEWHDR (LPCUR, LP)
```

```
      END
```

```
END;
```

```
(* SETLINE *)
```

```
(*****)
```

```

(*-25-----*)
{* SETJMPEXT SETS THE EXTERNAL JMP PTR FROM CFCUR@.JMPP *}
{* TO THE WEASIC LINUM INDICATED BY THE GOTO OR GOSUE *}
*-----*

```

```

PROCEDURE SETJMPEXT (LINUM : BASLINRNG) ;

```

```

BEGIN
  IF LINUM > LPCUR@.LINO THEN
    BEGIN
      (* FORWARD JUMP *)
      LPLEAD := LPCUR@.LPTR;
      LPTRAIL := LPCUR@;
      WHILE (LPTRAIL@.LINO < LINUM) AND (LPLEAD <> NIL) DO
        BEGIN
          LPLEAD := LPLEAD@.LPTR;
          LPTRAIL := LPTRAIL@.LPTR;
        END;
      IF LPTRAIL@.LINO = LINUM THEN
        BEGIN
          CFCUR@.JMPP := LPTRAIL@.CPTR; (* SET JMPP PTR *)
          LPTRAIL@.CPTR@.ADDR := 0 (* MARK JMPP TERMINAL *)
        END
      ELSE
        BEGIN
          LPTRAIL@.LPTR := GETNEWHDR (LINUM);
          LASTLP := LPTRAIL@.LPTR;
          CFCUR@.JMPP := LASTLP@.CPTR; (* SET JMPP PTR *)
          LASTLP@.CPTR@.ADDR := 0 (* MARK JMPP TERMINAL *)
        END
      END
    END
  ELSE
    BEGIN
      (* BACKWARD JUMP *)
      LPLEAD := FIRSTIP;
      WHILE (LPLEAD@.LINO <> LINUM) AND (LPLEAD <> NIL) DO
        LPLEAD := LPLEAD@.LPTR;
      IF LPLEAD@.LINO = LINUM THEN
        BEGIN
          CFCUR@.JMPP := LPLEAD@.CPTR; (* SET JMPP PTR *)
          LPLEAD@.CPTR@.ADDR := 0 (* MARK JMPP TERMINAL *)
        END
      END
    END
  END (* ASSUMES THAT A LINO = LINUM ALWAYS EXISTS ! *)
END; (* SETJMPEXT *)

```

```

{*****}
{*
      FUNCTION CALL ROUTINES
*}
{*****}

(*-40-*)
PROCEDURE PEXPR;                                FORWARD;

(*-41-*)
PROCEDURE PPRIMARY;                              FORWARD;

{*****}

{*--26-----}
{* GENFNC GENERATES THE CODE FOR SHORT, SIMPLE CALCULATOR *}
{* ARITHMETIC FUNCTIONS ("QUICK" FUNCTIONS). *}
{*-----}

PROCEDURE GENFNC (OPND : SLOTPTR);
VAR   I : 1..FNQLEN;
BEGIN
  SCAN (TCKNUM);
  IF TCKNUM <> OPARENTOK THEN                    (* ' ( ' *)
    PERRCR
  ELSE
    BEGIN
      SCAN (TOKNUM);
      PEXPR;
      I := 1;
      WITH CFND@ DO
        REPEAT
          GENKEY (FNQ(.I.));
          I := I + 1
        UNTIL (I >= FNQLEN) OR (FNQ(.I.) = K_NOP)
    END
  END;                                           (* GENFNC *)
{*****}

{*--27-----}
{* CHKFNLIST SEARCHES THE FNL USE LIST TO DETERMINE IF *}
{* THE FNL HAS EVER BEEN CALLED BEFORE; IF NOT, THEN IT *}
{* MUST BE ADDED TO THE USE LIST. *}
{*-----}

PROCEDURE CHKFNLIST (VAR IDSLOT : SLOTPTR);
VAR   LISTPTR, HCLDPTR : SLOTPTR;
      USED : BOOLEAN;
BEGIN
  LISTPTR := FNLLIST;                            (* GET THE LONG FN LIST *)
  IF LISTPTR <> NIL THEN                          (* TRAVERSE THE FNLLIST *)
    REPEAT
      USED := (IDSLOT@.IDENT = LISTPTR@.IDENT);
      LISTPTR := LISTPTR@.FNLLINK;
    UNTIL (USED) OR (LISTPTR = NIL);
  IF LISTPTR = NIL THEN                          (* IF NOT FOUND ON LIST, THEN *)
    BEGIN                                        (* ADD THIS LONG FUNCTION TO THE FNLLIST *)
      HCLDPTR := FNLLIST;
      FNLLIST := IDSLOT;
      IDSLOT@.FNLLINK := HCLDPTR
    END
  END;                                           (* CHKFNLIST *)
{*****}

```

```

(*-28-----*)
{* GENFNL GENERATES THE SBR CALL TO A PNL BODY AND THEN *}
{* CALLS CN PROCEDURE CHKFNLLIST TO SEE IF THE PNL HAS *}
{* BEEN USED BEFORE. *}
(*-----*)

```

```
PROCEDURE GENFNL (VAR IDSLOT : SLOTPTR) ;
```

```

BEGIN
  SCAN (TOKNUM) ;
  IF TOKNUM <> OPARENTOK THEN (* ' ( ' *)
    PERRCR
  ELSE
    BEGIN
      SCAN (TOKNUM) ;
      PEXFR ;
      GENKEY (K STO) ;
      GENKEY (PNIREG) ;
      GENKEY (K SBR) ;
      GENKEY (IDSLOT.LBL) ;
      CHKFNLLIST (IDSLOT)
    END
  END ;
(* GENFNL *)

```

```
(*****)
```

```

(*-29-----*)
{* NEWPARM IS CALLED WHEN A FIRST CALL TO A PARAMETER FN *}
{* IS ENCOUNTERED IN THE WBASIC SOURCE FILE; AT THIS TIME *}
{* NO REGISTERS HAVE BEEN DESIGNATED FOR THE FORMAL FN *}
{* PARMS; THIS PROCEDURE CREATES THE SYM TBL ENTRY AND *}
{* DESIGNATES THE RESPECTIVE REGISTER FOR THE NEW PARM *}
{* ENCOUNTERED IN THE FN CALL; NOTE THAT THE FORMAL PARM *}
{* IDENT NAME IS NOT YET KNOWN AND WILL BE ENTERED ONLY *}
{* AFTER THE FN DEFINITION STATEMENT IS ENCOUNTERED LATER *}
(*-----*)

```

```
FUNCTION NEWPARM : SICTPTR;
```

```
VAR PARMSICT : SLOTPTR;
```

```

BEGIN
  NEW (PARMSLOT) ;
  WITH PARMSLOT DO
    BEGIN
      TYF := VARID ;
      REGNO := NEWREG ;
      WRITELN (NAMEP, ' ':5, REGNO:5, ' ':5, ' (FN PARAMETER) ');
      SICT := NIL
    END ;
  NEWPARM := PARMSLOT
END ;
(* NEWPARM *)

```

```
(*****)
```

```

(*-30-----*)
(* GENPARM GENERATES THE FORMAL PARAMETER LIST FOR *)
(* PARAMETER FUNCTIONS; IT ALSO GENERATES THE CODE WHICH *)
(* WILL EVALUATE AN ACTUAL PARAMETER AND STORE IT IN THE *)
(* FORMAL PARAMETER STORAGE LOCATION PRIOR TO FN EXECUTE. *)
(*-----*)

```

```

PROCEDURE GENPARM (VAR IDSLOT : SLOTPTR);
VAR PARMPTR : SLOTPTR;
BEGIN
  SCAN (TOKNUM);
  IF TOKNUM = OPARENICK THEN
    BEGIN
      PARMPTR := IDSLOT@.FNP;
      IF PARMPTR = NIL THEN
        BEGIN
          PARMPTR := NEWPARM;
          IDSLOT@.FNP := PARMPTR;
        END;
      REPEAT
        SCAN (TOKNUM);
        PEXPR; (* STOP AT EACH ',' AND LAST ')' *)
        GENKEY (K_STC);
        GENKEY (PARMPTR@.REGNO);
        IF TOKNUM = CCMMATOK THEN
          BEGIN
            IF PARMPTR@.SLOT = NIL THEN
              PARMPTR@.SLOT := NEWPARM;
            PARMPTR := PARMPTR@.SLOT (* NEXT PARAMETER *)
          END;
      UNTIL (TOKNUM = CPARENTOK) (* PEXPR WILL FIND ')' *)
        OR (TOKNUM IN TRAILTOKS) (* OR WILL FIND END *)
    END;
END; (* GENPARM *)

```

```

(*****)

```

```

(*-31-----*)
(* GENFNP GENERATES CCDE SEQUENCE WHICH CALLS A PARM FNP. *)
(*-----*)

```

```

PROCEDURE GENFNP (VAR IDSLOT : SLOTPTR);
BEGIN
  GENFARM (IDSLOT);
  GENKEY (K_CPAREN);
  GENKEY (K_SBR);
  GENKEY (IDSLOT@.LBI);
  GENKEY (K_CPAREN);
END; (* GENFNP *)

```

```

(*****
{*
{*          FUNCTION DEFINITION ROUTINES
{*
*****

```

```

{*--32-----*
{* PUSHFN PUSHES A FNP SLOT ONTO THE FNP ACTIVATION STACK *
{*-----*

```

```

PROCEDURE PUSHFN (VAR FNSLOT : SLOTPTR);

```

```

BEGIN
  FNSLOT@.FNPLINK := FNSTACK;
  FNSTACK := FNSLOT;
  FNSTACKCT := FNSTACKCT + 1;
  IF FNSTACKCT > FNSTACKLIM THEN
    BEGIN
      FWARN;
      WRITE (LISTP, '***** SER STACK OVERFLOW...> ');
      WRITELN (LISTP, FNSTACKLIM:1, ' RETURN ADDRESSES. ');
    END
  END;
(* PUSHFN *)

```

```

(*****
{*--33-----*
{* PCPFN POPS A FNP SLOT OFF TOP OF FNP ACTIVATION STACK. *
{*-----*

```

```

PROCEDURE PCPFN;

```

```

VAR HOLDFTR : SLOTPTR;

```

```

BEGIN
  HOLDFTR := FNSTACK;
  FNSTACK := FNSTACK@.FNPLINK;
  HOLDFTR@.FNPLINK := NIL;
  FNSTACKCT := FNSTACKCT - 1;
  IF FNSTACKCT < 0 THEN
    BEGIN
      FWARN;
      WRITE (LISTP, '***** ATTEMPT TO POP RETURN ADDR ');
      WRITELN (LISTP, 'FROM EMPTY STACK...RESET CT = 0');
      FNSTACKCT := 0;
    END
  END;
(* POPFN *)

```

```

(*****

```



```

{* -36-----*
{* PDEF GENERATES THE CODE WHICH DEFINES THE SCOPE AND *
{* VISIBILITY FOR VARIABLE NAMES; RESETS THE VALUE OF THE *
{* REGISTER IN WHICH THE FUNCTION VALUE IS RETURNED. *
{*-----*

```

```
PROCEDURE PDEF;
```

```
VAR FNSICT : SLCTPTR;
```

```

BEGIN
  SCAN (TCKNUM);
  IF NOT FN_CHK (ACCUM) THEN
    PERRCR
  ELSE
    BEGIN
      FNSICT := IDLOCKUP (ACCUM, ACCINX);
      GENKEY (K_LBL);
      GENKEY (FNSLOT@.LBL);
      GENKEY (K_ZERO);
      GENKEY (K_STO);
      GENKEY (FNSLOT@.FNREGNO);
      PUSHFN (FNSLOT);
      SCAN (TOKNUM);
      IF TCKNUM = OPARENTOK THEN (* LOOKING FOR PARMS *)
        FILIPARMIDS (FNSLOT);
      IF TCKNUM = EQUALTOK THEN (* LOOK FOR ONE LINE FN *)
        BEGIN
          SCAN (TOKNUM);
          PEXPR;
          PFNEND (* GENERATE THE RETURN FROM ONE LINE FN *)
        END;
      CLCSELINE
    END
  END;
END; (* PDEF *)

```

```

(*****

```

```

{* -37-----*
{* GETFNIS IS CALLED AFTER ALL OTHER CODE HAS BEEN *
{* GENERATED; IT GENERATES THE CODE FOR THE BODIES OF ALL *
{* BUILT-IN LONG FUNCTIONS WHICH HAVE BEEN CALLED AT *
{* LEAST CNCE AND ARE, THUS, ON THE FNLLIST. *
{*-----*

```

```
PROCEDURE GETFNIS;
```

```

VAR LISTPTR : SLOTPTS;
    I : 1..FNLLEN + 1;

```

```

BEGIN
  LISTPTR := FNLLIST; (* GET LONG FN LIST *)
  WHILE LISTPTR <> NIL DO
    BEGIN (* INSERT CODE FOR NEXT LONG FN ON FNLLIST *)
      GENKEY (K_LBL);
      GENKEY (LISTPTR@.LBL);
      I := 1;
      WHILE (I <= FNLLEN) AND
        (LISTPTR@.FNL(.I.) <> K_NOP) DO
        BEGIN
          GENKEY (LISTPTR@.FNL(.I.));
          I := I + 1
        END;
      GENKEY (K_INVSER);
      LISTPTR := LISTPTR@.FNLLINK (* NEXT LONG FN ON LIST *)
    END
  END;
END; (* GETFNIS *)

```



```

(*****
*
*           EXPRESSION GENERATOR ROUTINES
*
*****
)

```

```

(*-38-----*)
* GENID GENERATES CCDE FOR A VARIABLE OR FUNCTION IDENT. *
*-----*
)

```

```

PROCEDURE GENID;
VAR   CFND : SLOTPTR;
BEGIN
  OPND := IDLOOKUP (ACCUM, ACCINX);
  CASE OPND.TYP OF
    VVARIC : BEGIN
      IF OPND.REGNO = -314 THEN (* REGNO FOR PI *)
        GENKEY (K_PI)
      ELSE
        BEGIN
          GENKEY (K_RCL);
          GENKEY (OPND.REGNO)
        END
      END;
    FNOID  : GENFNO (CFND);
    FNLIC  : GENFNL (CFND);
    FNPID  : GENFNP (CFND);
  END (* CASE *)
END;

```

```
(* GENID *)
```

```

(*****
)

```

```

(*-39-----*)
* GENNUM GENERATES II-59 EQUIVALENT CODE FOR A LITERAL *
* NUMERIC (BOTH INTEGER AND REAL). *
*-----*
)

```

```

PROCEDURE GENNUM;
VAR   I, DECPTLOC, ESIGNLOC : INTEGER;
BEGIN
  DECPTLOC := LCCUNT + 1;
  FOR I := 1 TO LCOUNT DO
    GENKEY (TRANSDIGIT (ACCUM(.I.)));
    IF RCCUNT > 0 THEN
      BEGIN
        GENKEY (K_DECPT);
        FOR I := (DECPTLOC + 1) TO (DECPTLOC + RCOUNT) DO
          GENKEY (TRANSDIGIT (ACCUM(.I.)))
        END;
      END;
    IF ECCUNT > 0 THEN
      BEGIN
        ESIGNLOC := LCCUNT + 1 + RCOUNT + 2;
        IF ACCUM(.DECPTLOC.) <> PERIOD THEN
          ESIGNLOC := ESIGNLOC - 1;
          GENKEY (K_EE);
          IF ACCUM(.ESIGNLOC.) = '-' THEN
            GENKEY (K_NEG);
          FOR I := (ESIGNLOC + 1) TO (ESIGNLOC + ECOUNT) DO
            GENKEY (TRANSDIGIT (ACCUM(.I.)))
          END;
        END;
      END;
    END;
  END;
END;

```

```
(* GENNUM *)
```

```

(*****
)

```

```
(*--40-----*)
{* PEXPR PARSSES AND GENERATES CODE FOR EXPRESSIONS. *}
(*-----*)
```

```
PROCEDURE PEXPR; (* FORWARD DECLARATION WITH *)
(* FUNCTION CALL ROUTINES *)
```

```
BEGIN
  GENKEY (K_OPAREN);
  PPRIMARY;
  WHILE TCKNUM IN BINOPTCKS DO
    BEGIN
      CASE TCKNUM OF
        PIUSTOK : GENKEY (K_ADDOP); (* '+' *)
        MINUSTOK : GENKEY (K_SUBCP); (* '-' *)
        MULTCK : GENKEY (K_MULTOP); (* '*' *)
        DIVTOK : GENKEY (K_DIVOP); (* '/' *)
        EXPTOK : GENKEY (K_POWR); (* '**' *)
      END; (* CASE *)
      IF NOT (TCKNUM IN TRAILTOKS) THEN BEGIN
        SCAN (TOKNUM);
        PPRIMARY
      END
    END;
  GENKEY (K_CPAREN)
END; (* PEXPR *)
```

```
(*****)
```

```
(*--41-----*)
{* PPRIMARY PARSSES AND GENERATES CODE FOR A PRIMARY ITEM *}
{* EXPECTED AS PART OF EXPRESSIONS. *}
(*-----*)
```

```
PROCEDURE PPRIMARY; (* FORWARD DECLARATION WITH *)
(* FUNCTION CALL ROUTINES *)
```

```
BEGIN
  CASE TCKNUM OF
    PIUSTCK : BEGIN (* '+' *)
      SCAN (TOKNUM);
      PPRIMARY
    END;
    MINUSTCK : BEGIN (* '-' *)
      SCAN (TOKNUM);
      PPRIMARY;
      GENKEY (K_NEG)
    END;
    OPARENTCK : BEGIN (* '(' *)
      SCAN (TOKNUM);
      PEXPR;
      SCAN (TOKNUM)
    END;
    IDENTCK : BEGIN
      GENID;
      IF (TOKNUM=IDENTCK) OR (TOKNUM=CPARENTCK)
      THEN SCAN (TOKNUM)
    END;
    NUMBERTCK : BEGIN
      GENNUM;
      SCAN (TOKNUM)
    END
  END (* CASE *)
END; (* PPRIMARY *)
```

```
(*****)
```

```
{*-42-----*}
{* PCCNDITICN PARSES AND GENERATES CODE FOR BOOLEAN EXPRS *}
{*-----*
```

```
PROCEDURE FCCNDITION;
```

```
VAR REICP : TOKENRNG;
    INVERT : BOOLEAN;
```

```
BEGIN
```

```
    INVERT := FALSE;
```

```
    SCAN (TCKNUM);
```

```
(* SCAN FOR "NOT" *)
```

```
    IF TOKNUM = NOTTOK THEN
```

```
        BEGIN
```

```
            INVERT := TRUE;
```

```
            SCAN (TOKNUM)
```

```
        END;
```

```
    PEXER;
```

```
    GENKEY (K_X T);
```

```
    IF INVERT THEN
```

```
        CASE TOKNUM OF
```

```
            EQUALTCK : TOKNUM := NOTEQTOK;
```

```
            NCTEQTOK : TOKNUM := EQUALTOK;
```

```
            GTTCK : TOKNUM := LITECTOK;
```

```
            GTECTCK : TOKNUM := LTTOK;
```

```
            LITCK : TOKNUM := GTECTOK;
```

```
            LITECTCK : TOKNUM := GTTOK;
```

```
        END; (* CASE *)
```

```
    REICP := TCKNUM;
```

```
(* BEGIN NEXT EXPR *)
```

```
    SCAN (TCKNUM);
```

```
    PEXER;
```

```
    CASE REICP OF
```

```
        EQUALTOK : BEGIN
```

```
            GENKEY (K_INV);
```

```
            GENKEY (K_IPXEQT)
```

```
        END;
```

```
        NCTEQTOK : GENKEY (K_IPXEQT);
```

```
        GTTCK : GENKEY (K_IPXGET);
```

```
        GTEQTCK : BEGIN
```

```
            GENKEY (K_X T);
```

```
            GENKEY (K_INV);
```

```
            GENKEY (K_IPXGET)
```

```
        END;
```

```
        LITECTCK : BEGIN
```

```
            GENKEY (K_INV);
```

```
            GENKEY (K_IPXGET)
```

```
        END;
```

```
        LITCK : BEGIN
```

```
            GENKEY (K_X T);
```

```
            GENKEY (K_IPXGET)
```

```
        END
```

```
    END (* CASE *)
```

```
(* PCONDITION *)
```

```
END;
```

```

{*****}
{*
      LOOPING ROUTINES
*}
{*****}

```

```

{*--43-----*}
{* PUSHCCDE PUSHES THE RCODE DATA NODE ONTO THE LOOP/IF *}
{* STACK DESIGNATED BY STACK. *}
{*-----*}

```

```

PROCEDURE PUSHCODE (RCODE : CODEPTR; VAR STACK : CODEPTR);
BEGIN
  RCODE@.SEQP := STACK;
  STACK := RCODE
END;
(* PUSHCODE *)

```

```

{*****}
{*--44-----*}
{* POPCODE PCPS AND RETURNS THE CODE DATA NODE ON THE TOP *}
{* OF THE LCCP/IF STACK DESIGNATED. *}
{*-----*}

```

```

FUNCTION PCPCODE (VAR STACK : CODEPTR) : CODEPTR;
BEGIN
  IF STACK = NIL THEN
    BEGIN
      WRITE (LISTF, '***** INCORRECT NESTING OF CONTRCL');
      WRITEIN (LISTF, ' STATEMENTS (IF, FOR, OR LOOP).');
      PERROB;
      POPCODE := NIL
    END
  ELSE
    BEGIN
      POPCODE := STACK;
      STACK := STACK@.SEQP
    END
  END;
(* POPCODE *)

```

```

{*****}
{*--45-----*}
{* SETFWDJMP IS USED TO SET THE JUMP POINTER (JMPP) OF *}
{* THE CURRENT CODE DATA NODE POINTING TO THE MOST RECENT *}
{* CODE DATA NODE ON THE DESIGNATED LOOP/IF STACK; THE *}
{* POTENTIAL ABSOLUTE ADDRESS SPACE IS GENERATED WITH THE *}
{* ASSUMPTION THAT THE CODE DATA NODE IN THE STACK TO *}
{* WHICH THE FIRST ADDRESS SPACE NODE IS POINTING WILL *}
{* LATER BE POPPED AND INSERTED INTO THE CODE AT THE *}
{* APPROPRIATE POSITION. *}
{*-----*}

```

```

PROCEDURE SETFWDJMP (STACK : CCDEPTR);
BEGIN
  CPCUR@.JMPP := STACK; (* SET JMPP TO NODE JUST PUSHED *)
  STACK@.ADDR := 0; (* MARK THE TERMINAL NODE OF JUMP *)
  GENKEY (-2); (* GEN SPACE FOR ABSOLUTE ADDRESS *)
  GENKEY (-2)
END;
(* SETFWDJMP *)

```

```

{*****}

```

```
(*-46-----*)
{* SETBAKJUMP IS SIMILAR TO SETFWDJUMP EXCEPT IN THIS CASE *}
{* THE FIRST NODE OF A POTENTIAL ADDRESS SPACE PAIR IS *}
{* PUSHED ONTO THE DESIGNATED LOOP/IF STACK AFTER ITS *}
{* JMPF HAS BEEN SET TO A CCDE DATA NODE INSERTED AS AN *}
{* ANCHOR FOR THIS BACK JUMP; THE POTENTIAL ADDRESS SPACE *}
{* NODE WILL LATER BE POPPED AND INSERTED (ALONG WITH ITS *}
{* THE INSERTION OF ANOTHER NODE TO COMPOSE AN ADDR PAIR) *}
(*-----*)
```

```
PROCEDURE SETBAKJMP (VAR STACK : CODEPTR);
VAR JCODE : CODEPTR;
BEGIN
  JCCDE := NEWCODE (-2);
  JCODE@.JMFF := CPCUR;
  CPCUR@.ADDR := 0;
  GENKEY (K_NOP);
  PUSHCODE (JCODE, STACK)
END; (* SETBAKJUMP *)
```

```
(*****)
```

```
(*-47-----*)
{* FLOCP GENERATES CCDE FOR THE LOOP COMMAND; *}
{* IT SETS UP THE START OF A LOOP CONSTRUCT BY GENERATING *}
{* AN ANCHOR NODE FOR THE BACK JUMP IN THE LOOP. *}
(*-----*)
```

```
PROCEDURE FLOCP;
BEGIN
  SETBAKJMP (LOOPSTACK);
  PUSHCODE (NEWCODE (K_NOP), ENDLOOPSTACK);
  SCAN (TCKNUM);
  CLOSELINE
END; (* FLOCP *)
```

```
(*****)
```

```
(*-48-----*)
{* PWHILE GENERATES CCDE FOR THE WHILE COMMAND; *}
{* IT IS SIMILAR TO FLOCP EXCEPT IT INSERTS CODE TO *}
{* EVALUATE A BOOLEAN EXPRESSION (CONDITION). *}
(*-----*)
```

```
PROCEDURE PWHILE;
BEGIN
  SETBAKJMP (LOOPSTACK);
  PCCNDITCN;
  PUSHCODE (NEWCODE (K_NOP), ENDLOOPSTACK);
  SETFWDJMP (ENDLOOPSTACK)
END; (* PWHILE *)
```

```
(*****)
```

```
(* -49 ----- *)
{* PENDLOCF POPS AND INSERTS CODE WHICH HAD BEEN STACKED *}
{* EARLIER AS A RESULT OF THE START OF A LOOP CONSTRUCT. *}
{* ----- *)
```

```
PROCEDURE PENDLOCF;
```

```
VAR JCODE : CODEPTR;
```

```
BEGIN
```

```
IF TCKNUM = ENDLOOPFLOK THEN
  GENKEY (K_GTO);
  JCODE := FOPCODE (LOPSTACK);
  PUTKEY (JCODE);
  GENKEY (-2);
  JCODE := POPCODE (ENDLOOPSTACK);
  PUTKEY (JCODE);
  CLCSELINE
```

```
END; (* PENDLOCF *)
```

```
(*****)
```

```
(* -50 ----- *)
{* PUNTIL GENERATES CODE TO EVALUATE A BOOLEAN EXPRESSION *}
{* AND CALLS PENDLOCF TO CLOSE OUT THE LOOP. *}
{* ----- *)
```

```
PROCEDURE PUNTIL;
```

```
BEGIN
```

```
PCCNDITION;
PENDLOCF
```

```
END; (* PUNTIL *)
```

```
(*****)
```

```
(* -51 ----- *)
{* PNEXT GENERATES CODE FOR THE NEXT COMMAND. *}
{* THIS ROUTINE IS WEAK IN SYNTAX ERROR CHECKING. *}
{* ----- *)
```

```
PROCEDURE PNEXT;
```

```
VAR ISICT : SLOTPTR;
    JCODE : CODEPTR;
```

```
BEGIN
```

```
SCAN (TCKNUM);
ISICT := IDLOOKUP (ACCUM, ACCINX);
IF ISICT@.AUXREG2 = -1 THEN
  GENKEY (K_1)
ELSE
  BEGIN
    GENKEY (K_RCL);
    GENKEY (ISLOT@.AUXREG2)
  END;
  GENKEY (K_SUM);
  GENKEY (ISLOT@.REGNC);
  GENKEY (K_GTO);
  JCODE := FOPCODE (FORSTACK);
  PUTKEY (JCODE);
  GENKEY (-2);
  JCODE := POPCODE (NEXTSTACK);
  PUTKEY (JCODE);
  CLCSELINE
```

```
END; (* PNEXT *)
```

```
(*****)
```

```

{* -52-----*
* PFOR GENERATES CODE FOR THE PFCR COMMAND. THIS ROUTINE *
* (AND THE PNEXT ROUTINE) IS WEAK IN SYNTAX ERROR CHECK- *
* ING. THERE ARE MANY PLACES WHERE SIMPLZ CHECKS FOR *
* CORRECT SYNTAX COULD HAVE BEEN PERFORMED BUT WERE NOT *
* BECAUSE CF COMPLEXITY SUCH CHECKS WOULD HAVE INDUCED. *
*-----*

```

```
PROCEDURE PFCR;
```

```
VAR ISICT : SLOTPTR;
```

```

BEGIN
  SCAN (TCKNUM);
  ISICT := IDLOOKUP (ACCUM, ACCINX);
  ISICT@.AUXREG1 := NEWREG;
  SCAN (TCKNUM);
  SCAN (TCKNUM);
  PEXER;
  GENKEY (K STO);
  GENKEY (ISLOT@.REGNC);
  SCAN (TCKNUM);
  PEXER;
  GENKEY (K STO);
  GENKEY (ISLOT@.AUXREG1);
  IF TCKNUM = STEPTOK THEN
    BEGIN
      SCAN (TOKNUM);
      PEXER;
      GENKEY (K STO);
      ISICT@.AUXREG2 := NEWREG;
      GENKEY (ISLOT@.AUXREG2)
    END;
  PUSHCODE (NEWCODE (-2), PFCRSTACK);
  FORSTACK@.JMPP := CPCUR;
  GENKEY (K RCL);
  GENKEY (ISLOT@.REGNO);
  GENKEY (K X T);
  GENKEY (K RCL);
  GENKEY (ISLOT@.AUXREG1);
  GENKEY (K INV);
  GENKEY (K IPXGET);
  PUSHCODE [NEWCODE (K NOP), NEXTSTACK];
  CPCUR@.JMPP := NEXISTACK;
  GENKEY (-2);
  GENKEY (-2);
  CLCSELINE
END;

```

```
(* PFOR *)
```

```

{*****}
{*
{*          IF-BRANCHING ROUTINES
{*
{*****}

```

```

{*53-----}
{* QUITERROR IS CALLED WHENEVER A QUIT STATEMENT IS
{* ENCOUNTERED WHILE NOT WITHIN THE SCOPE OF A LOOP.
{*-----}

```

PROCEDURE QUITERROR;

```

BEGIN
  WRITE (LISTF, '***** ATTEMPT TO "QUIT" WHILE NOT ');
  WRITELN (LISTF, 'INSIDE A LOOP. ');
  PERFOR
END;
(* QUITERRCR *)

```

```

{*****}

```

```

{*54-----}
{* PQUIT GENERATES POTENTIAL ADDRESS SPACE WHOSE JMPP
{* POINTS TO THE MOST CURRENT CODE NODE ON THE ENDLOOP
{* STACK; THUS, CONTROL WILL LEAVE THE MOST CURRENTLY
{* EXECUTING LOOP DURING TI-59 EXECUTION. NOTE THAT THIS
{* IMPLEMENTATION WILL NOT ALLOW LINE# TO FOLLOW 'QUIT'
{*-----}

```

PROCEDURE PQUIT;

```

BEGIN
  IF ENDLCOFSTACK = NIL THEN
    QUITERRCR
  ELSE
    BEGIN
      GENKEY (K-GTO);
      SETFWLJMP (ENDLCOFSTACK);
      SCAN (TOKNUM);
      IF TOKNUM = NUMEERTOK THEN
        BEGIN
          PSUBERROR;
          WRITE (LISTF, '***** "QUIT" DOES NOT ACCEPT ');
          WRITELN (LISTF, 'LINE NUMBERS THIS IMPLEMENT. ');
        END
      ELSE
        CLOSILINE
    END
END;
(* PQUIT *)

```

```

{*****}

```

```

{*56-*}
PROCEDURE PTHENLINE;
FORWARD;

```

```

{*****}

```



```

*-55-----*
* PTHENELSE DETERMINES WHETHER THE ELSE BRANCH OF AN *
* IF-THEN-ELSE IS LINE-ORIENTED (LINE#) OR LOOP-ORIENTED *
* ('QUIT'); APPROPRIATE ROUTINE IS CALLED TO SET JUMP. *
*-----*

```

PROCEDURE PTHENELSE;

```

BEGIN
  SCAN (TCKNUM);
  IF TCKNUM = NUMBERTOK THEN
    BEGIN
      GENKEY (K_GTO);
      PTHENLINE
    END
  ELSE IF TCKNUM = QUITTOK THEN
    PQUIT
  ELSE
    BEGIN
      WRITE (LISTP, '***** "IF-THEN-ELSE" LIMITED TO ');
      WRITELN (LISTP, '"QUIT" OR LINE NUMBERS. ');
      PERROR
    END
END;
(* PTHENELSE *)

```

(*****)

```

*-56-----*
* PTHENLINE SETS THE LINE-ORIENTED JUMPS FOR THE IF-THEN *
* OR IF-THEN-ELSE STATEMENTS. *
*-----*

```

PROCEDURE PTHENLINE;

```

BEGIN
  SETJMPFEXT (XNUMBER (ACCUM, ACCINX));
  GENKEY {-2};
  GENKEY {-2};
  SCAN (TCKNUM);
  IF TOKNUM = ELSETOK THEN
    PTHENELSE
  ELSE
    CIOSELINE
END;
(* PTHENLINE *)

```

(*****)

```

*-57-----*
* PTHENQUIT SETS THE LOOP-ORIENTED JUMPS FOR IF-THEN OR *
* IF-THEN-ELSE STATEMENTS. *
*-----*

```

PROCEDURE PTHENQUIT;

```

BEGIN
  IF ENDICOPSTACK = NIL THEN
    CUIFEROR
  ELSE
    BEGIN
      SETFWLJMP (ENDICOPSTACK);
      SCAN (TCKNUM);
      IF TOKNUM = ELSETOK THEN
        PTHENELSE
      ELSE
        CIOSELINE
    END
END;
(* PTHENQUIT *)

```

(*****)

```

{*--58-----*}
{* PIF DETERMINES THE TYPE OF 'IF' STATEMENT; IT WILL *}
{* CALL THE REQUIRED SET ROUTINES FOR UNSTRUCTURED JUMPS *}
{* (LINE OR LOOP ORIENTED) OR PERFORM THE SET UP ITSELF *}
{* FOR STRUCTURED JUMPS. *}
{*-----*}

```

PROCEDURE PIF;

```

BEGIN
  PCONDITION;
  IF TOKNUM = THENTOK THEN
    BEGIN
      SCAN (TOKNUM);
      IF TOKNUM = NUMBERTOK THEN
        ETHENLINE
      ELSE IF TOKNUM = QUITTOK THEN
        ETHENQUIT
      ELSE
        PERROR
    END
  ELSE
    BEGIN
      PUSHCCDE (NEWCCDE (K_NOP), ENDIFSTACK);
      CPCUR@.BAKP := ENDIFSTACK;
      ENDIFSTACK@.BAKP := CFCUR;
      PUSHCCDE (NEWCCDE (K_NOP), IPSTACK);
      SETFWDJMP (IPSTACK);
      CFCSELINE
    END
  END;

```

(* PIF *)

(*****)

```

{*--59-----*}
{* ELSE_ADJUST PERFORMS HOUSE-KEEPING ON THE VARIOUS 'IF' *}
{* STACKS DEPENDENT UPON THE FORM OF THE STRUCTURED 'IF' *}
{* STATEMENT ENCOUNTERED; IF-ENDIF REQUIRES A DIFFERENT *}
{* SEQUENCE OF PUSH/POP STACK THAN DOES IF-ELSE-ENDIF OR *}
{* IF-ELSEIF-ELSE-ENLIF. *}
{*-----*}

```

PROCEDURE ELSE_ADJUST;

```

BEGIN
  WITH ENDIFSTACK@ DC
  BEGIN
    IF BAKP <> NIL THEN
      BEGIN
        BAKP@.BAKP := NIL;
        BAKP@.JMPP@.BAKP := NIL;
        BAKP := NIL
      END
    END
  END;

```

(* ELSE_ADJUST *)

(*****)

```
(*-60-----*)
{* PELSEIF PERFORMS A SEQUENCE OF STACK MANIPULATIONS IN *}
{* ORDER TO GENERATE THE ADDRESS SPACES AND JUMPS WHICH *}
{* IMPLEMENT THE ELSEIF CONSTRUCT. *}
(*-----*)
```

```
PROCEDURE PELSEIF;
```

```
VAR JCODE : CODEPTR;
```

```
BEGIN
  ELSE ADJUST;
  GENKEY (K_GTO);
  SETFWDJMP (ENDIFSTACK);
  JCODE := POPCODE (IFSTACK);
  PUTKEY (JCODE);
  PCCNDITION;
  PUSHCODE (NEWCODE (K_NOP), IFSTACK);
  SETFWDJMP (IFSTACK);
  CLOSELINE
```

```
END; (* PELSEIF *)
```

```
(*****)
```

```
(*-61-----*)
{* PELSE IS SIMILAR TO ELSEIF EXCEPT IT DOES NOT PARSE/ *}
{* GENERATE CODE TO EVALUATE A BOOLEAN EXPRESSION. *}
(*-----*)
```

```
PROCEDURE PELSE;
```

```
VAR JCODE : CODEPTR;
```

```
BEGIN
  ELSE ADJUST;
  GENKEY (K_GTO);
  SETFWDJMP (ENDIFSTACK);
  JCODE := POPCODE (IFSTACK);
  PUTKEY (JCODE);
  SCAN (TKNUM);
  CLOSELINE
```

```
END; (* PELSE *)
```

```
(*****)
```

```

(*-62-----*)
(* PENDIF CLOSES UP THE SCOPE OF A STRUCTURED 'IF' *)
(* CCNSTRUCT BY POPPING THE APPROPRIATE STACKS AND *)
(* INSERTING AND DISCARDING CODE WHICH HAD BEEN STACKED; *)
(* DISCARDING/INSERTING IS DEPENDENT UPON THE PARTICULAR *)
(* TYPE OF 'IF' CONSTRUCT (IF-ENDIF OR IF-ELSE-ENDIF). *)
-----*)

PROCEDURE PENDIF;
VAR  DUMPC, JCODE : CODEPTR;
BEGIN
  WITH ENCIFSTACK@ DO
    BEGIN
      IF EAKP <> NIL THEN
        BEGIN
          (* NC ELSE/ELSEIF HAS BEEN SEEN *)
          EAKP@.EAKP := NIL;          (* NULLIFY POINTERS *)
          BAKP := NIL;
          DUMPC := PCFCODE (ENDIFSTACK);  (* CLEAR STACK *)
          DUMPC@.ADDR := -1;
          JCODE := PCFCODE (IFSTACK);      (* INSERT ENDIF *)
          PUTKEY (JCODE)
        END
      ELSE
        (* ELSE/ELSEIF HAS BEEN SEEN *)
        BEGIN
          JCODE := PCFCODE (ENDIFSTACK);  (* INSERT ENDIF *)
          JCODE@.ADDR := 0; (* MARK TERMINAL NODE OF JUMP *)
          PUTKEY (JCODE)
          (* ELSE ADJUST HAS ALREADY *)
          (* CLEARED ENCIFSTACK. *)
        END
      END;
    SCAN (TCKNUM);
    CLCSELINE
  END;
END;
(* PENDIF *)

```

```

{*****}
{*
{*          I/C COMMAND ROUTINES
{*
{*****}

```

```

(*-84-*)
PROCEDURE WRITLN (VAR WFILE, MSGFILE : TEXT;
                 MSG_NO : INTEGER);          FORWARD;

```

```

{*****}

```

```

{*-----*}
{* PDATA GENERATES SINGLE NCP TO PROTECT FROM LINE- *}
{* ORIENTED JUMP REPERENCES; THIS COMMAND IS INTENDED FOR *}
{* USE AT THE START OF A PROGRAM SINCE ITS USE WITHIN *}
{* LCOPS, SER'S, OR FUNCTIONS WOULD RENDER THE DATA TO *}
{* READ MAPPING MEANINGLESS. THIS ROUTINE READS THE DATA *}
{* VALUES FOUND AS ITS PARAMETERS, COUNTS THEM, AND *}
{* STORES THEM IN AN ARRAY OF RECORDS WHICH IS ACCESSED *}
{* BY SUBSEQUENT READ COMMANDS. *}
{*-----*}

```

```

PROCEDURE FLATA;

```

```

VAR DATASIGN : CHAR;

```

```

BEGIN
  GENKEY (K NOP);
  SCAN (TCKNUM);
  WHILE (TOKNUM IN SIGNTOKS) OR (TOKNUM = NUMBERTOK) DO
    BEGIN
      DATASIGN := BLANK;
      IF TOKNUM IN SIGNTOKS THEN
        BEGIN
          IF TOKNUM = MINUSTOK THEN
            DATASIGN := '-';
          SCAN (TOKNUM)
        END;
      IF DATAIX = REGEASE + 1 THEN
        BEGIN
          PWARN;
          WRITE (LISTF, '***** EXCEEDED DATASTORE ');
          WRITE (LISTF, 'CAP = ', REGBASE : 1);
          WRITELN (LISTF, '...RESET DATA INDEX TO 1. ');
          DATAIX := 1
        END;
      CATALIST (.DATAIX.) .NUMB := ACCUM;
      CATALIST (.DATAIX.) .SIGN := DATASIGN;
      DATAIX := DATAIX + 1;
      SCAN (TOKNUM);
      IF TOKNUM = COMMATOK THEN
        SCAN (TOKNUM)
    END;
  CLCSFLINE
END;

```

```

(* PDATA *)

```

```

{*****}

```

```

*-64-----*
* PREAD ONLY GENERATES A NCP INSTRUCTION TO ALLOW FOR *
* LINE-ORIENTED JUMP REFERENCES; OTHERWISE, THIS CCMAND *
* WRITES THE READP FILE WHICH INDICATES DATA VALUES FOR *
* RESPECTIVE REGISTERS AND THEIR WBASIC VARIABLE NAMES; *
* THE READP FILE IS USED TO INPUT DATA PRIOR TO PROGRAM *
* EXECUTION ON THE II-59, THUS, SAVING PROGRAM STEPS. *
* THE CCNSTRUCT IS INTENDED FOR USE AT THE START OF A *
* PROGRAM. IF NESTED WITHIN THE PROGRAM, LOOPS, SBR'S, *
* AND FN'S WOULD RENDER THE DATA/READ MAP MEANINGLESS. *
*-----*

```

```
PROCEDURE PREAD;
```

```
VAR IDSICT : SLCTPTR;
```

```
BEGIN
```

```
  GENKEY (K NOP);
```

```
  IF FIRSTREAD THEN
```

```
    BEGIN
```

```
      REWRITE (READP, 'NAME=READP.WBASIC.A');
```

```
      WRITELN (READP, MSGF, 9);
```

```
      FIRSTREAD := FALSE
```

```
    END;
```

```
  IF NOT INDEXERROR THEN
```

```
    BEGIN
```

```
      SCAN (TOKNUM);
```

```
      WRITELN (READP);
```

```
      WHILE TOKNUM = IDENTOK DO
```

```
        BEGIN
```

```
          ICSLOT := IDLOOKUP (ACCUM, ACCINX);
```

```
          WITH IDSLOT DO
```

```
            BEGIN
```

```
              IF READIX >= DATAIX THEN
```

```
                BEGIN
```

```
                  PWAEN;
```

```
                  WRITE (LISTP, '***** READ PAST DATA');
```

```
                  WRITE (LISTP, 'INDEX...IGNORING');
```

```
                  WRITELN (LISTP, 'SUBSEQUENT READ/DATA.');
```

```
                  WRITELN (READP);
```

```
                  WRITE (READP, '***** READ PAST DATA');
```

```
                  WRITE (READP, 'INDEX...SUBSEQUENT');
```

```
                  WRITELN (READP, 'READ/DATA IGNORED.');
```

```
                  INDEXERR := TRUE;
```

```
                  PRECOVER
```

```
                END
```

```
              ELSE
```

```
                BEGIN
```

```
                  WRITE (READP, ' ':5);
```

```
                  WRITE (READP, DATALIST (.READIX.) .SIGN);
```

```
                  WRITE (READP, DATALIST (.READIX.) .NUMB);
```

```
                  WRITE (READP, ' ':2);
```

```
                  ZEECPAD (READP, REGNO, 2);
```

```
                  WRITELN (READP, ' ':3, IDENT);
```

```
                  READIX := READIX + 1;
```

```
                  SCAN (TOKNUM);
```

```
                  IF TOKNUM = COMMATOK THEN
```

```
                    SCAN (TCKNUM)
```

```
                END
```

```
            END
```

```
          END
```

```
        PRECOVER;
```

```
      CLOSELINE
```

```
    END;
```

```
(* PREAD *)
```

```
(*****)
```

```

(*-65-----*)
{* PRESTORE GENERATES A SINGLE NOP TO PROTECT FROM LINE- *}
{* ORIENTED JUMP REPERENCES. IN THIS IMPLEMENTATION *}
{* THIS CCNSTRUCT IS NOT OF GREAT VALUE SINCE DATA/READ *}
{* STATEMENTS ARE SUGGESTED FOR USE AT THE START OF A *}
{* PROGRAM CNLY; THIS ROUTINE RESETS THE READ INDEX TO *}
{* ITS INITIAL VALUE (=1). *}
(*-----*)

```

PROCEDURE PRESTORE;

```

BEGIN
  GENKEY (K_NOP);
  READIX := -1;
  SCAN (TCKNUM);
  CLOSELINE
END; (* PRESTORE *)

```

(*****)

```

(*-66-----*)
{* PINPUT PARSES A LIMITED FORM OF THE WBASIC "INPUT" *}
{* STATEMENT; THE LIST OF INPUT PARAMETERS MAY CONSIST OF *}
{* VARIAELE NAMES ONLY. *}
(*-----*)

```

PROCEDURE PINPUT;

```

VAR TENDIG : -1..9;
    INPVAR : SLCTPTB;

```

```

BEGIN
  SCAN (TOKNUM);
  GENKEY (K_CE);
  TENDIG := -1; (* FLAG CHECKS IF INPUT VARS ARE LISTED *)
  WHILE TCKNUM = IDENTCK DO
    BEGIN
      INPVAR := IDLOCKUP (ACCUM, ACCINX);
      TENDIG := INPVAR@.REGNO DIV 10;
      GENKEY (TENDIG); (* REG IN WHICH INPUT TO BE STORED *)
      GENKEY (INPVAR@.REGNO - (TENDIG * 10));
      GENKEY (K_INT); (* CLOSES DISPLAY REG *)
      GENKEY (K_RS);
      GENKEY (K_STO);
      GENKEY (INPVAR@.REGNO);
      SCAN (TOKNUM);
      IF TOKNUM = COMMATOK THEN (* PARAMETERS SEPARATED *)
        SCAN (TOKNUM) (* BY COMMAS OR BLANKS. *)
      END;
      IF TENDIG = -1 THEN (* GENERATES A R/S IF "INPUT" *)
        GENKEY (K_RS); (* IS USED WITHOUT A VAR LIST *)
      CLOSELINE
    END;
  END; (* PINPUT *)

```

(*****)

```

{* -67-----*}
{* PPRINT PARSES A LIMITED FORM OF THE WBASIC "PRINT" *}
{* STATEMENT; IT ALLCWS EXPRESSIONS, VARIABLE NAMES, AND *}
{* LITERAL NUMERICS IN THE LIST OF PARAMETERS. *}
{*-----*}

```

```
PROCEDURE PPRINT;
```

```

BEGIN
  SCAN (TCKNUM);
  WHILE TCKNUM IN BEGIN_EXPR TOKS DO
    BEGIN
      PEXPR;
      IF FC100 THEN (* WITH FC100 *)
        GENKEY (K_PRT)
      ELSE (* WITHOUT FC100 *)
        BEGIN
          GENKEY (K_PAUSE);
          GENKEY (K_RS)
        END;
      IF TOKNUM = COMMATOK THEN (* CAN SEPARATE ITEMS BY *)
        SCAN (TOKNUM) (* COMMAS OR BLANKS. *)
      END;
      IF FC100 THEN (* WITH FC100 *)
        GENKEY (K_ADV);
      CLOSELINE
    END;
  END; (* PPRINT *)

```



```

*****
*
*   OTHER COMMAND WORD PARSE/GENERATE ROUTINES
*
*****

```

```

*-68-----*
* PCPTICN SETS/RESETS THE OPTICN TOGGLES WHICH ALLOW THE *
* USER OUTSIDE CONTROL OF COMPILER OUTPUTS; NOTE THAT *
* THIS CCNSTRUCT IS NOT THE SAME AS THE OPTION STATEMENT *
* OF WBASIC; IT IS INTENDED FOR USE AT THE BEGINNING OF *
* THE SCURCE PROGRAM; USE ANYWHERE ELSE MAY PRODUCE *
* UNEXPECTED RESULTS AND/OR OPERATING SYSTEM ERRORS. *
*-----*

```

PROCEDURE PCPTION;

```

VAR SWITCH : BOOLEAN;
    TOGGLE : INTEGER;

```

```

BEGIN
  GENKEY (K NOP);
  SCAN (TOKNUM);
  WHILE NOT (TOKNUM IN TRAILTOKS) DO
    BEGIN
      SWITCH := TRUE;
      IF TOKNUM IN SIGNTOKS THEN
        BEGIN
          IF TOKNUM = MINUSTOK THEN
            SWITCH := FALSE;
          SCAN (TOKNUM)
        END;
      IF TOKNUM = NUMBERTOK THEN
        BEGIN
          TOGGLE := XNUMBER (ACCUM, ACCINX);
          IF TOGGLE IN (.0..8.) THEN
            CASE TOGGLE OF
              0 : LINK59 := TRUE;
              1 : PC100 := SWITCH;
              2 : OPTIFAR := SWITCH;
              3 : OPTINOP := SWITCH;
              4 : CODUMP := SWITCH;
              5 : SYLUMP := SWITCH;
              6 : DSDUMP := SWITCH;
              7 : TOKCUT := SWITCH;
              8 : TOKLIS := SWITCH;
            END (* CASE *)
          ELSE
            BEGIN
              PWARN;
              WRITE (LISTP, '***** NO SUCH OPTION...');
              WRITELN (LISTP, ACCUM)
            END
          END
        ELSE
          BEGIN
            PWARN;
            WRITE (LISTP, '***** OPTION PARAMETERS ARE ');
            WRITELN (LISTP, '-8..0..+8 ONLY. ');
          END;
        SCAN (TOKNUM)
      END;
    CLCSELINE
  END;
  (* POPTICN *)
  (*****

```

```
(*--69-----*
{* PNOLET PARSES AND GENERATES CODE FOR AN ASSIGNMENT *
{* STATEMENT WHICH DOES NOT BEGIN WITH THE 'LET' COMMAND. *}
*-----*)
```

```
PROCEDURE PNOLET;
```

```
VAR RESULT : SLOTPTR;
```

```
BEGIN
  RESULT := IDLOOKUP (ACCUM, ACCINX);
  SCAN (TCKNUM);
  IF TCKNUM <> EQUALTOK THEN (* '=' *)
  ELSE
  BEGIN
    SCAN (TCKNUM);
    PEXPR;
    GENKEY (K STO);
    IF RESULT@.TYP = VARID THEN
      GENKEY (RESULT@.REGNO)
    ELSE IF RESULT@.TYP = FNPID THEN
      GENKEY (RESULT@.FNREGNO)
    ELSE
      PERROR
  END
END; (* PNOLET *)
```

```
(*****)
```

```
(*--70-----*
{* PLET PARSES AND GENERATES CODE FOR A 'LET' STATEMENT. *}
*-----*)
```

```
PROCEDURE PLET;
```

```
BEGIN
  SCAN (TCKNUM);
  PNCLET
END; (* PLET *)
```

```
(*****)
```

```
(*--71-----*
{* PREM GUARDS AGAINST USE OF A GOTO DIRECTED TO A REM BY *
{* CAUSING GENERATION OF A NODE (LOADED W/ A NOP INSTRU) *
{* WHICH CAN BE REFERENCED BY A JMPP POINTER; THE SCANNER *
{* HAS RESPONSIBILITY FOR SKIPPING OVER THE CMT TEXT. *}
*-----*)
```

```
PROCEDURE PREM;
```

```
BEGIN
  GENKEY (K NOP);
  SCAN (TCKNUM)
END; (* PREM *)
```

```
(*****)
```

```

(*-72-----*)
(* PGCTO GENERATES THE TI-59 GTO STATEMENT AND ITS      *)
(* POTENTIAL ADDRESS SPACE; THE JUMP POINTER FROM THE 1ST *)
(* NODE OF THIS ADDRESS SPACE IS POINTED TO THAT NODE IN *)
(* THE CODE DATA STRUCTURE WHICH IS THE START (OR, IN THE *)
(* CASE OF FORWARD JUMPS, THE POTENTIAL START) OF CODE *)
(* GENERATED FOR THE WBASIC LINE NUMBER REFERENCED IN THE *)
(* GOTO COMMAND.                                         *)
-----*)

```

PROCEDURE PGOTO;

```

BEGIN
  GENKEY (K GTO);
  SCAN (TCKNUM);
  IF TOKNUM <> NUMBERTOK THEN
    PERRCR
  ELSE
    BEGIN
      SETJMFEXT (XNUMBER (ACCUM, ACCINX));
      GENKEY (-2);
      GENKEY (-2);
    END;
  CLCSELINE
END; (* FGOTO *)

```

(*****)

```

(*-73-----*)
(* PGCSUB GENERATES A CALL TO A SUBROUTINE REFERENCED BY *)
(* WBASIC LINE NUMBER; NOTE THAT ALTHOUGH WBASIC CALLS *)
(* SUBROUTINES BY LINE NUMBER, THE TI-59 CODE GENERATED *)
(* CALLS A SUBROUTINE BY A LABEL NAME; AN EXTERNAL JUMP *)
(* IS SET (AS IN THE GOTO), HOWEVER, RESOLUTION OF THE *)
(* JUMP WILL BE MADE BY INSERTING THE LABEL USED IN THE *)
(* CALL IN FRONT OF THE NODE REFERENCED BY THE JMPP; THIS *)
(* INSERTION IS DONE AFTER ALL CODE HAS BEEN GENERATED. *)
(* NOTE THAT THIS ROUTINE NEITHER CHECKS FOR NOR DOES IT *)
(* KNOW OF THE EXISTENCE OF A RETURN STATEMENT IN THE *)
(* SEQUENCE OF SOURCE CODE ASSUMED TO BE THE GOSUB BODY; *)
(* IF THE USER DOES NOT PROVIDE A RETURN STATEMENT, THEN *)
(* NO CORRESPONDING TI-59 INVSBR (SBR RETURN) WILL BE *)
(* GENERATED, AND THE SBR RETURN REGISTER IN THE *)
(* CALCULATOR WILL NEVER BE CLEARED OF THAT SBR CALL. *)
-----*)

```

PROCEDURE PGOSUB;

```

BEGIN
  GENKEY (K SBR);
  SCAN (TCKNUM);
  IF TCKNUM <> NUMBERTOK THEN
    PERRCR
  ELSE
    BEGIN
      SETJMFEXT (XNUMBER (ACCUM, ACCINX));
      GENKEY (NEWLBL)
    END;
  CLCSELINE
END; (* PGOSUB *)

```

(*****)

```

*-74-----*)
* PRETURN GENERATES THE RETURN FROM A SUBROUTINE. *)
* STRUCTURED PROGRAMMING DISCIPLINE DEMANDS A RETURN FOR *)
* EACH SUBROUTINE CALL; NOTE THAT THE TI-59 HAS A LIMIT *)
* OF SBR RETURN ADDRESSES WHICH CAN BE STACKED; THE USER *)
* SHOULD REMEMBER THAT THE WBASIC RETURN STATEMENT IS *)
* THE ONLY ONE WHICH WILL GENERATE THE TI-59 INVSBR *)
* FOR A GCSUB GENERATED SBR CALL (FUNCTIONS GENERATE SBR *)
* AND INVSBR ALSO, BUT THEY DO THIS AS A RESULT OF THE *)
* FNEED STATEMENT OF A ONE-LINE FUNCTION). *)
-----*)

```

PROCEDURE PRETURN;

```

BEGIN
  GENKEY (K_INV_SBR);
  CLOSILINE
END;
(* PRETURN *)

```

(*****)

```

*-75-----*)
* PPAUSE GENERATES (82) (31) WHICH ARE ACTUALLY A VOID *)
* CODE AND THE 'LRN' KEY; WHEN ENTERING HIS PROGRAM INTO *)
* THE CALCULATOR THE USER MUST ENTER 'STO 31' INSTEAD OF *)
* (82) (31) WHICH CANNOT BE ENTERED DIRECTLY ANYWAY; THEN *)
* THE USER MUST BACKSTEP AND CHANGE THE ORIGINAL 'STO 31' *)
* BY ISSUING THE FOLLOWING EDITING KEY STROKE SEQUENCE *)
* TO THE CALCULATOR IMMEDIATELY AFTER ENTERING THE *)
* 'STO 31': BST,BST,NOP,SST. THIS WILL REVISE THE *)
* ORIGINAL 'STO 31' TO 'NOP 31'; WHEN ENCOUNTERED BY THE *)
* CALCULATOR THESE 2 INSTRUCTIONS WILL STOP EXECUTION BY *)
* SHIFTING THE CALCULATOR INTO THE LEARN (LRN) MODE; *)
* IN ORDER TO RESUME EXECUTION, THE USER MUST ENTER *)
* 'LRN' (PLACING THE DISPLAY REG BACK INTO VIEW) *)
* FOLLOWED BY 'R/S' (WHICH RESUMES THE PROCESSING MODE); *)
* THIS INTERRUPTION OF EXECUTION DOES NOT CAUSE ANY SIDE *)
* EFFECTS AND PROVIDES AN ACCURATE INDICATION OF THE *)
* LOCATION OF ANY 'PAUSE' STATEMENTS PLACED IN THE *)
* WBASIC SOURCE CODE; THIS IMPLEMENTATION OF THE 'PAUSE' *)
* INSTRUCTION PROVIDES A CONVENIENT AND RECOGNIZABLE *)
* DEBUGGING/TRANSLATION TOOL WHICH CARRIES A LOW *)
* OVERHEAD IN TERMS OF REGISTER/PROGRAM STEP USE. *)
-----*)

```

PROCEDURE PPAUSE;

```

BEGIN
  GENKEY (82);
  GENKEY (31);
  CLOSILINE
END;
(* PPAUSE *)

```

(*****)

```

(*-76-----*)
{* FSTOP GENERATES CCDE WHICH CAUSES THE TI-59 TO HALT *}
{* EXECUTION AND DISPLAY '888' THUS SIGNALING THAT A *}
{* PROGRAM STOP HAS BEEN ENCOUNTERED INSTEAD OF A DATA *}
{* INPUT OR MAGNETIC CARD LINKING INSTRUCTION. *}
(*-----*)

```

```
PROCEDURE FSTOP;
```

```
VAR I : 1..4;
```

```
BEGIN
```

```
  GENKEY (K_CE);
  FOR I := 1 TO 3 DO
    GENKEY (8);
  GENKEY (K_RS);
  CLOSELINE
```

```
END; (* FSTOP *)
```

```
(*****)
```

```

(*-77-----*)
{* PEND ASSUMES THAT THE END OF THE WBASIC SOURCE FILE *}
{* HAS BEEN ENCOUNTERED AND WILL INSERT THE END OF FILE *}
{* CHAR INTO THE TOKEN STREAM, CAUSING IMMEDIATE *}
{* TERMINATION OF THE COMPILATION PROCESS. *}
(*-----*)

```

```
PROCEDURE PEND;
```

```
BEGIN
```

```
  GENKEY (K_NOP);
  TOKNUM := -ENDFILTK
```

```
END; (* PEND *)
```

```

{*****}
{*
{*          CODE RESOLUTION ROUTINES
{*
{*****}

```

```

{*--78-----*}
{* PUTGOSUBLBL USES PROCEDURE INSERTKEY TO ENTER THE LBL *}
{* REFERENCED BY THE GOSUB CALL INTO THE CODE SEQUENCE AT *}
{* LOCATICN POINTED TO BY TO BY THE JMPP (LBLP). *}
{*-----*}

```

```

PROCEDURE PUTGOSUBLBL (LBL : LELRNG; VAR LBLP : CODEPTR);
BEGIN
  WITH IELP@.BAKP@ DO
    BEGIN
      INSERTKEY (LBL, SEQP);
      INSERTKEY (K_LEL, SEQP)
    END
END;
(* PUTGOSUBLBL *)

```

```

{*****}

```

```

{*--79-----*}
{* FINDGCSUBLBL SEARCHES THE CODE DATA STRUCTURE TO FIND *}
{* SBR CALLS FOR WHICH THE JMPP HAS BEEN SET; THESE WILL *}
{* CORRESPOND TO WBASIC GOSUB STATEMENTS; THE JMPP IS *}
{* PCLLCWED AND THE CORRECT LABEL IS INSERTED INTO THE *}
{* CODE SEQUENCE USING THE BAKP AND PROCEDURE PUTGOSUBLBL *}
{*-----*}

```

```

PROCEDURE FINDGOSUBLBL (VAR START : CODEPTR);
VAR TRAVELP, TAILP : CODEPTR;
BEGIN
  TRAVELP := START@.SEQP;
  TAILP := START;
  WHILE TRAVELP <> ENDCP@.SEQP DO
    BEGIN
      WITH TRAVELP@ DC
        IF (JMPP <> NIL) AND (TAILP@.KEY = K_SBR) THEN
          BEGIN (* FIRST CHECK FOR REDUNDANT GOSUB CALL *)
            IF JMPP@.BAKP@.SEQP@.KEY = K_LBL THEN
              KEY := JMPP@.BAKP@.SEQP@.SEQP@.KEY
            ELSE
              PUTGOSUBLBL (KEY, JMPP); (* INSERT A LABEL *)
              JMPP@.ADDR := -1; (* UNMARK JMPP ADDR *)
              JMPP := NIL (* RESET JMPP TO NIL *)
            END;
          TAILP := TRAVELP;
          TRAVELP := TRAVELP@.SEQP
        END
    END;
END;
(* FINDGOSUBLBL *)

```

```

{*****}

```

```

(*-80-----*)
{* OSQPAREN (OPTIMIZE SQUEEZE PARENTHESSES) REMOVES *}
{* UNNECESSARY PARENTHESSES (IN PAIRS) FROM THE CODE DATA *}
{* STRUCTURE FOR THE MOST COMMON CASES, NAMELY '(RCL NN)' *}
{* AND '<LITERAL NUMERIC>' *}
(*-----*)

PROCEDURE OSQPAREN (START : CODEPTR);
VAR CPEN, CLOSE, TAILP, MOVEP : CODEPTR;
    CPENCT, CLOSECT : INTEGER;

(*-----*)

(*-80-01-----*)
{* CCUNTREP COUNTS THE NUMBER OF SEQUENTIAL OCCURENCES OF *}
{* KEYC AT A PARTICULAR LOCATION IN THE CODE DATA STRUCTR *}
{* STRUCTURE; NOTE THAT IT ALSO CHECKS FOR JMP P POINTERS *}
{* TO THESE KEYS. *}
(*-----*)

FUNCTION CCUNTREP (VAR MOVEP:CODEPTR; KEYC:INTEGER):INTEGER;
VAR COUNT : INTEGER;
BEGIN
    COUNT := 0;
    WHILE (MOVEP@.KEY = KEYC) AND (MOVEP@.ADDR = -1) DO
        BEGIN
            MOVEP := MOVEP@.SEQP;
            COUNT := COUNT + 1;
        END;
    CCUNTREP := COUNT
END;
(* COUNTREP *)

(*-----*)

(*-80-02-----*)
{* NUMBERUN MOVES ITS POINTER PARAMETER PASSED ANY NODE *}
{* WHICH CONTAINS A NUMERIC LITERAL KEY CODE AND HAS NO *}
{* POINTER FEPERENCE; IT IS ASSUMED HERE THAT NO JUMP *}
{* POINTER IS EVER SET IN THE MIDDLE OF A NUMERIC LITERAL *}
{* KEY SEQUENCE, ELSE PART OF THE NUMBER MAY BE REMOVED. *}
(*-----*)

PROCEDURE NUMBERUN (VAR MOVEP : CODEPTR);
BEGIN
    WHILE (MOVEP@.KEY IN NUMERICKEY) AND (MOVEP@.ADDR = -1) DO
        MOVEP := MOVEP@.SEQP;
    END;
(* NUMBERUN *)

(*-----*)

```

```

(*-30-03-----*)
{* REMOVEPAREN TAKES PAIRS OF NODES OUT OF THE CODE DATA *}
{* STRUCTURE: NOTE THAT THIS PROCEDURE DOES NOT KNOW WHAT *}
{* CODE IT IS REMOVING; THAT IS DEFINED BY OSQPAREN. *}
-----*

```

```

PROCEDURE REMOVEPAREN (VAR OPEN, CLOSE : CODEPTR;
                      OPENCT, CLOSECT : INTEGER);
BEGIN
  REPEAT
    OPEN@.SEQP := OPEN@.SEQP@.SEQP;
    OPENCT := OPENCT - 1;
    CLOSE@.SEQP := CLOSE@.SEQP@.SEQP;
    CLOSECT := CLOSECT - 1
  UNTIL (OPENCT = 0) OR (CLOSECT = 0)
END;
(* REMOVEPAREN *)

```

```

-----*

```

```

BEGIN (* CSQPAREN MAIN *)
  MOVEP := START;
  WHILE MCVEP@.SEQP <> NIL DO
    BEGIN
      IF (MOVEP@.KEY = K_OPAREN) AND (MOVEP@.ADDR = -1) THEN
        BEGIN
          OPEN := TAILP; (* SET OPEN PTR *)
          OPENCT := CCUNTRP (MOVEP, K_OPAREN);
          IF (MOVEP@.KEY=K_RCL) AND (MOVEP@.ADDR=-1) THEN
            BEGIN
              CLOSE := MOVEP@.SEQP; (* SET CLOSE PTR *)
              MOVEP := MOVEP@.SEQP@.SEQP; (* MOVE AHEAD *)
              CLOSECT := CCUNTRP (MOVEP, K_CPAREN);
              IF CLOSECT > 0 THEN (* IF EXTRAS, DELETE *)
                REMOVEPAREN (OPEN, CLOSE, OPENCT, CLOSECT)
            END
          ELSE IF (MOVEP@.KEY IN NUMERICKEY) AND
                 (MCVEP@.ADDR = -1) THEN
            BEGIN
              WHILE MCVEP@.SEQP@.KEY IN NUMERICKEY DO
                MOVEP := MOVEP@.SEQP; (* PASS OVER NUMBER *)
                CLOSE := MOVEP; (* SET CLOSE PTR *)
                MOVEP := MOVEP@.SEQP; (* MOVE AHEAD *)
                CLOSECT := CCUNTRP (MOVEP, K_CPAREN);
                IF CLOSECT > 0 THEN (* IF EXTRAS, DELETE *)
                  REMOVEPAREN (OPEN, CLOSE, OPENCT, CLOSECT)
            END
          END;
          TAILP := MOVEP;
          MCVEP := MOVEP@.SEQP
        END
      END;
    END;
  END;
(* OSQPAREN *)

```

```

(*****

```



```

(*-81-----*)
(* OSCNCF (OPTIMIZE SCUEEZE NOP) LOCATES ALL 'NOP' KEY *)
(* CODES, RESETS POINTER REFERENCES TO THEM IF THEY EXIST, *)
(* AND THEN PINCHES THEM OUT OF THE CODE DATA STRUCTURE. *)
(*-----*)

PROCEDURE OSCNOP (VAR START : CODEPTR);
VAR  CUR   : CODEPTR;
     I     : 0..3;
     INDEX : CTEXTRNG;

BEGIN
  CUR := START;
  WHILE CUR <> NIL DO
    BEGIN
      IF CUR@.JMPP <> NIL THEN
        (* RESET JMPPS PAST NOPS *)
        (* ASSUMES THAT NO JMP *)
        (* IS SET ON POTENTIAL *)
        (* ADDR SPACE NOPS. *)
        WHILE (CUR@.JMPP@.KEY = K_NOP)
          AND (CUR@.JMPP <> ENDCP) DO
          BEGIN
            CUR@.JMPP@.ADDR := -1;
            CUR@.JMPP@.SEQP := CUR@.JMPP@.SEQP;
            CUR@.JMPP@.ADDR := 0;
          END;
        CUR := CUR@.SEQP;
      END;
    END;
  CUR := START;
  WHILE CUR@.SEQF <> NIL DO
    BEGIN
      INDEX := CUR@.KEY;
      (* FIX THE INDEX TO CTEXT *)
      FOR I := 1 TO (CTEXT(.INDEX).UNIT) DO
        CUR := CUR@.SEQP;
        (* BYPASS REG/ADDR SPACES *)
        IF (CUR@.SEQP@.KEY = K_NOP) AND (CUR@.SEQP@.ADDR = -1)
          THEN CUR@.SEQF := CUR@.SEQP@.SEQP
            (* REMOVE NOP *)
          ELSE
            CUR := CUR@.SEQP
            (* NEXT NODE *)
          END;
      END;
    END;
  (* OSCNOP *)
  (*****

```

```
(*--82-----*)
(* RESOLVE ADDR FILLS THE ADDR FIELDS OF ALL TI-59 CODE *)
(* NODES LINKED IN THE CODE DATA STRUCTURE, AND THEN *)
(* FILLS THE KEY FIELDS OF NODES WHICH HAVE NON-NIL *)
(* JMPF'S WITH THE ABSOLUTE ADDR POINTED TO BY THOSE *)
(* JMPF'S; JMPP'S ARE THEN SET BACK TO NIL. *)
(*-----*)
```

```
PROCEDURE RESOLVE_ADDR (START : CODEPTR);
```

```
VAR TRAVEL : CODEPTR;
    I : INTEGER;
```

```
(*-----*)
(*--82-01-----*)
(* INSERT JMPADDR CONVERTS THE ADDR FOUND AT THE NODE *)
(* REFERENCED BY JMPP PTR INTO A TI-59 MACHINE CODE ADDR *)
(* (2 INTEGERS IN RANGE 0..99), AND INSERTS IT INTO THE *)
(* THE KEY FIELDS (OCCUPIED BY -2'S) OF THE NODES FROM *)
(* WHICH THE JMPP ORIGINATES. *)
(*-----*)
```

```
PROCEDURE INSERT_JMPADDR (JADDR : INTEGER);
```

```
VAR HIPART, LOPART : INTEGER;
```

```
BEGIN
  HIPART := JADDR DIV 100;
  LOPART := JADDR - HIPART * 100;
  TRAVEL@.KEY := HIPART;
  TRAVEL@.SEQP@.KEY := LOPART
END;
(* SPLIT ADDR INTO *)
(* HI/LO PARTS; *)
(* OVERWRITE NOP'S *)
(* W/ ABS ADDR'S. *)
(* INSERT_JMPADDR *)
```

```
(*-----*)
```

```
BEGIN (* RESOLVE_ADDR MAIN *)
  TRAVEL := START;
  I := 0;
  WHILE TRAVEL <> ENDCP@.SEQP DO (* INSERT ABSOLUTE ADDR *)
    BEGIN
      TRAVEL@.ADDR := I;
      TRAVEL := TRAVEL@.SEQP;
      I := I + 1
    END;
  TRAVEL := START;
  WHILE TRAVEL <> ENDCP@.SEQP DO (* FIND/JUSTIFY JMP ADDR *)
    BEGIN
      WITH TRAVEL@ DC
        IF JMPP <> NIL THEN (* FIND JMPP'S WHICH ARE SET *)
          BEGIN
            INSERT_JMPADDR (JMPP@.ADDR);
            JMPP := NIL (* SET JMPP BACK TO NIL *)
          END;
      TRAVEL := TRAVEL@.SEQP
    END
  END;
END; (* RESOLVE_ADDR *)
```

```

(*****
*
*           OUTPUT DUMP ROUTINES
*
*****
)

```

```

(*-83-----*)
* FINDMSG LOCATES THE START OF THE CORRECT MESSAGE IN
* THE MSGFILE.
*-----*)

```

```

PROCEDURE FINDMSG (VAR MSGFILE : TEXT; VAR ESCHAR : CHAR;
                  MSG_NO : INTEGER);
VAR  CH : CHAR;
     I : INTEGER;

```

```

BEGIN
  RESET (MSGFILE, 'NAME=MSGF.FASCAL.A');
  READLN (MSGFILE, ESCHAR);
  REPEAT
    READ (MSGFILE, CH);
    IF CH = ESCHAR THEN (* CHECK FOR ESCAPE CHAR & MSG NO *)
      READLN (MSGFILE, I)
    ELSE
      READLN (MSGFILE)
  UNTIL ((EOF (MSGFILE)) OR
        ((CH = ESCHAR) AND (I = MSG_NO)))
END;                                     (* FINDMSG *)

```

```

(*****

```

```

(*-84-----*)
* WRITLN WRITES A FULL MESSAGE FROM '$N' TO '$N' AS FOUND
* IN THE MSGFILE.
*-----*)

```

```

PROCEDURE WRITLN; (* FWD DECL WITH I/O COMMAND ROUTINES *)
VAR  CH, ESCHAR : CHAR;
     I : INTEGER;

```

```

BEGIN
  FINDMSG (MSGFILE, ESCHAR, MSG_NO);
  REPEAT
    READ (MSGFILE, CH);
    IF CH=ESCHAR THEN (* CHECK FOR EMBEDDED ESCAPE CHARS *)
      READLN (MSGFILE, I) (* AND DISCARD IF FOUND. *)
    ELSE
      WRITE (WFILE, CH);
      IF ECIN (MSGFILE) THEN (* NEXT LINE *)
        BEGIN
          READLN (MSGFILE);
          WRITELN (WFILE)
        END
      END
  UNTIL ((EOF (MSGFILE)) OR
        ((CH = ESCHAR) AND (I = MSG_NO)))
END;                                     (* WRITLN *)

```

```

(*****

```

```
(*-85-----*)
{* WRIT WRITES A ONE-LINE MESSAGE OR THE FIRST LINE OF A *}
{* MESSAGE FROM THE MSGFILE. *}
(*-----*)
```

```
PROCEDURE WRIT (VAR WFILE, MSGFILE : TEXT;
                MSG_NO : INTEGER);
VAR CH, ESCHAR : CHAR;
    I : INTEGER;
```

```
BEGIN
  FINDMSG (MSGFILE, ESCHAR, MSG_NO);
  REPEAT
    READ (MSGFILE, CH);
    WRITE (WFILE, CH);
  UNTIL ECLN (MSGFILE)
END; (* WRIT *)
```

```
(*****)
```

```
(*-86-----*)
{* REPORT COMPUTES AND WRITES THE REGISTER/LABEL SUMMARY. *}
(*-----*)
```

```
PROCEDURE REPORT (VAR WFILE : TEXT);
```

```
VAR LCTAL : LBLRNG;
    RTCTAL : INTEGER;
```

```
BEGIN
  WRITLN (WFILE, MSGF, 3);
  WRITELN (WFILE, ERRCRCT:7, ' FATAL ERRORS. ');
  WRITELN (WFILE, WARNCT:7, ' WARNING MSGS. ');
  IF ERRCRCT > 0 THEN (* CALCULATIONS INCOMPLETE *)
    WRITLN (WFILE, MSGF, 14);
  RTCTAL := NEXTREG - STARTREG;
  LTOTAL := LBLCT - 1;
  WRITELN (WFILE);
  WRITELN (WFILE, NEXTREG:1, ' IS NEXT AVAILABLE REGISTER ');
  WRITELN (WFILE, ' TOTAL REGISTERS RESERVED = ', RESERVECT:1);
  WRITELN (WFILE, ' TOTAL REGISTERS USED = ', RTOTAL:1);
  WRITELN (WFILE, ' TOTAL LABELS USED = ', LTOTAL:1);
  WRITLN (WFILE, MSGF, 4)
END; (* REPORT *)
```

```
(*****)
```

```

(*-87-----*)
(* CODEDUMP WRITES THE TI-59 CODE STORED IN THE CODE DATA *)
(* STRUCTURE AND APPLIES THE CTEXTF FILE TO EACH STEP TO *)
(* PRODUCE THE LITERAL TEXT OF FOR THE KEY STROKES. *)
(*-----*)
PROCEDURE CODEDUMP (VAR WFILE : TEXT; VAR TICODE : CODEPTR);
VAR CUR, HCLD : CODEPTR;
    I : 0..3;

(*-----*)
(*-87-01-----*)
(* WRCODE WRITES THE NUMERICAL FORM OF TI-59 ADDR AND KEY *)
(*-----*)
PROCEDURE WRCODE (VAR CUR : CODEPTR);
BEGIN
    WRITE (WFILE, ' ':5);
    ZEROPAD (WFILE, CUR@.ADDR, 3);
    WRITE (WFILE, ' ':3);
    ZEROPAD (WFILE, CUR@.KEY, 2);
    WRITE (WFILE, ' ':3)
END;
(* WRCODE *)

(*-----*)
BEGIN (* CODEDUMP MAIN *)
    CUR := TICODE;
    WRITLN (WFILE, MSGF, 5);
    WRITE (WFILE, '$');
    (* 'S' MUST BE WRITTEN HERE, *
    (* ELSE WILL INTERFERE W/ WRITLN *)
    WRITLN (WFILE, MSGF, 6);
    WHILE CUR@.SEQ# <> NIL DO
        BEGIN
            WRCCDE (CUR);
            WRITELN (WFILE, CTEXT (.CUR@.KEY.).CODECHAR);
            IF CUR@.KEY IN (.K SBR, K LBL.) THEN
                BEGIN (* MUST NOT TAKE SBR'S OR LBL'S LITERALLY *)
                    CUR := CUR@.SEQ#;
                    WRCCDE (CUR);
                    WRITELN (WFILE, CTEXT (.CUR@.KEY.).CODECHAR)
                END
            ELSE
                BEGIN
                    HCLD := CUR;
                    FOR I := 1 TO (CTEXT (.HOLD@.KEY.).UNIT) DO
                        BEGIN (* UNIT FIELD DEFINES TYPE INSTRUCTION *)
                            CUR := CUR@.SEQ#;
                            WRCCDE (CUR);
                            ZEROPAD (WFILE, CUR@.KEY, 2);
                            WRITELN (WFILE)
                        END
                    END;
                    CUR := CUR@.SEQ#
                END
        END;
END;
(* CODEDUMP *)

(*****

```

```

(*-88-----*)
(* LINK INTERFACE CREATES THE SCRATCH FILE WHICH PROVIDES *)
(* THE LINKER WITH ALL THE INFORMATION IT MUST HAVE TO *)
(* SEGMENT THE TI-59 CODE; ENTRIES IN SCRATCH ARE IN THE *)
(* FORM OF SUB-FILES (MESSAGES) DELIMITED BY "$N". *)
(*-----*)

PROCEDURE LINK_INTERFACE;

(*-88-01-----*)
(* LOGTO IS USED BY LINK INTERFACE TO READ AND WRITE *)
(* FILES TO THE SCRATCH FILE (COPY). *)
(*-----*)

PROCEDURE LOGTO (VAR WFILE, RFILE : TEXT; MSGNO : INTEGER);
VAR CH : CHAR;

BEGIN
  WRITE (WFILE, '$', MSGNO: 1); (* WRITE MSG DELIMITER *)
  WHILE NOT EOF (RFILE) DO
    BEGIN
      WHILE NOT EOLN (RFILE) DO (* COPY THE FILE TO SCRATCH *)
        BEGIN
          READ (RFILE, CH);
          WRITE (WFILE, CH);
        END;
      WRITELN (WFILE);
      IF NOT EOF (RFILE) THEN
        READLN (RFILE);
    END;
  WRITELN (WFILE, '$', MSGNO: 1); (* WRITE MSG DELIMITER *)
  WRITELN (WFILE); (* LOGTO *)
END;

(*-----*)

BEGIN (* LINK_INTERFACE MAIN *)

  REWRITE (SCRATCH, 'NAME=SCRATCH.PASCAL.A');

  WRITELN (SCRATCH, '$1'); (* NEXT REGISTER = MSG $1 *)
  WRITELN (SCRATCH, NEXTREG: 1, ' IS NEXT AVAILABLE REG. ');
  WRITELN (SCRATCH, '$1'); WRITELN (SCRATCH);

  WRITELN (SCRATCH, '$2'); (* TI-59 CODE = MSG $2 *)
  CODEDUMP (SCRATCH, BEGINCF);
  WRITELN (SCRATCH, MSGF, 7); (* END CODE MSG *)
  WRITELN (SCRATCH, '$2'); WRITELN (SCRATCH);

  RESET (NAMEF, 'NAME=NAMEF.WBASIC.A'); (* REG/NAME MAP *)
  LOGTO (SCRATCH, NAMEF, 3); (* = MSG $3 *)

  IF NOT FIRSTREAD THEN (* DATA/READ MAP = MSG $4 *)
    BEGIN
      RESET (READF, 'NAME=READF.WBASIC.A');
      LOGTO (SCRATCH, READF, 4);
    END

END; (* LINK_INTERFACE *)

(*****)

```

```

(*-89-----*)
{* SYMTELDUMP IS A SPECIAL PURPOSE ROUTINE USED FOR          *}
{* DEBUGGING: IT WILL DUMP THE ENTIRE CONTENTS OF THE      *}
{* COMPILER SYMBOL TABLE BUCKET BY BUCKET; THIS ROUTINE  *}
{* IS TOGGLED USING OPTION NUMBER 5.                      *}
(*-----*)

```

```
PROCEDURE SYMTELDUMP (VAR WFILE : TEXT; BUCKET : HASH);
```

```
VAR I LOCK : INTEGER;
    LOCK : SLOTPTR;
```

```

BEGIN
  WRITLN (WFILE, MSGF, 10); (* HEADER MSG *)
  FOR I := 0 TO HASHEASE DO
    IF BUCKET(.I.) <> NIL THEN (* SKIP EMPTY BUCKETS *)
      BEGIN
        ZEROPAD (WFILE, I, 2);
        WRITLN (WFILE, MSGF, 11); (* BUCKET BOUNDARY MSG *)
        LOCK := BUCKET(.I.);
        REPEAT (* UNTIL LOCK = NIL *)
          WRITE (WFILE, ' ':11, LOCK@.IDENT, ' ':1);
          WITH LOCK@ DO
            CASE TYP OF
              VARID : BEGIN
                IF REGNO < 0 THEN (* PI = -314 *)
                  WRITELN (WFILE, '.. CONSTANT');
                ELSE
                  BEGIN
                    ZEROPAD (WFILE, REGNO, 2);
                    WRITELN (WFILE, ' GLOBAL VAR');
                    IF AUXREG1<>-1 THEN (* IF USED *)
                      BEGIN
                        WRITE (WFILE, ' ':32);
                        ZEROPAD (WFILE, AUXREG1, 2);
                        WRITELN (WFILE, ' AUXREG 1');
                      END;
                    IF AUXREG2<>-1 THEN (* IF USED *)
                      BEGIN
                        WRITE (WFILE, ' ':32);
                        ZEROPAD (WFILE, AUXREG2, 2);
                        WRITELN (WFILE, ' AUXREG 2');
                      END;
                  END;
                END;
              FNQID : WRITELN (WFILE, '.. QUICK FN');
              FNLID : BEGIN
                ZEROPAD (WFILE, FNLREG, 2);
                WRITELN (WFILE, ' LONG FN');
              END;
              FNPID : BEGIN
                ZEROPAD (WFILE, FNREGNO, 2);
                WRITELN (WFILE, ' PARAMETER FN');
              END;
            END;
          (* CASE *)
          LOCK := LOCK@.SLOT
        UNTIL LOCK = NIL
      END;
    WRITLN (WFILE, MSGF, 12) (* END SYMTBL MSG *)
  END; (* SYMTELDUMP *)

```

```
(*****)
```

```

(*-90-----*)
(* SEARCH IS A SPECIAL PURPOSE DEBUGGING TOOL; *)
(* THIS PROCEDURE FOLLOWS AND PRINTS THE CONTENTS OF ALL *)
(* POINTERS IN THE CODE DATA STRUCTURE (LINE AND CODEPTR) *)
(* THIS ROUTINE CAN BE TOGGLED USING OPTION NUMBER 6. *)
(*-----*)

PROCEDURE SEARCH (VAR WFILE : TEXT; LSTART : LINEPTR);
VAR LPSEARCH : LINEPTR; CODE : CODEPTR;
BEGIN
  WRITLN (WFILE, MSGF, 13); (* HEADER MSG *)
  LPSEARCH := LSTART;
  REPEAT
    WRITE (WFILE, 'LINUM = ');
    ZEROPAD (WFILE, LPSEARCH@.LINO, 5); (* WBASIC LINE NO *)
    WRITELN (WFILE);
    CCDP := LPSEARCH@.CPTR;
    REPEAT (* TI-59 CODE ATTACHED TO WBASIC LINE NO *)
      WRITE (WFILE, ' ':2);
      ZEROPAD (WFILE, CCDP@.ADDR, 3);
      WRITE (WFILE, ' ':2);
      ZEROPAD (WFILE, CCDP@.KEY, 2);
      WRITELN (WFILE);
      CODE := CCDP@.SEQP
    UNTIL (CCDP = LPSEARCH@.LPTR@.CPTR) OR (CCDP = NIL);
    LPSEARCH := LPSEARCH@.LPTR
  UNTIL LPSEARCH@.LINO = MAXBASLIN (* MAXBASLIN IS END *)
END; (* SEARCH *)

```



```

(*****
*)
*)          INITIALIZATION ROUTINE
*)
(*****

*-91-----*
*) INITIALIZE SETS UP ALL FILES, DATA STRUCTURES, SETS,
*) AND INITIAL VARIABLE VALUES REQUIRED TO BEGIN THE
*) READING AND COMPILATION OF THE WBASIC SOURCE CODE, AND
*) THE OUTPUT OF THE TRANSLATED TI-59 CODE AND LISTINGS.
*)-----*

PROCEDURE INITIALIZE;
VAR I : INTEGER;

(*-----*)

*-91-01-----*
*) LOADRW READS THE RWTBLF FILE (RESERVED WORD TABLE) AND
*) LOADS THE RESERVED WORD CHAR/INDEX ARRAYS; NOTE THAT
*) THE ARRAYS ARE STATIC FIXED AND ARE DEFINED BY THE
*) SYSTEM PARAMETERS RWCHARCT, RWORDCT, RWLENGCT IN THE
*) CONSTANT DECLARATION BLOCK AT THE FRONT OF THE PROGRAM
*)-----*

PROCEDURE LCADRW (VAR RWTBLF : TEXT);
VAR  CHINX, STARTCHINX : 0..RWCHARCT + 1;
     WINX      : 0..RWORDCT + 1;
     LINX, LENG : 0..RWLENGCT + 1;
     CH        : CHAR;

BEGIN
  LINX := 0;          (* INIT LENGTH INDEX *)
  CHINX := 1;        (* INIT CHAR INDEX *)
  WHILE NCT EOF(RWTBLF) DO
    BEGIN
      STARTCHINX := CHINX;
      READ (RWTBLF, WINX); (* READ WORD INDEX (INTEGER) *)
      READ (RWTBLF, CH, CH); (* READ OFF 2 BLANK SPACES *)
      RWCRE (.WINX.) := CHINX;
      REPEAT (* READ CHARS OF ONE WORD INTO CHAR ARRAY *)
        READ (RWTBLF, RWCHAR (.CHINX.));
        CHINX := CHINX + 1;
      UNTIL EOLN (RWTELF);
      REACLN (RWTBLF); (* NEXT WORD *)
      LENG := CHINX - STARTCHINX;
      IF LENG > LINX THEN (* IF LENGTH CHANGE, THEN *)
        BEGIN (* INDEX ITS LOCATION IN *)
          LINX := LENG; (* THE LENGTH ARRAY. *)
          RWLENG (.LINX.) := WINX
        END
      END;
      RWCHAR (.RWCHARCT + 1.) := BLANK; (* SET DELIMITERS FOR *)
      RWCEB (.RWORDCT + 1.) := RWCHARCT + 1; (* ARRAYS AND *)
      RWLENG (.RWLENGCT + 1.) := RWORDCT + 1 (* INDICES *)
    END;
  (* LCADRW *)
(*-----*)

```

```
(*--91-02-----*)
{* LOADLIB READS PREDEFINED FUNCTION LIBRARIES IN BIFNLF *}
{* AND BIFNQF FILES; MAKES APPROPRIATE SYM TBL ENTRIES. *}
(*-----*)
```

```
PROCEDURE LOADLIB (VAR LIBFILE : TEXT; FNTYPE : IDTYP;
                  SEQLEN : INTEGER);
VAR  IDSICT : SLCTPTR;
     I : INTEGER;

BEGIN
  READLN (LIBFILE); READLN (LIBFILE); (* SKIP HEAD LINES *)
  WHILE NOT EOF(LIBFILE) DO
  BEGIN
    ACCINX := 0; (* INIT ACCUM INDX *)
    REPEAT (* READ NAME OF FN *)
      ACCINX := ACCINX + 1;
      READ (LIBFILE, ACCUM(.ACCINX.));
    UNTIL ACCUM(.ACCINX.) = BLANK; (* TO 1ST BLANK *)
    FOR I := ACCINX TO MAXTOKLEN DO (* FILL REST BLANK *)
      ACCUM(.I.) := BLANK;
    ACCINX := ACCINX - 1; (* SET INDEX BACK TO NAME LEN *)
    IDSLOT := GETSLOT(ACCUM, ACCINX); (* ENTER IN SYMTBL *)
    IDSLOT@.TYP := FNTYPE; (* SET IDENT TYPE *)
    FOR I := 1 TO SEQLEN DO (* READ KEY CODES *)
      CASE FNTYPE OF
        FNCID : READ (LIBFILE, IDSLOT@.FNQ(.I.));
        FNLID : BEGIN
          IDSLOT@.FNLLINK := NIL;
          READ (LIBFILE, IDSLOT@.FNL(.I.))
        END
      END; (* CASE *)
    READLN (LIBFILE) (* SKIP TO NEXT LN *)
  END;
END; (* LOADLIB *)
```

```
(*-----*)
(*--91-03-----*)
{* LOADCTEXT READS THE CTEXTF FILE AND LOADS THE DATA *}
{* STRUCTURE WHICH WILL PROVIDE THE TRANSLATIONS OF TI-59 *}
{* KEY CCDES DURING THE FINAL CODE DUMP. *}
(*-----*)
```

```
PROCEDURE LOADCTEXT;
VAR  I, K : INTEGER;
     J, CH : 1..TEXTLEN + 1;
     CH : CHAR;

BEGIN
  READLN (CTEXTF); READLN (CTEXTF);
  WHILE NOT EOF(CTEXTF) DO
  BEGIN
    READ (CTEXTF, I, CTEXT(.I.).UNIT); (* SKIP TWO BLANKS *)
    READ {CTEXTF, CH, CH};
    J := 1;
    WHILE NOT EOLN(CTEXTF) DO
    BEGIN
      READ {CTEXTF, CTEXT(.I.).CODECHAR(.J.)};
      J := J + 1
    END;
    FOR K := J TO TEXTLEN DO
      CTEXT(.I.).CCDECHAR(.K.) := BLANK;
    READLN (CTEXTF);
  END;
END; (* LOADCTEXT *)

(*-----*)
```

BEGIN (* INITIALIZE MAIN *)

(*-----*)
(* OPEN ALL FILES AND WRITE OUTPUT FILE HEADERS. *)
(*-----*)

```
TERMOU1 (OUTFILE);
RESET (EASGF, 'NAME=BASICF.WBASIC.A');
RESET (MSGF, 'NAME=MSGF.PASCAL.A');
RESET (RWTBLF, 'NAME=RWTBLF.PASCAL.A');
RESET (LABELF, 'NAME=LABELF.PASCAL.A');
RESET (CTEXTF, 'NAME=CTEXTF.PASCAL.A');
RESET (EIFNOF, 'NAME=BIFFNOF.PASCAL.A');
RESET (EIFNLF, 'NAME=BIFFNLF.PASCAL.A');
REWRITE (LISTF, 'NAME=LISTF.WBASIC.A');
REWRITE (NAMEF, 'NAME=NAMEF.WBASIC.A');
WRITLN (LISTF, MSGF, 2); (* HEADER MSG TO LISTF *)
WRITLN (OUTFILE, MSGF, 1); (* TERMINAL INVOKE MSG *)
WRITLN (NAMEF, MSGF, 8); (* HEADER MSG TO NAMEF *)
```

(*-----*)
(* INITIALIZE OPTION TOGGLES *)
(*-----*)

```
LINK59 := FALSE; (* OPTION 0 *)
PC100 := TRUE; (* OPTION 1 *)
OPTPAR := TRUE; (* OPTION 2 *)
OPTNOF := TRUE; (* OPTION 3 *)
CODUMP := TRUE; (* OPTION 4 *)
SYDUMP := FALSE; (* OPTION 5 *)
DSDUMP := FALSE; (* OPTION 6 *)
TOKCUT := FALSE; (* OPTION 7 *)
TOKLIS := FALSE; (* OPTION 8 *)
```

(*-----*)
(* INITIALIZE RESERVED WORD ARRAY INDEXES. *)
(*-----*)

LOADRW (RWTBLF);

(*-----*)
(* INITIALIZE CHARACTER SETS. *)
(*-----*)

```
LETTERS := (.A..I.) + (.J..R.) + (.S..Z.);
DIGITS := (.0..9.);
ALFANUM := LETTERS + DIGITS;
SIGNS := (.+.) + (.-.) + (.*) + (./.);
DOUELE1 := (.<.) + (.>.) + (.!) + (.@.);
DOUELE2 := (.>.) + (.=.) + (.*) + (.@.);
SPECIALS := (.+.) + (.-.) + (.*) + (./.);
SUBERRCR := (.ENDLIN.) + (.ENDFIL.) + (.ENDFIL.);
CRITICAL := (.ENDLIN.) + (.ENDFIL.) + (.ENDFIL.);
TRAILTCK := (.CMTCKEXC.) + (.ENDLINTOK.) + (.ENDFILTOK.);
BINCPICKS := (.PLUSTOK.) + (.MINUSTOK.) + (.MULTOK.);
RELCPICKS := (.EQUALTOK.) + (.NOTEQTOK.) + (.GTEQTCK.);
NUMERICKEY := (.LTECTOK.) + (.LTTOCK.) + (.GTTOK.);
NUMERICKEY := (.K_DECP.) + (.K_EE.) + (.K_NEG.);
NUMERICKEY := (.K_ZERO..K_9.);
SIGNTOKS := (.PLUSTOK.) + (.MINUSTOK.);
BEGIN_EXPRTOKS := SIGNTOKS + (.IDENTOK.) + (.OPARENTOK.);
```

```

{*-----*}
{* INITIALIZE HASH TABLE AND REGISTER COUNT. *}
{*-----*}

FOR I := 0 TO HASHEASE DO
    EUCKET (.I.) := NIL;
NEXTREG := STARTREG;

{*-----*}
{* INITIALIZE ARRAY HOLDING OUTPUT TEXT OF TI-59 CODE. *}
{*-----*}

LOADCTEXT;

{*-----*}
{* INITIALIZE BUILT-IN FUNCTION LIBRARY. *}
{*-----*}

LOADLIB (EIPNOF, FNCID, FNOLEN);
LOADLIB (EIPNLF, FNIID, FNLEN);

{*-----*}
{* ENTER 'PI' = 3.14159265359 IN SYMBOL TABLE. *}
{*-----*}

ACCUM (.1.) := 'P'; ACCUM (.2.) := 'I';
ACCINX := 2;
FOR I := 3 TO MAXTCKLEN DC
    ACCUM (.I.) := BLANK;
IDSLOT := GETSLOT (ACCUM, ACCINX);
IDSLOT@.TYP := VARID;
IDSLOT@.REGNO := -314; (* SPECIAL REGNO FOR 'PI' *)

{*-----*}
{* INITIALIZE LABEL STACK (ARRAY OF INTEGER KEY CODES). *}
{*-----*}

READLN (LABELF); (* SKIP HEAD LINE *)
FOR I := 1 TO LBLBASE DO
    READ (LABELF, CLAEEL (.I.));
LBLCT := 1;

{*-----*}
{* INITIALIZE RESERVED REGISTER SET. *}
{*-----*}

READLN (LABELF); READLN (LABELF); READLN (LABELF);
RESERVECT := 0;
RESERVE REG := (.); (* INITIALIZE TO EMPTY SET *)
WHILE NOT EOF (LABELF) DO
    BEGIN
        WHILE NOT EOLN (LABELF) DO
            BEGIN
                READ (LABELF, I);
                RESERVE REG := RESERVE REG + (.I.); (* MAKE SET *)
                RESERVECT := RESERVECT + 1 (* COUNT MEMBERS *)
            END;
        READLN (LABELF)
    END;

{*-----*}
{* INITIALIZE FNF ACTIVATION STACK AND FNL USE LIST. *}
{*-----*}

FNSTACKCT := 0;
FNSTACK := NIL;
FNLIST := NIL;

```

```
{*-----*}
{* INITIALIZE FIRST CALL TO SCAN. *}
```

```
LINEUF (.0.) := BLANK;
LINEUF (.1.) := ENDLIN;
TOKNUM := ENDLINTCK;
LNEINX := 0;
ERRCRCT := 0;
WARNCT := 0;
FLAGCMT := FALSE;
LINUM := 0;
LLINUM := 0;
CIINUM := 0;
```

```
{*-----*}
{* INITIALIZE LINKED DATA STRUCTURE FOR TI-59 CODE. *}
```

```
FIRSTLP := GETNEWHDR (LINUM); (* SET A COMMON REF NODE *)
LPCUR := FIRSTLP; (* ANCHOR ALL MARKER PTRS TO IT *)
LASTLP := FIRSTLP;
ENDCP := FIRSTLP@.CPTR;
BEGINCF := ENDCP;
CPCUR := ENDCP;
SETLINE (LPCUR, LP); (* SET UP FOR MAIN PROCEDURE LABEL *)
GENKEY (R_LBL);
GENKEY (NEWLBL); (* MAIN PROCEDURE = LBL A *)
BEGINCF := BEGINCP@.SEQ; (* BYPASS THE HEADER NODE *)
```

```
{*-----*}
{* INITIALIZE LOOP/BRANCH STACKS. *}
```

```
IFSTACK := NIL;
ENDIFSTACK := NIL;
LOOFSTACK := NIL;
ENDLOCFSTACK := NIL;
FORSTACK := NIL;
NEXTSTACK := NIL;
```

```
{*-----*}
{* INITIALIZE READ/DATA STATEMENT INDEXES/FLAGS. *}
```

```
READIX := 1;
DATAIX := 1;
INDEXERROR := FALSE;
FIRSTREAD := TRUE
```

END;

(* INITIALIZE *)

```

*****
*
*           FAX59:  MAIN DRIVER
*
*****

```

```
BEGIN  (* FAX59 MAIN *)
```

```
  INITIALIZE;
```

```
  REPEAT  (* UNTIL TCKNUM = ENDFILTOK *)
```

```

    SCAN (TCKNUM);          (* SCAN FIRST WORD OF NEW LINE *)
    IF ERRORCT = 0 THEN    (* PARSING IS DISCONTINUED AFTER *)
      BEGIN                (* FIRST FATAL ERROR ENCOUNTERED *)
        SETLINE (LPCUR, LP); (* NEW WBASIC LINE NO & LINE *)
        CASE TOKNUM CF     (* RECURSIVE DESCENT PARSE PRCC *)

```

```

-----*
* KEYWORDS MARKED IN RIGHT CMT COLUMN BY ASTERISKS MUST
* ALWAYS RESULT IN A PARSE ERROR IF USED AS A COMMAND
* (IE. 1ST WORD ON A LINE) REGARDLESS OF IMPLEMENTATION:
* ** IMPLEMENTED IN THIS SUBSET
* *** NOT IMPEMENTED IN THIS SUBSET
-----*

```

```

ERRORTOK      : BEGIN END;          (* SCAN ERROR *)
CMTOKEXC      : PREM;               (* EXCLAM *)
2,3,4,5,6     : ERROR;              (* 1-CHAR SYMBOLS *)
7,8,9,10,11   : ERROR;              (* 1-CHAR SYMBOLS *)
12,13,14      : BEGIN END;          (* SCAN ERROR *)
15,16,17,18   : ERROR;              (* 2-CHAR SYMBOLS *)

```

```

19            : PIF;                  (* IF *)
20            : FERROR;               (* TO ** *)
21            : FERROR;               (* OR *** *)
22            : PSUBERROR;            (* ON *)
23            : FLET;                 (* LET *)
CM:TOK REM    : FREM;                 (* REM *)
25            : FFOR;                 (* FOR *)
26            : FEND;                 (* END *)
27            : FDEF;                 (* DEF *)
28            : FERROR;               (* NOT ** *)
29            : FSUBERROR;            (* DIM *)
30            : FERROR;               (* AND *** *)
31,32         : PSUBERROR;            (* THEN ** *)
33            : FERROR;               (* ELSE *)
34            : FELSE;                (* GOTO *)
35            : FGOTO;                (* LOOP *)
36            : FLLOOP;               (* NEXT *)
37            : FNEXT;                (* QUIT *)
38            : FQUIT;                (* STOP *)
39            : FSTOP;                (* DATA *)
40            : FDATA;                (* READ *)
41            : FREAD;                (* STEF ** *)
42            : FERROR;               (* *)
43,44,45      : PSUBERROR;            (* ENDFIF *)
46            : FENDIF;               (* FNEND *)
47            : PFNEND;               (* GOSUB *)
48            : FGOSUB;               (* INPUT *)
49            : PINPUT;               (* UNTIL *)
50            : FUNTIL;               (* WHILE *)
51            : FWHILE;               (* PAUSE *)
52            : FPAUSE;               (* PRINT *)
53            : FPRINT;               (* *)
54,55,56,57,58 : FSUBERROR;

```


APPENDIX D

RWTBLF FILE--ORDERED RESERVED WORDS

1	!
2	=
3	+
4	*
5	/
6	(
7)
8	<
9	>
10	~
11	^
12	%
13	~
14	~
15	~
16	~
17	<>
18	<=>
19	~
20	~
21	~
22	~
23	~
24	~
25	~
26	~
27	~
28	~
29	~
30	~
31	~
32	~
33	~
34	~
35	~
36	~
37	~
38	~
39	~
40	~
41	~
42	~
43	~
44	~
45	~
46	~
47	~
48	~
49	~
50	~
51	~
52	~
53	~
54	~
55	~
56	~
57	~
58	~
59	~

60 RETURN
61 CPTICN
62 LINFUT
63 BEMCVE
64 RENAME
65 RESUME
66 UNLCK
67 ENDLOOP
68 RESTORE
69 SCRATCH
70 TAGSORT
71 ENDGUESS
72 RANDOMIZE

APPENDIX E

LABEL FILE--TI-59 LABELS/RESERVED REGISTERS

KEY CCDES FOR TI-59 LABELS:

11	12	13	14	15	16	17	18	19	10
20	22	23	24	25	27	28	29	30	32
33	34	35	36	37	38	39	42	43	44
45	47	48	49	50	52	53	54	55	57
58	59	60	61	65	66	67	68	69	70
71	75	76	77	78	79	80	81	85	86
87	88	89	90	91	93	94	95	96	97
98	99								

REGISTERS RESERVED BY USER:

00 01 02 03 04 05 06 07 08 09 10

APPENDIX F

BIPNQF/BIPNLF FILES--BUILT-IN FUNCTIONS

BUILT-IN "QUICK" FUNCTION NAMES AND TI-59 KEY CODES:

ABS	50	68	68	68
ACOS	27	39	68	68
ASIN	27	38	68	68
ATN	27	30	68	68
COS	39	68	68	68
COT	30	35	68	68
CSC	38	35	68	68
EXP	22	23	68	68
FP	27	59	68	68
IP	59	68	68	68
LOG	23	68	68	68
LOG 10	28	68	68	68
SEC	39	35	68	68
SIN	38	68	68	68
SQR	34	68	68	68
TAN	30	68	68	68

BUILT-IN "LCNG" FUNCTION NAMES AND TI-59 KEY CODE SEQUENCES:

RND 36 15 10 43 10 36 15 15 36 15 71 88 68 68 68

APPENDIX G
CTEXTF FILE--TI-59 KEYCODE TRANSLATIONS

TI-59 KEY CCDE TEXT.....

-2	0	UNRESOLVED ADDR \$\$\$\$
-1	00	UNFILLED CODE \$\$\$\$\$\$
00	00	0
01	00	1
02	00	2
03	00	3
04	00	4
05	00	5
06	00	6
07	00	7
08	00	8
09	00	9
10	00	2ND F'
11	00	A'
12	00	B'
13	00	C'
14	00	D'
15	00	E'
16	00	2ND A'
17	00	2ND B'
18	00	2ND C'
19	00	2ND D'
20	00	2ND CLR
21	00	2ND \$\$\$\$ ERROR \$\$\$\$\$
22	00	INV
23	00	LNK
24	00	CE
25	00	CLR
26	00	2ND \$\$\$\$ ERROR \$\$\$\$\$
27	00	2ND INV
28	00	2ND LCG
29	00	2ND CP
30	00	2ND TAN
31	00	LEN (DEBUGGING TOOL)
32	00	X<=>T
33	00	X**2
34	00	SCRT (X)
35	00	1/X
36	1	2ND FGM
37	00	2ND P=>R
38	00	2ND SIN
39	00	2ND COS
40	00	2ND IND
41	00	SSI \$\$\$\$ ERROR \$\$\$\$\$
42	1	SIC
43	1	RCI
44	1	SUM
45	00	Y**X
46	00	INS \$\$\$\$ ERROR \$\$\$\$\$
47	00	2ND CMS
48	00	2ND EXC
49	1	2ND PRD
50	00	IXI
51	00	EST \$\$\$\$ ERROR \$\$\$\$\$
52	00	EE
53	00	(
54	0)

55	0	/				
56	0	DEL	\$\$\$\$	ERROR	\$\$\$\$	
57	0	2ND	ENG			
58	0	2ND	FIX			
59	0	2ND	INT			
60	0	2ND	DEG			
61	2	GTC				
62	1	2ND	PGM	2ND	IND	
63	1	2ND	EXC	2ND	IND	
64	1	2ND	PRD	2ND	IND	
65	0	*				
66	0	2ND	PAUSE			
67	2	2ND	X=T			
68	0	2ND	NOP			
69	1	2ND	CF			
70	0	2ND	RAD			
71	2	SER				
72	1	SIC	2ND	IND		
73	1	RCL	2ND	IND		
74	1	SUM	2ND	IND		
75	0	-				
76	0	2ND	ISL			
77	2	2ND	X>=T			
78	C	2ND	SUMMATION			
79	0	X-EAR				
80	0	2ND	GRAD			
81	0	RST				
82	0	\$\$\$	VOID CODE	\$\$\$\$		
83	1	GTC	2ND	IND		
84	1	2ND	OP	2ND	IND	
85	0	+				
86	1	2ND	STFLG			
87	3	2ND	IFFLG			
88	0	2ND	D.MS			
89	0	2ND	FI			
90	0	2ND	LIST			
91	0	R/S				
92	0	INV	SBR			
93	0	.				
94	0	+/-				
95	0	=				
96	0	2ND	WRITE			
97	3	2ND	CSZ			
98	0	2ND	ADV			
99	0	2ND	FRT			

APPENDIX I
LINKER SOURCE CODE

```

*****
* PURPOSE: THIS PROGRAM TAKES AS INPUT A TI-59 PROGRAM. *
* IT SEGMENTS THE PROGRAM SO THAT IT WILL FIT *
* INTO THE TI-59 CALCULATOR. INSTRUCTIONS AND *
* CODE LISTINGS ARE PROVIDED AS OUTPUT. *
* COMMENT: PROGRAM MAY LOOP INFINITELY IF SMALL; *
* LIMIT IS USED BECAUSE OF DIVIDE *
* ALGORITHM. *
* COMMENT: FILEDEFS WERE USED FOR THIS PROGRAM *
* CONSEQUENTLY THEY WERE NOT DEFINED *
* IN THE PROGRAM. SPECIFIC FILEDEFS *
* FOLLOW: SCRATCH--"SCRATCH PASCAL" *
* PASSED FROM COMPILER *
* OUTFILE--"ANY DESIRED NAME" *
* YOUR OUTPUT FILE *
* TEMPFILE--"ANY DESIRED NAME" *
* A TEMPORARY SCRATCH PAD *
* MESSAGEFILE--MESSAGEFILE FILE *
* LINKER'S MESSAGES *
*****

```

PROGRAM TSDRIVER (INPUT,OUTPUT) ;

```

*****
* DECLARATIONS: *
*****

```

```

(*-----*)
CONST
  FJUMCONST      = 10;      (* NUM STEPS FOR F JUMP CODE *)
  SERCCNST       = 15;      (* NUM STEPS FOR SBR BRK CCDE *)
  SERCCNTCONST   = 7;      (* NUM STEPS FOR SBR BRK RTN *)

  STO = 42;              (* TIS9 KEYCODES *)
  LEL = 76;
  RCLINE = 73;
  STOIND = 72;
  CF = 69;
  DECIMAL = 93;
  BS = 91;
  CE = 24;              (* END KEYCODES *)

  DISPLAYREGSTORE = 00;    (* TEMP STORE OF THE DISPLAY *)
  RTNREGNUM = 6;          (* NUMBER OF MANUAL RETURN REGISTERS *)
  MANRTNREG = 08;        (* MANUAL SER RETURN REGISTER *)

  NCNE = 101;            (* MESSAGE NUMS *)
  ASTER = 102;          (* SEE MESSAGEFILE FOR TRANSLATION *)
  YES = 103;
  MCDFFCMFTS = 100;
  RTNRGTOP = 104;
  CCDENUM = 0;
  STOINRG = 105;
  FGMEPARTIS = 106;
  PARINUMIS = 107;
  MCDN = 108;
  CARD1 = 109;
  CARD2 = 110;

```



```

SIDE1 = 111;
SIDE2 = 112;
EAXINSTR = 81;
SPECIFICS = 82;
ENDLEI = 83;
FAILINSTR = 84;
UNSEGCCLBL = 5;
PSEQ = 6;
PMANETN = 7;
PFWDJ = 8;
PSERINV = 9;
REGMAF = 3;
CATAREAD = 4;
RGCT = 1;
ALPHALBL = 99;
(* END MESSAGE NUMS *)
(*-----*)
TYPE
(*-----*)
LABELS = PACKED ARRAY (.1..15.) OF CHAR;
TYPELABELS = ARRAY (.0..99.) OF LABELS; (* TI59 KEYS *)
(*-----*)
CODEPTR = @CODERCD; (* THIS RECORD IS BUILT_CODE *)
CCDCRCD = RECORD
    AADDR:INTEGER;
    MEMNUM:INTEGER;
    RADDR:INTEGER;
    KEY:INTEGER;
    JMP:CCDEPTR;
    SEQ:CODEPTR;
END; (* SHOULD HAVE MADE A VARI *)
(*-----*)
INSTR_SET = SET CF 0 .. 99 ; (* RANGE INSTRUCTION SET*)
(*-----*)
NCDE = (TABLE, SBFPTR, SBRBREAK, FWD_JUMP, MEMODULE, CODE);
IPLPTR = @NODES;
NCDES = RECORD
    CASE TAG: NODE OF
        TABLE: (NEST:INTEGER;
            START_ADDR:INTEGER;
            STOP_ADDR:INTEGER;
            LENGTH:INTEGER;
            INCLUDED:BOOLEAN;
            COALESCED:BOOLEAN;
            SBRLIST:TBLPTR;
            NUM_F:INTEGER;
            F_JUMLIST:TBLPTR;
            TABLELIST:TBLPTR); (* NEXT TABLE *)
        SBRPTR: (SBR:TBLPTR;
            PRGM:INTEGER;
            NEXT_SBR:TBLPTR); (* NEXT PTR *)
        MEMODULE: (MEMNUM:INTEGER;
            OFFSET:INTEGER;
            HIGHOFFSET:INTEGER;
            LOWOFFSET:INTEGER;
            RETURNCODE_NEEDED:BOOLEAN;
            SEGTLBLS:TBLPTR; (* TABLE *)
            CODELIST:TBLPTR;
            NEXT:TBLPTR);

```

```

CODE: (ADDRESS:INTEGER;
      ABS_ADDR:INTEGER;
      KEYCODE:INTEGER;
      SEQUENTIAL:TBLPTR);

SBREREAK: (SBRZ:TBLPTR);

FWD_JUMP: (JUMP_ADDRFR:INTEGER;
          JUMP_ADDRTO:INTEGER;
          JUMP_ADDRTO1:INTEGER; (*HUNDREDS*)
          JUMP_ADDRTO2:INTEGER; (*TEN/UNIT*)
          MEM_ADDR:INTEGER; (*MEMNUM*)
          JUMP_INTADDRTO1:INTEGER; (*LOCAL*)
          JUMP_INTADDRTO2:INTEGER;
          NEXT_PJUMP:TBLPTR);

```

END;

```

(*----- (* SEGMENT TABLE STRUCTURE *)-----*)
(*-----*)
VAR
  CUTFILE:TEXT; (* OUTPUT FILE *)
  TEMPFILE:TEXT; (* NEST DIAGS TEMP FILE *)
  SCRATCH:TEXT; (* INFORMATION FROM COMPILER FILE *)
  MESSAGEFILE:TEXT; (* MESSAGE INPUT FILE *)
(*-----*)

  PARTIICN:REAL; (* CALCULATOR PARTITION INFO *)
  REGCCUNT:INTEGER;
  GOOD_SEGMENT:BOOLEAN;
  SERINVNEST:INTEGER; (* SBR NEST LEVEL CHECK *)
  NUMBANKS,PARI_NUM:INTEGER;
  LIMIT:INTEGER; (* MEMORY SIZE LIMIT *)
(*-----*)

  BUILT_CCDE,CURCP:CODEPTR; (* CODE TBL VARS*)
  BUILT_CCDE_COUNT:INTEGER;
  HDRPTR,SEG_TBL:TBLPTR; (* TABLE VARS *)
(*-----*)

  STEP_0,STEP_1,STEP_2,STEP_3:INSTR_SET; (* INS SET VAR *)
(*-----*)

  TLEL:TYPELABELS; (* PROGRAM LAEELS *)
(*-----*)

```

```

{*****}
{* PRCCEDURES AND FUNCTIONS: *}
{*****}

```

```

{=====}
{* FCLLOWING ROUTINES ARE USED AS UTILITIES SUCH AS PRINT *}
{* AND SCRATCHFILE AND MESSAGE FILE MANIPULATORS *}
{=====}

```

```

{*-----*}
{* DUMP SEGTBL: DUMPS THE SEGMENT TABLE. USED FOR DEEUG *}
{* ANE IS NOT CALLED IN THIS PROGRAM. *}
{*-----*}
PROCEDURE DUMP_SEGTBL(VAR OUTFILE:TEXT; HDRPTR:TBLPTR);
VAR F_JMPLINK,SBRLINK,SBR,CURTP,SBRTP:TBLPTR;
BEGIN
  SBRTP:=HDRPTR;
  WHILE SBRTP <> NIL DO
    BEGIN
      WRITELN(OUTFILE);
      WRITELN(OUTFILE);
      WRITELN(OUTFILE,'=====');
      WRITE(OUTFILE,' SBR CODE NUMBER ');
      WRITELN(OUTFILE,SBRTP@.STOP_ADDR:2);
      WRITELN(OUTFILE,'=====');
      CURTP:=SBRTP@.TABLELIST;
      WHILE CURTP <> NIL DO
        BEGIN
          WITH CURTP@ DO
            BEGIN
              WRITELN(OUTFILE);
              WRITELN(OUTFILE,'NEST',NEST:3);
              WRITELN(OUTFILE,'START',START_ADDR:4);
              WRITELN(OUTFILE,'STOP',STOP_ADDR:4);
              WRITELN(OUTFILE,'LENGTH',LENGTH:5);
            END;
          SBRLINK:=CURTP@.SERLIST;
          F_JMPLINK:=CURTP@.F_JUMPLIST;
          WHILE (SERLINK <> NIL) OR (F_JMPLINK <> NIL) DO
            BEGIN
              IF SERLINK <> NIL THEN
                BEGIN
                  SBR:=SERLINK@.SBR;
                  CASE SBR@.TAG OF
                    TABLE: SBR:=SBR;
                    SBRBREAK:
                      BEGIN
                        SBR:=SBR@.SBRZ;
                        WRITELN(OUTFILE,'*** BREAk ***');
                      END;
                  END;
                  WRITE(OUTFILE,'SBR INVOKE FROM',
                    SBRLINK@.FROM:5,' TO ',
                    SBR@.START_ADDR:5,' *** ');
                  SBRLINK:=SBRLINK@.NEXT_SBR;
                END
              ELSE
                WRITE(OUTFILE,'
                WRITE(OUTFILE,' *** ');
            IF F_JMPLINK <> NIL THEN
              BEGIN
                WRITE(OUTFILE,'JUMP FROM ',
                  F_JMPLINK@.JUMP_ADDR:5,' TC ',
                  F_JMPLINK@.JUMP_ADDRTO:5);
                F_JMPLINK:=F_JMPLINK@.NEXT_FJUMP;
              END;
            WRITELN(OUTFILE);
          END;
        END;
      END;
    END;
  END;

```

```

        CURIP:=CURTP@.TABLELIST;
        END;
        SBRTIP:=SBRTF@.SBRLIST;
    END;
(* DUMP_SEGTBL TEST ROUTINE *)
(*-----*)

```

```

(*-----*)
* DUMP MEMODULENODES: PRINTS OUT THE CONTENTS OF THE *
* MEMODULENODE LIST FORMED. THIS IS A DEBUGGING *
* ROUTINE AND IS NOT INVOKED IN THE PROGRAM. *
(*-----*)
PROCEDURE DUMP MEMODULENODES (HEAD_MEMODULE:TBLPTR);
VAR S:TBLPTR;
BEGIN
    S:=HEAD MEMODULE;
    WHILE S<>NIL DO
        BEGIN
            WITH S@ DO
                BEGIN
                    WRITELN(OUTFILE);
                    WRITE(OUTFILE,'MEMNUM OFFSET HIGH LOW');
                    WRITELN(OUTFILE,' SEGTBLSTART');
                    WRITELN(OUTFILE, MEMNUM:6, OFFSET:8, HIGHOFFSET:6,
                        LOWOFFSET:5, SEGTBLS@.START_ADDR:10);
                    WRITELN(OUTFILE);
                END;
            S:=S@.NEXT;
        END;
    END;
(* DUMP_MEMODULENODES *)
(*-----*)

```

```

(*-----*)
* WRITE LEADZERO: PADS INTEGER FIELD WITH LEADING *
* ZEROS *
(*-----*)
PROCEDURE WRITE LEADZERO (VAR OUTFILE:TEXT; NUM, PLD:INTEGER);
VAR I, TN:INTEGER;
BEGIN
    TN:=NUM;
    REPEAT
        TN:=TN DIV 10;
        FLD:=PLD-1;
    UNTIL (TN=0);
    FOR I:=1 TO PLD DO
        WRITE(OUTFILE,'0');
        WRITE(OUTFILE, NUM:1);
    END;
(* WRITE_LEADZERO *)
(*-----*)

```

```

{ *-----* }
{ * WRITELEL: WRITES OUT THE TI-59 CODED LABELS * }
{ *-----* }
PROCEDURE WRITELEL (VAR OUTFILE:TEXT; CODESS:INTEGER) ;
  BEGIN
    WRITELN (OUTFILE,TILBL (.CODESS.)) ;
  END; (* WRITELEL *)
{ *-----* }

```

```

{ *-----* }
{ * WRITCCDES: WRITES THE ADDRESS AND KEYCODE TO LINE * }
{ *-----* }
PROCEDURE WRITCCDES (VAR OUTFILE:TEXT; CUR:CODEPTR) ;
  BEGIN
    WRITE LEADZERO (OUTFILE,CUR@.AADDR, 3) ;
    WRITE (OUTFILE, ' ');
    WRITE LEADZERO (OUTFILE,CUR@.KEY, 2) ;
    WRITE (OUTFILE, ' ');
  END; (* WRITCCDES *)
{ *-----* }

```

```

{ *-----* }
{ * WRITENUM: WRITES KEYC DE AS A NUMBER NOT A LABEL * }
{ *-----* }
PROCEDURE WRITENUM (VAR OUTFILE:TEXT; CUR:CODEPTR) ;
  BEGIN
    WRITE LEADZERO (OUTFILE,CUR@.KEY, 2) ;
  END; (* WRITENUM *)
{ *-----* }

```

```

{ *-----* }
{ * HANDLE_#STEPS: PRINTS OUT DIFFERENT CASES OF CODES, * }
{ * EG. -WHETHER ONE OR TWO STEP INSTRUCTION. * }
{ * USED FOR CODEPTR TYPE OF NODES. * }
{ *-----* }
PROCEDURE HANDLE_0STEP (VAR OUTFILE:TEXT;
  VAR CUR:CODEPTR) ;
  BEGIN
    WRITE (OUTFILE, ' ');
    WRITCCDES (OUTFILE,CUR) ;
    WRITE (OUTFILE, ' ');
    WRITELEL (OUTFILE,CUR@.KEY) ;
  END;

PROCEDURE HANDLE_1STEP (VAR OUTFILE:TEXT;
  VAR CUR:CODEPTR) ;
  BEGIN
    CUR:=CUR@.SEC;
    WRITE (OUTFILE, ' ');
    WRITCCDES (OUTFILE,CUR) ;
    WRITE (OUTFILE, ' ');
    WRITENUM (OUTFILE,CUR) ;
    WRITELN (OUTFILE) ;
  END; (* HANDLE_1STEP *)

```

```

PROCEDURE HANDLE_2STEP (VAR OUTFILE:TEXT;
                       VAR CUR:CODEPTR);
  VAR I:INTEGER;
  BEGIN
    FOR I:=1 TO 2 DO
      BEGIN
        CUR:=CUR@.SEQ;
        WRITE (OUTFILE, ' ');
        WRITECODES (OUTFILE,CUR);
        WRITE (OUTFILE, ' ');
        WRITENUM (OUTFILE,CUR);
        WRITELN (OUTFILE);
      END;
    END; (* HANDLE_2STEP *)

PROCEDURE HANDLE_3STEP (VAR OUTFILE:TEXT;
                       VAR CUR:CODEPTR);
  VAR I:INTEGER;
  BEGIN
    FOR I:=1 TO 3 DO
      BEGIN
        CUR:=CUR@.SEQ;
        WRITE (OUTFILE, ' ');
        WRITECODES (OUTFILE,CUR);
        WRITE (OUTFILE, ' ');
        WRITENUM (OUTFILE,CUR);
        WRITELN (OUTFILE);
      END;
    END; (* HANDLE_3STEP *)
(*-----*)

(*-----*)
(* PRINT CODELIST: PRINTS OUT THE TI-59 CODE FOR *)
(* CODEPTR NODES ONLY. *)
(*-----*)
PROCEDURE PRINT_CODELIST (VAR OUTFILE:TEXT;
                          VAR BUILTCODE:CODEPTR);
  VAR CUR:CODEPTR;
  BEGIN
    CUR:=BUILTCODE;
    WHILE CUR <> NIL DO
      BEGIN
        HANDLE_0STEP (OUTFILE,CUR);
        IF CUR@.KEY IN (.71,76.) THEN
          BEGIN
            CUR:=CUR@.SEQ;
            HANDLE_0STEP (OUTFILE,CUR);
          END
        ELSE
          BEGIN
            IF CUR@.KEY IN STEP 1 THEN
              HANDLE_1STEP (OUTFILE,CUR);
            IF CUR@.KEY IN STEP 2 THEN
              HANDLE_2STEP (OUTFILE,CUR);
            IF CUR@.KEY IN STEP 3 THEN
              HANDLE_3STEP (OUTFILE,CUR);
          END;
        IF CUR <> NIL THEN
          CUR:=CUR@.SEQ;
        END;
      END;
    END; (* PRINT CODELIST *)
(*-----*)

```

```

(*-----*)
{* FIND_MSG: SEARCHES INPUT FILE TO FIND MSG NUMBER. *}
(*-----*)
PROCEDURE FIND_MSG (VAR MESSAGEFILE:TEXT; MSG:INTEGER);
VAR C1:CHAR; DIGIT:INTEGER;
BEGIN
  RESET (MESSAGEFILE);
  C1:=' ';
  DIGIT:=-1;
  REPEAT
    READ (MESSAGEFILE,C1);
    IF C1 = 'S' THEN
      READLN (MESSAGEFILE,DIGIT)
    ELSE
      READLN (MESSAGEFILE);
  UNTIL ((C1='S') AND (DIGIT=MSG));
END;
(*-----*)
(* FIND_MSG *)

```

```

(*-----*)
{* INIT_SETS: INITIALIZES IMPORTANT DATA SUCH AS KEY- *}
{* CCTE LABEL ARRAY, STEP SETS, KEY VARIABLES AND *}
{* INITIALIZES THE SCRATCH FILE *}
(*-----*)
PROCEDURE INIT_SETS (VAR TEMPFILE:TEXT; VAR STEP 0,STEP 1,
STEP 2,STEP 3:INSTR SET;
VAR GOOD SEGMENT:BOCLEAN; VAR MESSAGEFILE:TEXT;
VAR TILBI:TYPELABELS; VAR SBRINVNEST:INTEGER);
VAR C:CHAR;
DIGIT,J,L,K,I:INTEGER;

```

```

(*-----*)
{* GET_REGCOUNT: GOES TO SCRATCH FILE AND FINDS THE *}
{* MESSAGE NUMBER CONTAINING THE REGISTER COUNT *}
(*-----*)
PROCEDURE GET_REGCOUNT (VAR REGCOUNT:INTEGER);
BEGIN
  FIND_MSG (SCRATCH,REGCT);
  READLN (SCRATCH,REGCOUNT);
END;
(*-----*)
(* GET_REGCOUNT *)

```

```

BEGIN
  SBRINVNEST:=0; (* INITIALIZES THE INVCKE NEST CHECK *)

  RESET (MESSAGEFILE); (* INITIALIZE TILABELS *)
  DIGIT:=-1; (* LABELS IN MESSAGEFILE *)
  L:=1;
  REPEAT
    READ (MESSAGEFILE,C);
    IF C = 'S' THEN
      READLN (MESSAGEFILE,DIGIT)
    ELSE
      READLN (MESSAGEFILE);
  UNTIL (C = 'S') AND (DIGIT = ALPHALBL);
  L:=0;
  FOR I:=0 TO ALPHALBL DO
    BEGIN
      IF NOT (I IN (.21,26,31,41,46,51,56,82.)) THEN
        BEGIN
          READ (MESSAGEFILE,TILBL (.I.));
          L:=L+1;
          IF L = 4 THEN
            BEGIN
              READLN (MESSAGEFILE);
              L:=0;
            END;
        END;
    END;
END

```

```

        ELSE
        TILBL (. I.) := 'BLANK          ' ;
    END;

    GET_REGCOUNT (REGCOUNT) ;

    REWRITE (TEMPPFILE) ;
    WRITELN (TEMPPFILE, ' 39' ) ; (* OPENNING AND MARKING *)
    REWRITE (OUTFILE) ;          (* INIT OUTPUTFILE *)

    GOCD_SEGMENT := TRUE ;

    STEP_3 := (. 87, 97.) ;      (* STEP TYPES OF INSTRUCTIONS *)
    STEP_2 := (. 61, 67, 77.) ;
    STEP_1 := (. 36, 40, 42, 43, 44, 48, 49, 58, 62, 63, 64, 69, 72, 73, 74,
               83, 84, 86.) ;
    STEP_0 := (. 0. . 99.) - (STEP_3 + STEP_2 + STEP_1) ;
    END ;                          (* INIT_SETS *)
(*-----*)

{ * ADVANCE CODEPTR: MOVES ALONG CODE SKIPPING 1, 2, OR 3* }
{ * STEP INSTRUCTIONS AND STOPS ON NEXT COMMAND INSTR. * }
{ * TREATS 71 AND 76 AS SINGLE STEPS. * }
{ *-----* }
PROCEDURE ADVANCE_CODEPTR (VAR CUR: CODEPTR) ;
    VAR L: INTEGER ;
    BEGIN
        IF CUR@ . KEY IN STEP_3 THEN
            BEGIN
                FOR L := 1 TO 4 DO
                    IF CUR@ . SEQ <> NIL THEN
                        CUR := CUR@ . SEQ
                    END
                END
            END
        ELSE
            IF CUR@ . KEY IN STEP_2 THEN
                BEGIN
                    FOR L := 1 TO 3 DO
                        IF CUR@ . SEQ <> NIL THEN
                            CUR := CUR@ . SEQ
                        END
                    END
                END
            ELSE
                IF CUR@ . KEY IN STEP_1 THEN
                    BEGIN
                        FOR L := 1 TO 2 DO
                            IF CUR@ . SEQ <> NIL THEN
                                CUR := CUR@ . SEQ
                            END
                        END
                    END
                ELSE
                    IF CUR@ . SEQ <> NIL THEN
                        CUR := CUR@ . SEQ ;
                    END
                END
            END ;
    END ;
(*-----*)
(* END ADVANCE_CODEPTR *)

```



```

(*-----*)
* PRINTLN MSG: PRINTS A SPECIFIC MSG FROM ONE FILE *
* TO ANOTHER FILE. THIS ROUTINE WILL TAKE THE *
* WHOLE MESSAGE AND PRINT IT. IT EXECUTES A WRITELN *
* AT THE END OF THE PRINT *
(*-----*)
PROCEDURE PRINTLN_MSG (VAR OUTFILE, MESSAGEFILE:TEXT;
                      MSG:INTEGER);
  VAR C1:CHAR;
  BEGIN
    FIND MSG (MESSAGEFILE, MSG);
    READ MESSAGEFILE, C1);
    WHILE C1<>'$' DO
      BEGIN
        WRITE (OUTFILE, C1);
        WHILE NOT EOLN (MESSAGEFILE) DO
          BEGIN
            READ (MESSAGEFILE, C1);
            WRITE (OUTFILE, C1);
          END;
        READLN (MESSAGEFILE);
        WRITELN (OUTFILE);
        READ (MESSAGEFILE, C1);
      END;
    END;
  END;
(*-----*)
(* PRINTLN_MSG *)

```

```

(*-----*)
* PRINT MSGLINE: PRINTS A SPECIFIC ONE-LINE MESSAGE *
* TO ANOTHER FILE. DOES NOT WRITELN TO FILE. *
* USED FOR LINE LABELS OF GENERATED DATA. *
(*-----*)
PROCEDURE PRINT_MSGLINE1 (VAR OUTFILE, MESSAGEFILE:TEXT;
                          MSG:INTEGER);
  VAR C1:CHAR;
  BEGIN
    FIND MSG (MESSAGEFILE, MSG);
    READ MESSAGEFILE, C1);
    WHILE C1<>'$' DO
      BEGIN
        WRITE (OUTFILE, C1);
        READ (MESSAGEFILE, C1);
      END;
    END;
  END;
(*-----*)
(* PRINT_MSGLINE1 *)

```

```

(*-----*)
* DET LIMIT: DETERMINES MEMORY LIMITS BASED ON REG CNT.*
* ONLY THREE PARTITIONS WERE CONSIDERED. THIS WAS *
* BECAUSE ANY OTHER PARTITION SPLITS THE SIDE *
* OF A MAG CARD BETWEEN REGISTERS AND PROGRAM. *
* THIS WOULD CHANGE REGISTERS DURING REPROGRAMMING *
* AND IS THEREFORE UNACCEPTABLE. *
(*-----*)
PROCEDURE DET_LIMIT (VAR REGCOUNT, LIMIT, NUMBANKS,
                    PART_NUM:INTEGER;
                    VAR PARTITION:REAL);
  BEGIN
    IF REGCOUNT+RTNREGNUM IN (.0..29.) THEN
      BEGIN
        NUMBANKS:=3;
        PARTITION:=719.29;
        PART_NUM:=3;
        LIMIT:=719;
      END
    END;
  END;

```

```

ELSE
  IF REGCOUNT+RINREGNUM IN (.30..59.) THEN
    BEGIN
      NUMBANKS:=2;
      PARTITION:=479.59;
      PART_NUM:=6;
      LIMIT:=479;
    END
  ELSE
    BEGIN
      NUMBANKS:=1;
      PARTITION:=239.89;
      PART_NUM:=9;
      LIMIT:=239;
    END
  END;
END;
(*-----(* DET_LIMIT *)-----*)

(*-----*)
(* CLEAN: REMOVES SAME F JUMPS AND SAME SBRS IN A SEG *)
(* TEL NODE. ALSO GIVES DELETE COUNT FOR F_JUMPS *)
(* INCLUDED IN THE CONFINES OF THE SEGMENT. DOUBLE *)
(* DUTY ROUTINE. USED BY COMBINE AND BY SET_LENGTH *)
(*-----*)
PROCEDURE CLEAN(VAR CURTP:TBLPTR; VAR DELETE:INTEGER);
VAR F,S:TBLPTR;

(*-----*)
(* PRUNE_SAMEF: REMOVES SAME PJUMP ADDRESSTO FROM TEL *)
(*-----*)
PROCEDURE PRUNE_SAMEF(VAR F:TBLPTR);
VAR T,S:TBLPTR;
BEGIN
  WHILE F@.NEXT_FJUMP<>NIL DO
    BEGIN
      S:=F@.NEXT_PJUMP;
      T:=F;
      WHILE (S<>NIL) DO
        IF S@.JUMP_ADDRTO = F@.JUMP_ADDRTO THEN
          BEGIN
            T@.NEXT_FJUMP:=S@.NEXT_PJUMP;
            DISPOSE(S,FWD_JUMP);
            S:=T@.NEXT_FJUMP;
          END
        ELSE
          BEGIN
            T:=T@.NEXT_FJUMP;
            S:=S@.NEXT_PJUMP;
          END;
        IF F@.NEXT_FJUMP<>NIL THEN
          F:=F@.NEXT_PJUMP;
        END;
      END;
    END;
  END;
END;
(*-----(* PRUNE_SAMEF *)-----*)

```

```

{ *-----*
* PRUNE_GREATOR: REMOVES FJUMPS CONTAINED IN SEGTBL *
*-----*
PROCEDURE PRUNE_GREATOR (VAR F,S:TBLPTR;
                        VAR DELETE:INTEGER);
BEGIN
  WHILE F<>NIL DO
    BEGIN
      IF F@.JUMP_ADDRT@<=CURTP@.STOP_ADDR THEN
        BEGIN
          S@.NEXT_FJUMP:=F@.NEXT_FJUMP;
          DISPOSE(F,FWD_JUMP);
          F:=S@.NEXT_FJUMP;
          DELETE:=DELETE+1;
        END
      ELSE
        BEGIN
          S:=S@.NEXT_FJUMP;
          F:=F@.NEXT_FJUMP;
        END;
      END;
    IF CURTP@.F_JUMLIST@.JUMP_ADDRT@<=CURTP@.STOP_ADDR
      THEN
      BEGIN
        F:=CURTP@.F_JUMLIST;
        CURTP@.F_JUMLIST:=F@.NEXT_FJUMP;
        DISPOSE(F,FWD_JUMP);
        DELETE:=DELETE+1;
      END;
    END;
  END;
  (* PRUNE GREATOR *)
{ *-----*

```

```

{ *-----*
* PRUNE_SAMES: REMOVES SAME SBR INVOKES FROM SEGTBL *
*-----*
PROCEDURE PRUNE_SAMES (VAR F:TBLPTR);
  VAR S,T,SS,PF:TBLPTR;

{ *-----*
* PASS_BRK: PASSES OVER THE SBR BREAK NODE *
*-----*
FUNCTION PASS_BRK (F:TBLPTR):TBLPTR;
BEGIN
  CASE F@.TAG OF
    SBRBREAK: PASS_BRK:=F@.SBRZ;
    TABLE: PASS_BRK:=F;
  END;
END;
  (* PASS_BRK *)
{ *-----*

```

```

BEGIN
  WHILE F<>NIL DO
    BEGIN
      S:=F@.NEXT_SBR;
      IF S<>NIL THEN
        BEGIN
          PF:=PASS_BRK (F@.SBR);
          SS:=PASS_BRK (S@.SBR);
        END;
      T:=F;
      WHILE (S<>NIL) DO
        IF SS = PF THEN
          BEGIN
            T@.NEXT_SBR:=S@.NEXT_SBR;
            DISPOSE(S,SBRPTR);
            S:=T@.NEXT_SBR;
            IF S<>NIL THEN
              SS:=PASS_BRK (S@.SBR);
            END
          END
        END
      END
    END
  END

```

```

ELSE
  BEGIN
    T:=T@.NEXT_SBR;
    S:=S@.NEXT_SBR;
    IF S<>NIL THEN
      SS:=PASS_BRK(S@.SBR);
    END;
    F:=P@.NEXT_SBR;
  END;
END;
(* PRUNE_SAMESER *)
-----
BEGIN
  DELETE:=0;
  IF CURTP@.F_JUMPLIST<>NIL THEN
    BEGIN
      F:=CURTP@.F_JUMPLIST;
      PRUNE_SAMEF(F);
    END;
  IF CURTP@.F_JUMPLIST<>NIL THEN
    BEGIN
      S:=CURTP@.F_JUMPLIST;
      F:=S@.NEXT_FJUMP;
      PRUNE_GREATOR(F,S,DELETE);
    END;
  IF CURTP@.SBRLIST<>NIL THEN
    BEGIN
      F:=CURTP@.SERLIST;
      PRUNE_SAMES(F);
    END;
  END;
(* CLEAN *)
-----

(*-----*)
(* DIAGS_NEST1SBRBRK: DIAGNOSTIC PRINTOUT IF THERE IS A *)
(* A SER BREAK WITHIN AN ITERATIVE LOOP. NEEDS TO SET *)
(* GOOD_SEGMENT VARIABLE FALSE *)
(*-----*)
PROCEDURE DIAGS_NEST1SBRBRK(VAR TEMPPFILE:TEXT; SEG:TBLPTR;
  VAR GOOD_SEGMENT:BOOLEAN);
  VAR IS_BRK_BELOW:BOOLEAN;

  (*-----*)
  (* BELOW_BREAK: SEARCHES OUT BELOW TO SEE IF A BREAK *)
  (* IS PRESENT SO DIAGS_NEST1 CAN CHECK FOR A BREAK *)
  (* WITHIN A LOOP *)
  (*-----*)
  PROCEDURE BELOW_BREAK(SEG:TBLPTR; VAR IS_BRK_BELOW:
    BOOLEAN);
    VAR SER,SBRL:TBLPTR;
    BEGIN
      IF NOT IS_BRK_BELOW THEN
        BEGIN
          IF SEG@.SERLIST<>NIL THEN
            BEGIN
              SBRL:=SEG@.SBRLIST;
              WHILE SBRL<>NIL DO
                BEGIN
                  SER:=SBRL@.SBR;
                  IF SBR@.TAG=SBRBREAK THEN
                    IS_BRK_BELOW:=TRUE;
                  ELSE
                    BELOW_BREAK(SBR,IS_BRK_BELOW);
                  SERL:=SBRL@.NEXT_SBR;
                END;
            END;
          END;
        END;
    END;
  END;
(* BELOW_BREAK *)

```

```

(*-----*)
BEGIN
  IS BRK BELOW:=FALSE;
  BEICW BREAK(SEG,IS BRK BELOW);
  IF (IS BRK_BELOW) AND (SEG@.NEST=1) THEN
    BEGIN
      GCOD_SEGMENT:=FALSE;
      WRITELN(TEMPFILE);
      WRITE(TEMPFILE,');
      WRITELN(TEMPFILE, '* SBR BREAK WITHIN A LOOP');
      WRITE(TEMPFILE,');
      WRITELN(TEMPFILE, 'LOOP BOUNDS',SEG@.START_ADDR:4
        , ' TO ',SEG@.STOP_ADDR:4);
    END;
  END;
(*-----*)
(* DIAGS_NEST1SBRBRK *)

```

```

(*-----*)
{ * DIAGS_NEST1LENGTHCHK: PRINTS OUT DIAGNOSTIC IF THERE * }
{ * EXISTS AN ITERATIVE LOOP OF TOO GREAT A LENGTH. * }
{ * TAKES INTO ACCUNT OUT OF LOOP JUMPS. NEEDS TO * }
{ * SET GOOD_SEGMENT FALSE IF ENCOUNTERED * }
(*-----*)
PROCEDURE DIAGS_NEST1LENGTHCHK (VAR TEMPFILE:TEXT;
  CUR:TBLPTR; VAR GOOD_SEGMENT:BOOLEAN);
  BEGIN
    IF (CUR@.LENGTH > LIMIT) AND (CUR@.NEST=1) THEN
      BEGIN
        GCOD_SEGMENT:=FALSE;
        WRITELN(TEMPFILE);
        WRITE(TEMPFILE,');
        WRITELN(TEMPFILE, '* BACK JUMP NEST TOO LONG');
        WRITE(TEMPFILE,');
        WRITELN(TEMPFILE, 'LOOP BOUNDS',
          CUR@.START_ADDR:4, ' TO ',CUR@.STOP_ADDR:4);
      END;
    END;
  END;
(*-----*)
(* DIAGS_NEST1SERBRK *)

```

```

(*-----*)
* DIAGS_NEST6SBRINVCHK: CHECKS THAT THE SBR NEST LEVEL *
* DOES NOT EXCEED 6 *
*-----*)
PROCEDURE DIAGS_NEST6SBRINVCHK(VAR TEMPPFILE:TEXT;
                               CUR:TBLPTR; VAR GOOD_SEGMENT:BOOLEAN;
                               SBRINVNEST:INTEGER);
BEGIN
  IF SBRINVNEST > 7 THEN
    BEGIN
      GCOD_SEGMENT:=FALSE;
      WRITE(TEMPPFILE,' ');
      WRITELN(TEMPPFILE,'* SBR INVOKE NEST LEVEL > 6');
      WRITE(TEMPPFILE,' ');
      WRITE(TEMPPFILE,' CALLES ROUTINE STARTS ');
      WRITELN(TEMPPFILE,'AT ABS ADDR ',CUR@.START_ADDR:3)
    END;
  END;
(*-----*) (* DIAGS_NEST6SBRINVCHK *)

```

```

(*-----*)
* RESET_INCLUDED: SETS ALL INCLUDES TO FALSE. DOES SO *
* FOR ALL ROUTINES ON THE SBRLIST AND BELOW SBRS *
*-----*)
PROCEDURE RESET_INCLUDED(VAR SBRL:TBLPTR);
VAR SBRLST,SBR:TBLPTR;
BEGIN
  SBRLST:=SBRL;
  WHILE SBRLST<>NIL DO
    BEGIN
      SER:=SERLIST@.SBR;
      CASE SER@.TAG OF
        TABLE:
          SBR:=SER;
          SBRBREAK:
          SBR:=SER@.SBR2
      END;
      IF (SER@.SERLIST<>NIL) AND (SER@.COALESCED=
                                     TRUE) THEN
        RESET_INCLUDED(SER@.SERLIST);
      IF (SER@.COALESCED=TRUE) THEN
        SER@.INCLUDED:=FALSE;
      SERLIST:=SERLIST@.NEXT_SBR;
    END;
  END;
(*-----*) (* RSET_INCLUDED *)

```

```

(*-----*)
* INPUT: *
* PURPOSE: TO READ AN INPUT FILE AND FORM SEQ LINKS. *
* THIS FORMS THE INTERNAL CODE STRUCTURE WHICH WILL *
* BE MANIPULATED *
*-----*)
PROCEDURE INPUT(VAR SCRATCH:TEXT; VAR BUILT_CODE:CODEPTR;
                VAR BUILT_CODE_COUNT:INTEGER);
VAR ADDRESS:INTEGER;
    TEMP,COUNT:INTEGER;
    CUR,TRAIL:CODEPTR;
BEGIN
  FIND_MSG(SCRATCH,CODENUM);
  READ(SCRATCH,TEMP);
  IF TEMP > -1 THEN
    BEGIN
      NEW(CUR);
      BUILT_CODE:=CUR;
    END;
  END;

```

```

        COUNT:=0;
        CUR@.AADDR:=TEMP;
        CUR@.RADDR:=COUNT;
        READLN(SCRATCH,CUR@.KEY);
        TRAIL:=CUR;
    END; (*IF*)

REPEAT
    NEW(CUR);
    COUNT:=COUNT+1;
    READ(SCRATCH,CUR@.AADDR);
    IF CUR@.AADDR<>-1 THEN
        BEGIN
            CUR@.RADDR:=COUNT;
            READLN(SCRATCH,CUR@.KEY);
            TRAIL@.SEQ:=CUR;
            TRAIL@.JMP:=NIL;
            TRAIL:=CUR;
        END;
    UNTIL (CUR@.AADDR = -1);
    BUILT_CODE_COUNT:=COUNT-1;
    TRAIL@.JMP:=NIL;
    TRAIL@.SEQ:=NIL

END;
(*----- (* INPUT *)-----*)

(*-----*)
* SETJMP:
* PURPCSE: TO SET THE JUMP POINTER OF THE BUILT CODE
*-----*)
PROCEDURE SETJMPS (VAR BUILT_CODE:CODEPTR);
    VAR CUR:CODEPTR;

    (*-----*)
    * SETJMP_PTR: SETS THE JUMPTR OF THE CURRENT NODE
    *-----*)
    PROCEDURE SETJMP_PTR (VAR BUILT_CODE:CODEPTR;
                           CUR:CODEPTR);
        VAR MARKER,SEARCH:CODEPTR;
            ADDRESS:INTEGER;
        BEGIN
            MARKER:=CUR@.SEQ;
            IF CUR@.KEY IN STEP 3 THEN
                MARKER:=MARKER@.SEQ;
                ADDRESS:=100*MARKER@.KEY;
                ADDRESS:=ADDRESS+MARKER@.SEQ@.KEY;
                SEARCH:=BUILT_CODE;
                WHILE (SEARCH@.AADDR<>ADDRESS) DO
                    SEARCH:=SEARCH@.SEQ;
                    MARKER@.SEQ@.JMP:=SEARCH;
                END;
            (* SETJMPS_PTR *)
        END;
    BEGIN
        CUR:=BUILT_CODE;
        WHILE CUR@.SEQ <> NIL DO
            BEGIN
                IF CUR@.KEY IN (STEP 2+STEP 3) THEN
                    SETJMP_PTR(BUILT_CODE,CUR);
                IF CUR@.KEY IN (.77,76.) THEN
                    CUR:=CUR@.SEQ@.SEQ;
                ELSE
                    ADVANCE_CODEPTR (CUR);
                END;
            END;
        END;
    (*----- (* TEST_SETJMP *)-----*)

```

```

(*=====*)
{* BUILD SEGMENT TABLE ROUTINES: ON THIS TABLE ALL *}
{* OF THE COALESCING IS DONE, AND NOT THE CODE. *}
(*=====*)
PROCEDURE BUILD_SEG_TBL (BUILT_CODE:CODEPTR;
                        VAR SEGTBL:TBLPTR; LIMIT:INTEGER;
                        BUILT_CODE_COUNT:INTEGER);
    VAR HDRPTR:TELPTR;

    (*-----*)
    (* BLK PRIMSEG_TBL: RESULTS IN A TABLE WITH CRITICAL *)
    (* POINTS IDENTIFIED. THESE ARE BACK JUMP POINTS *)
    (* TO AND FROM LOCATIONS. STOP IS STORED IN STOP *)
    (* ADDRESS OF THE FIRST NODE. *)
    (*-----*)
    PROCEDURE BLK_PRIMSEG_TBL (BUILT_CODE:CODEPTR;
                              VAR HDRPTR:TBLPTR);
        VAR CURCP:CODEPTR;
            CURTP:TBLPTR;

        (*-----*)
        (* PROCESS_SBRLEL: STORES THE SBRLEL IN THE HEADER *)
        (* SBRLEL AND PRODUCES THE FIRST SEGMENT OF THE *)
        (* SEGMENT TABLE FOR EACH SBR. *)
        (* THIS IS CONFUSING IN THAT THE SAME TYPE OF CODE *)
        (* IS USED TO STORE THE LABEL NAME AS IS USED FOR *)
        (* THE SEGTABLE. KEY FIELD REDEFINITIONS FOR THIS *)
        (* FUNCTION SO THAT THE NAME GOES INTO THE FIELD *)
        (* STOP_ADDR. THESE LABEL NODES ARE NEEDED TO BE *)
        (* ABLE TO SET THE SBR INVOKE POINTERS LATER ON. *)
        (*-----*)
        PROCEDURE PROCESS_SBRLEL (VAR CURCP:CODEPTR;
                                  VAR CURTP:TBLPTR);
            VAR TRAILTP:TELPTR;
                BEGIN
                    TRAILTP:=CURTP;
                    NEW(CURTP, TABLE);
                    CURTP@.TAG:=TABLE;
                    CURTP@.TABLELIST:=NIL;
                    CURTP@.COALESCED:=FALSE;
                    CURTP@.INCLUDED:=FALSE;
                    CURTP@.SBRLEL:=NIL;
                    CURTP@.START_ADDR:=CURCP@.ADDR;
                    CURTP@.STOP_ADDR:=-1;
                    IF CURCP@.KEY = 76 THEN
                        BEGIN
                            CURTP@.STOP_ADDR:=CURCP@.SEQ@.KEY;
                            CURCP:=CURCP@.SEQ@.SEQ;
                        END
                    ELSE
                        CURTP@.STOP_ADDR:=-1;
                    IF TRAILTP <> NIL THEN
                        TRAILTP@.SBRLEL:=CURTP;
                    END;
                (* PROCESS_SBRLEL *)
            (*-----*)

```



```

(*-----*)
* PROCESS SBRCCDE: PROCESS THE TI-59 SBR CODE FOR *
* CRITICAL INFO AND BUILDS THE PRIMITIVE SEGS *
(*-----*)
PROCEDURE PROCESS_SBRCCDE (VAR CURCP:CODEPTR;
                          VAR SBRHDRTP:TBLPTR);
  VAR TOPTP,CURTP:TBLPTR;

  (*-----*)
  * IS_BCK_JMP: DETERMINES IF THE JUMP IS BACKWARDS *
  (*-----*)
  FUNCTION IS_BCK_JMP (CURCP:CODEPTR):BOOLEAN;
  VAR ADDRESS:INTEGER;
  BEGIN
    IF (CURCP@.KEY IN STEP_2) THEN
      BEGIN
        ADDRESS:=CURCP@.SEQ@.SEQ@.JMP@.AADDR;
        IF ADDRESS > CURCP@.AADDR THEN
          IS_BCK_JMP:=FALSE
        ELSE
          IS_BCK_JMP:=TRUE;
        END;
      END;
    ELSE
      IF (CURCP@.KEY IN STEP_3) THEN
        BEGIN
          ADDRESS:=CURCP@.SEQ@.SEQ@.SEQ@.JMP@.AADDR;
          IF ADDRESS > CURCP@.AADDR THEN
            IS_BCK_JMP:=FALSE
          ELSE
            IS_BCK_JMP:=TRUE;
          END;
        END;
      ELSE
        IS_BCK_JMP:=FALSE;
      END;
    END;
  (*-----*)
  (* IS_BCK_JMP *)
  (*-----*)

```

```

(*-----*)
* APND JMP_TBL: DETERMINES ALL OUT OF CODE JUMPS *
* FROM A "FROM" ADDRESS TO A "TO" ADDRESS. *
(*-----*)
PROCEDURE APND_JMP_TBL (CURCP:CODEPTR;
                      VAR TOPTP:TBLPTR);
  VAR ADDRESSFR,ADDRESSTO:INTEGER;

  (*-----*)
  * INSERT CRITS: PLACES CRITICALS IN SEGTBL. *
  * CRITICAL IS ADDRESS WHERE A BACK JUMP NEST *
  * LEVEL CHANGE TAKES PLACE, IE START OR STOP. *
  (*-----*)
  PROCEDURE INSERT_CRITS (ADDRESS:INTEGER;
                        VAR TOPTP:TBLPTR);
  VAR CURTP,TRAILTP,INSERTTP:TBLPTR;
  BEGIN
    TRAILTP:=TOPTP;
    CURTP:=TCPTP@.TABLELIST;
    WHILE (CURTP@.START_ADDR < ADDRESS) AND
          (CURTP@.TABLELIST <> NIL) DO
      BEGIN
        TRAILTP:=CURTP;
        CURTP:=CURTP@.TABLELIST;
      END;
    NEW (INSERTTP, TABLE);
    INSERTTP@.TAG:=TABLE;
    INSERTTP@.START_ADDR:=ADDRESS;
    INSERTTP@.STOP_ADDR:=-2;
    IF (CURTP@.TABLELIST=NIL) AND
       (CURTP@.START_ADDR < ADDRESS) THEN

```

```

        BEGIN
            CURTP@.TABLELIST:=INSERTTP;
            INSERTTP@.TABLELIST:=NIL
        END
    ELSE
        IF (CURTP@.TABLELIST=NIL) AND
            (CURTP@.START_ADDR > ADDRESS) THEN
            BEGIN
                TRAILTP@.TABLELIST:=INSERTTP;
                INSERTTP@.TABLELIST:=CURTP
            END
        ELSE
            IF CURTP@.START_ADDR > ADDRESS THEN
                BEGIN
                    INSERTTP@.TABLELIST:=CURTP;
                    TRAILTP@.TABLELIST:=INSERTTP
                END
            ELSE
                DISPOSE(INSERTTP, TABLE);
            END;
        (* INSERT_CRITS *)
    (*-----*)

    (*-----*)
    (* SET_NESTS: SEARCHES THE PRIM SEGTBL AND MARKS *)
    (* AS 1 ALL OVERLAPPING BACK JUMPS TO DESIGNATE *)
    (* THAT THEY ARE IN A NO BREAK AREA *)
    (*-----*)
    PROCEDURE SET_NESTS (ADDRESSFR, ADDRESSTO:INTEGER;
                        VAR TOPTP:TELPTR);
        VAR CURTP:TELPTR;
        BEGIN
            CURTP:=TCPTP;
            WHILE CURTP@.START_ADDR <> ADDRESSTO DO
                CURTP:=CURTP@.TABLELIST;
            WHILE CURTP@.START_ADDR <> ADDRESSFR DO
                BEGIN
                    CURTP@.NEST:=1;
                    CURTP:=CURTP@.TABLELIST
                END;
            CURTP@.NEST:=0;
            IF ((CURTP@.TABLELIST <> NIL) AND (CURTP@.NEST =
                1)) THEN
                CURTP@.NEST:=1
            END;
        (* SET_NESTS *)
    (*-----*)
    BEGIN
        IF CURCP@.KEY IN STEP_2 THEN
            BEGIN
                ADDRESSTO:=CURCP@.SEQ@.SEQ@.JMP@.AADDR;
                ADDRESSFR:=CURCP@.SEQ@.SEQ@.AADDR
            END
        ELSE
            BEGIN
                ADDRESSTO:=CURCP@.SEQ@.SEQ@.SEQ@.JMP@.AADDR;
                ADDRESSFR:=CURCP@.SEQ@.SEQ@.SEQ@.AADDR
            END;
            INSERT_CRITS (ADDRESSFR, TOPTP);
            INSERT_CRITS (ADDRESSTO, TOPTP);
            SET_NESTS (ADDRESSFR, ADDRESSTO, TOPTP)
        END;
    (* APND_JMP_TEL *)
    (*-----*)
    BEGIN
        NEW (CURTP, TABLE);
        CURTP@.TAG:=TABLE;
        CURTP@.START_ADDR:=SBRHDRTP@.START_ADDR;
        CURTP@.NEST:=0;
        CURTP@.COALESCED:=FALSE;

```

```

CURTP@.INCLUDED:=FALSE;
CURTP@.TABLELIST:=NIL;
CURTP@.SERLIST:=NIL;
SERHDR:P@.TABLELIST:=CURTP;
TCPTP:=SERHDRTP;
WHILE ((CURCP@.KEY <> 76) AND (CURCP@.SEQ <> NIL ))
  DC
  IF IS_BACK_JMP(CURCP) THEN
    BEGIN
      APND_JMP_TBL(CURCP,TOPTP);
      ADVANCE_CODEPTR(CURCP)
    END
  ELSE
    ADVANCE_CCDEPTR(CURCP);
  IF CURCP@.SEC = NIL THEN
    TOPTP@.TABLELIST@.STOP_ADDR:=CURCP@.AADDR
  ELSE
    TOPTP@.TABLELIST@.STOP_ADDR:=CURCP@.AADDR-1;
END;
(*-----*)
(* PROCESS_SBRCODE *)
BEGIN
  CURCP:=BUILT_CODE;
  CURTP:=NIL;
  PROCESS_SBRLBL(CURCP,CURTP);
  HDRPTR:=CURTP;
  WHILE (CURCP@.SEQ <> NIL) DO
    BEGIN
      PROCESS_SERCODE(CURCP,CURTP);
      IF CURCP@.KEY = 76 THEN
        PROCESS_SBRLBL(CURCP,CURTP);
      END;
    END;
END;
(*-----*)
(* BLD_PRIMSEGTEL *)

{
* BLD_ADVSEGTEL: FILLS IN THE STOPS AND MERGES SAME
* NESTED LEVELS INTO ONE SEGMENT.
* STOPS ARE STOP_ADDR FIELD
*
}
PROCEDURE BLD_ADVSEGTEL(VAR HDRPTR:TBLPTR);
VAR SBRTT:TBLPTR; STOP:INTEGER;

{
*-----*
* MERGE_ONES: COMBINES SAME NESTED ADJACENT 1 SEGS *
*-----*
}
PROCEDURE MERGE_ONES(VAR SBRTT:TBLPTR);
VAR MARK,ZERO,ONE:TBLPTR;

{
*-----*
* MERGE: DOES ACTUAL MERGING OF ADJACENT SEGMENTS*
*-----*
}
PROCEDURE MERGE(VAR ONE,ZERO,MARK:TBLPTR);
VAR DIS:TBLPTR;
BEGIN
  ONE@.STOP_ADDR:=ZERO@.START_ADDR;
  DIS:=ONE@.TABLELIST;
  WHILE ONE@.TABLELIST <> ZERO DO
    BEGIN
      ONE@.TABLELIST:=DIS@.TABLELIST;
      DISPCSE(DIS,TABLE);
      DIS:=ONE@.TABLELIST;
    END;
  IF ZERO@.TABLELIST <> NIL THEN
    ONE@.TABLELIST:=ZERO@.TABLELIST
  ELSE
    ONE@.TABLELIST:=NIL;
  MARK:=ONE;

```

```

DISPOSE (LIS, TABLE) ;
END;
----- (* MERGE *)
BEGIN
MARK := SBRTPa.TABLELIST;
WHILE MARK@.TABLELIST <> NIL DO
BEGIN
IF (MARK@.NEST = 0) AND (MARK@.TABLELIST <> NIL)
THEN MARK := MARK@.TABLELIST;
ONE := MARK;
WHILE (MARK@.NEST = 1) AND (MARK@.TABLELIST <> NIL)
DO MARK := MARK@.TABLELIST;
ZERO := MARK;
IF ONE <> ZERO THEN
BEGIN
MERGE (ONE, ZERO, MARK);
MARK := ONE;
END;
IF MARK@.TABLELIST <> NIL THEN
MARK := MARK@.TABLELIST;
END;
END;
----- (* MERGE_ONES *)

```

```

----- (*
* ADD_ZEROS: FILLS IN GAPS IN TABLE WITH 0 SEG
* ----- *)
PROCEDURE ADD_ZEROS (VAR SBRTP: TBLPTR);
VAR CUR, TRAIL, INSERT: TBLPTR; STOP: INTEGER;
BEGIN
TRAIL := SBRTP;
CUR := SBRTPa.TABLELIST;
STOP := CUR@.STOP_ADDR;
WHILE CUR@.TABLELIST <> NIL DO
BEGIN
CUR := CUR@.TABLELIST;
WHILE TRAIL@.TABLELIST <> CUR DO
TRAIL := TRAIL@.TABLELIST;
IF TRAIL@.NEST <> CUR@.NEST THEN
TRAIL@.STOP_ADDR := CUR@.START_ADDR - 1
ELSE
BEGIN
NEW (INSERT, TABLE);
INSERT@.NEST := 0;
INSERT@.START_ADDR := TRAIL@.STOP_ADDR + 1;
INSERT@.STOP_ADDR := CUR@.START_ADDR - 1;
INSERT@.TABLELIST := CUR;
TRAIL@.TABLELIST := INSERT
END;
END;
IF CUR@.STOP_ADDR <> STOP THEN
BEGIN
NEW (INSERT, TABLE);
INSERT@.NEST := 0;
INSERT@.START_ADDR := CUR@.STOP_ADDR + 1;
INSERT@.STOP_ADDR := STOP;
INSERT@.TABLELIST := NIL;
CUR@.TABLELIST := INSERT
END;
END;
END;
----- (* ADD_ZEROS *)
BEGIN
SBRTP := HDRPTR;
WHILE SBRTP <> NIL DO
BEGIN
MERGE_ONES (SBRTP);
ADD_ZEROS (SBRTP);

```

```

        SB RTP:=SERTP@.SRLIST;
    END;
END;
(*-----* BLD_ADVSEGTEL *)
(*-----*)
(* BLD_FINSEGTL: PROCESS CODE FOR SBR INVOKES AND
 * FJUMPS. WHEN ENCOUNTERED IT PLACES INTO SEGTL.
 * THESE WILL INCLUDE ONLY ONE INVOKE PER SEGMENT
 * AND ONLY ONE FJUMP TO SAME LOCATIONS. REPEATS
 * WILL BE IGNORED. LENGTHS OF SEGMENTS WILL ALSO
 * BE CALCULATED. LENGTHS DO NOT INCLUDE CODE FOR
 * SER INVOKES/PROMPT CODE. ONLY SEQUENTIAL CONTIN-
 * UATION CODE IS INCLUDED IN LENGTH CALCULATION TO-
 * GETHER WITH FJUMP PROMPT CODE.
 *-----*)
PROCEDURE BLD_FINSEGTL (BUILTCODE:CODEPTR;
                       VAR HDRPTR:TBLPTR; LIMIT:INTEGER);
VAR CURCP:CODEPTR;
    SERTP:TBLPTR;

(*-----*
 * PROCESS SBRSEGTL: PLACES SBRS & FJMP INTO SEGTL*
 *-----*)
PROCEDURE PROCESS_SBRSEGTL (VAR CURCP:CODEPTR;
                            VAR HDRPTR,SB RTP:TBLPTR);
VAR CURTP,SB RINVOKE,FJMP:TBLPTR;

(*-----*
 * HANDLE_FWDJMP: INSESTS FWD JUMPS INTO TABLE.
 *-----*)
PROCEDURE HANDLE_FWDJMP (CURCP:CODEPTR;
                        VAR HDRPTR,CURTP,FJMP:TBLPTR);
VAR ADDRESSSTO,ADDRESSFR:INTEGER;
    INSERT:TBLPTR;
BEGIN
    IF CURCP@.KEY IN STEP_3 THEN
        BEGIN
            ADDRESSFR:=CURCP@.SEQ@.SEQ@.AADDR;
            ADDRESSSTO:=CURCP@.SEQ@.SEQ@.JMP@.AADDR;
        END
    ELSE
        BEGIN
            ADDRESSFR:=CURCP@.SFQ@.SEQ@.AADDR;
            ADDRESSSTO:=CURCP@.SEQ@.SEQ@.JMP@.AADDR;
        END;
        NEW (INSERT,FWD JUMP);
        INSERT@.TAG:=FWD JUMP;
        INSERT@.JUMP_ADDRFR:=ADDRESSFR;
        INSERT@.JUMP_ADDRTO:=ADDRESSSTO;
        INSERT@.JUMP_ADDRTO1:=-1;
        INSERT@.JUMP_ADDRTO2:=-2;
        INSERT@.MEM_ADDR:=-1;
        INSERT@.JUMP_INTADDRTO1:=-1;
        INSERT@.JUMP_INTADDRTO2:=-2;
        INSERT@.NEXT_FJUMP:=NIL;
        IF CURTP@.F_JUMPLIST = NIL THEN
            BEGIN
                CURTP@.F_JUMPLIST:=INSERT;
                FJMP:=INSERT;
            END
        ELSE
            BEGIN
                FJMP@.NEXT_FJUMP:=INSERT;
                FJMP:=INSERT;
            END;
        END;
END;
(*-----* HANDLE_FWDJMP *)
(*-----*)

```

```

(*-----*)
{* IS_FWD_JMP: BOOLEAN TRUE IF KEYCODE IS FWD JMP *}
{*-----*}
FUNCTION IS_FWD_JMP (CURCP:CODEPTR) : BOOLEAN;
VAR ADDRESS:INTEGER;
BEGIN
  IF (CURCP@.KEY IN STEP_2) THEN
    BEGIN
      ADDRESS:=CURCP@.SEQ@.SEQ@.JMP@.AADDR;
      IF ADDRESS > CURCP@.AADDR THEN
        IS_FWD_JMP:=TRUE
      ELSE
        IS_FWD_JMP:=FALSE;
    END
  ELSE
    IF (CURCP@.KEY IN STEP_3) THEN
      BEGIN
        ADDRESS:=CURCP@.SEQ@.SEQ@.SEQ@.JMP@.AADDR;
        IF ADDRESS > CURCP@.AADDR THEN
          IS_FWD_JMP:=TRUE
        ELSE
          IS_FWD_JMP:=FALSE;
        END
      ELSE
        IS_FWD_JMP:=FALSE;
    END;
  END;
  (* IS_FWD_JMP *)
(*-----*)

```

```

(*-----*)
{* HANDLE SBRINVOKE: PLACES SBR CALL INTO TABLE *}
{*-----*}
PROCEDURE HANDLE_SBRINVOKE(VAR CURCP:CODEPTR;
  VAR HDRPTR,CURTP,SBRINVOKE:TBLPTR);
VAR TOSBR,INSERT:TBLPTR;
KEYY:INTEGER;
BEGIN
  TOSBR:=HDRPTR;
  KEYY:=CURCP@.SEQ@.KEY;
  WHILE TOSBR@.STOP_ADDR<> KEYY DO
    TOSBR:=TOSBR@.SERLIST;
  NEW(INSERT,SBRPTR);
  INSERT@.TAG:=SBRPTR;
  INSERT@.FROM:=CURCP@.AADDR+1;
  INSERT@.SER:=TOSBR@.TABLELIST;
  INSERT@.NEXT_SBR:=NIL;
  IF CURTP@.SERLIST = NIL THEN
    BEGIN
      CURTP@.SERLIST:=INSERT;
      SBRINVOKE:=INSERT
    END
  ELSE
    BEGIN
      SBRINVOKE@.NEXT_SBR:=INSERT;
      SBRINVOKE:=INSERT;
    END;
  END;
  (* HANDLE_SBRINVOKE *)
(*-----*)

```

```

BEGIN
  CURTP:=SBRTP@.TABLELIST;
  REPEAT
    CURTP@.SERLIST:=NIL;
    CURTP@.F_JUMPLIST:=NIL;
    SBRINVOKE:=NIL;
    FJMP:=NIL;
  WHILE CURCP@.AADDR < CURTP@.STOP_ADDR DO
    BEGIN
      IF CURCP@.KEY = 71 THEN

```

```

        HANDLE_SBRINVOKE(CURCP, HDRPTR, CURTP,
                          SBRINVOKE);
        IF IS_FWD_JMP(CURCP) THEN
            HANDLE_FWDJMP(CURCP, HDRPTR, CURTP, FJMP);
        IF (CURCP@.KEY=76) OR (CURCP@.KEY=71) THEN
            CURCP:=CURCP@.SEQ@.SEQ
        ELSE
            ADVANCE_CODEPTR(CURCP);
        END;
        CURTP:=CURTP@.TABLELIST
    UNTIL (CURTF = NIL);
END;
(*-----(* PROCESS_SBRSEGTEL *)-----*)

```

```

(*-----*)
* SET LENGTH: ENSURES LENGTH IS WITHIN MEMORY LIMIT*
* IF NOT WILL DIVIDE THE SEGMENT IN HALF AND
* RESET ALL SBRLISTS AND PJUMPLISTS THEN CONTINUE*
* NOTE: MAY LEAD TO PROBLEMS IS LIMIT IS
* ARBITRARILY SMALL.
*-----*)
PROCEDURE SET_LENGTH(BUILT_CODE:CODEPTR;
                    VAR SBRTPT:TBLPTR; LIMIT:INTEGER);
    VAR CURTP:TBLPTR;
        LENGTH,DELETE,L_POSSBR:INTEGER;

    (*-----*)
    (* CALCULATE: DETERMINES LENGTH OF A SEGMENT. WILL*
    (* NOT ADD ADDITICNAL STEPS FOR DUPLICATE FJMP *
    (* ADDRTO
    (*-----*)
    PROCEDURE CALCULATE(CURTP:TBLPTR;
                       VAR LENGTH:INTEGER);
        VAR S,F:TBLPTR;
            ADDITICNS:INTEGER;
        BEGIN
            LENGTH:=CURTP@.STOP_ADDR-CURTP@.START_ADDR
                    +FJUMPCONST+1;
            IF CURTP@.F_JUMPLIST<>NIL THEN
                BEGIN
                    ADDITIONS:=0;
                    F:=CURTP@.F_JUMPLIST;
                    IF F@.JUMP_ADDRTO > CURTP@.STOP_ADDR THEN
                        ADDITIONS:=1;
                    S:=F;
                    F:=F@.NEXT_PJUMP;
                    WHILE F<>NIL DO
                        BEGIN
                            IF F@.JUMP_ADDRTO>CURTP@.STOP_ADDR THEN
                                BEGIN
                                    ADDITIONS:=ADDITIONS+1;
                                    S:=CURTP@.F_JUMPLIST;
                                    WHILE ((S<>F) AND (S@.JUMP_ADDRTO<>
                                                                    F@.JUMP_ADDRTO)) DO
                                        S:=S@.NEXT_PJUMP;
                                    IF S<>F THEN
                                        ADDITIONS:=ADDITIONS-1;
                                    END;
                                END;
                            F:=F@.NEXT_PJUMP;
                        END;
                    LENGTH:=LENGTH+(ADDITIONS)*(FJUMPCONST);
                END;
            END;
        END;
    (*-----(* CALCULATE *)-----*)

```

```

(*-----*)
{* L_POSSBRK: CALCULATES ANY SBR INVOKES AS A *}
{* _POSSIBLE BREAK FOR DIVISION PURPOSES. DOES *}
{* NOT INCLUDE MULTIPLE INVOKES OF SAME SBR *}
(*-----*)
PROCEDURE LENGTH_SBRBRKS (CURTP:TBLPTR;
                          VAR L_POSSBR:INTEGER);
VAR F,T:TBLPTR; COUNT:INTEGER;
BEGIN
  COUNT:=0;
  IF CURTP@.SBRLIST<>NIL THEN
    BEGIN
      F:=CURTP@.SBRLIST;
      WHILE F<>NIL DO
        BEGIN
          IF NOT (F@.SBR@.INCLUDED) THEN
            BEGIN
              COUNT:=COUNT+1;
              F@.SBR@.INCLUDED:=TRUE;
            END;
          F:=F@.NEXT_SBR;
        END;
      F:=CURTP@.SBRLIST;
      WHILE F<>NIL DO
        BEGIN
          F@.SBR@.INCLUDED:=FALSE;
          F:=F@.NEXT_SBR;
        END;
      END;
      L_POSSBR:=COUNT*SBRCONST;
    END;
  (* L_SBRBRKS *)
(*-----*)

```



```

*-----*
* DIVIDE: DIVIDES A SEG IN HALF AND RESETS FJMF *
* AND SBR PCINTERS. *
*-----*

```

```

PROCEDURE DIVIDE(BUILT_CODE:CODEPTR;
                 VAR CURTP:TBLPTR);
VAR INSERT:TELPTR;
    S,F:TBLPTR;
    NEW_STOP:INTEGER;

```

```

*-----*
* SETT: ENSURES THAT DIVIDE CALCULATED STOP IS *
* NOT SPLITTING A 1,2,3 PART INSTRUCTION. *
*-----*

```

```

PROCEDURE SETT(BUILT_CODE:CODEPTR;
              VAR NEW_STOP:INTEGER);
VAR P,T:CCDEPTR;
BEGIN
  P:=BUILT_CODE;
  WHILE (P@.AADDR <= NEW_STOP) DO
  BEGIN
    T:=P;
    IF (P@.KEY = 76) OR (P@.KEY = 71) THEN
      F:=P@.SEQ@.SEQ
    ELSE
      ADVANCE_CODEPTR(P);
    END;
    IF (T@.KEY = 76) OR (T@.KEY = 71) THEN
      T:=T@.SEQ@.SEQ;
    NEW_STOP:=T@.AADDR-1;
  END;
  (* SETT *)

```

```

*-----*
* FIND INSSBRLIST: DIVIDES UP THE SBRLIST *
* BETWEEN THE OLD AND NEW SEGMENTS *
*-----*

```

```

PROCEDURE FIND_INSSBRLIST(VAR CURTP,S:TBLPTR;
                          NEW_STOP:INTEGER);
VAR LIMIT:INTEGER;
    T:TBLPTR;
BEGIN
  LIMIT:=NEW_STOP+1;
  IF CURTP@.SBRLIST<> NIL THEN
  BEGIN
    S:=CURTP@.SBRLIST;
    IF (S@.NEXT_SBR<>NIL) AND (S@.FROM < LIMIT) THEN
    BEGIN
      S:=S@.NEXT_SBR;
      T:=CURTP@.SBRLIST;
      WHILE (S@.FROM<LIMIT) AND (S@.NEXT_SBR<>NIL) DO
      BEGIN
        S:=S@.NEXT_SBR;
        T:=T@.NEXT_SBR;
      END;
      IF S@.FROM >=LIMIT THEN
        T@.NEXT_SBR:=NIL
      ELSE
        S:=NIL
    END
  ELSE
  BEGIN
    IF S@.FROM >= LIMIT THEN
      CURTP@.SBRLIST:=NIL
    ELSE
      S:=NIL
  END

```

```

        END;
    END
  ELSE
    S:=NIL;
  END;
  (* FIND_INSSBRLIST *)
  -----*)

  (*-----*)
  (* FIND_INSFJUMPLIST: DIVIDES FJUMPLIST BETWEEN *)
  (* OLD AND NEW SEGMENTS. *)
  (*-----*)
  PROCEDURE FIND_INSFJUMPLIST (VAR CURTP,F:TBLPTR;
                               NEW_STOP:INTEGER);
    VAR LIMIT:INTEGER;
        T:TBLPTR;
    BEGIN
      LIMIT:=NEW_STOP+1;
      IF CURTP@.F_JUMPLIST<> NIL THEN
        BEGIN
          F:=CURTP@.F_JUMPLIST;
          IF (F@.NEXT_FJUMP<>NIL) AND (F@.JUMP_ADDRFR
                                       < LIMIT) THEN
            BEGIN
              F:=F@.NEXT_FJUMP;
              T:=CURTP@.F_JUMPLIST;
              WHILE (F@.JUMP_ADDRFR<LIMIT) AND
                    (F@.NEXT_FJUMP<>NIL) DO
                BEGIN
                  F:=F@.NEXT_FJUMP;
                  T:=T@.NEXT_FJUMP;
                END;
              IF F@.JUMP_ADDRFR >=LIMIT THEN
                T@.NEXT_FJUMP:=NIL
              ELSE
                F:=NIL
            END
          END
        ELSE
          BEGIN
            IF F@.JUMP_ADDRFR >= LIMIT THEN
              CURTP@.F_JUMPLIST:=NIL
            ELSE
              F:=NIL
            END;
          END;
        END
      ELSE
        F:=NIL;
      END;
    (* FIND_INSFJUMPLIST *)
  -----*)
  BEGIN
    NEW_STOP:=(((CURTP@.STOP_ADDR-CURTP@.START_ADDR
                  +1)DIV 2))+CURTP@.START_ADDR;
    SETT (BUILT_CODE,NEW_STOP);
    NEW (INSERT_TABLE);
    INSERT@.NEST:=CURTP@.NEST;
    INSERT@.START_ADDR:=NEW_STOP+1;
    INSERT@.STOP_ADDR:=CURTP@.STOP_ADDR;
    FIND_INSSERLIST (CURTP,S,NEW_STOP);
    FIND_INSFJUMPLIST (CURTP,F,NEW_STOP);
    INSERT@.SERLIST:=S;
    INSERT@.F_JUMPLIST:=F;
    INSERT@.TABLELIST:=CURTP@.TABLELIST;
    CURTP@.TABLELIST:=INSERT;
    CURTP@.STOP_ADDR:=NEW_STOP;
  END;
  (* DIVIDE *)
  -----*)
  BEGIN
    CURTP:=SBRTF@.TABLELIST;
  REPEAT

```

```

CALCULATE (CURTP, LENGTH);
IF (CURTP@.NEST = 0) THEN
  BEGIN
    LENGTH SBRBRKS (CURTP, L POSSBR);
    IF LENGTH+L POSSBR > LIMIT THEN
      DIVIDE (BUILT_CODE, CURTP)
    ELSE
      BEGIN
        CLEAN (CURTP, DELETE);
        DELETE:=0; (*CALCULATE HAS THIS COVERED*)
        CURTP@.LENGTH:=LENGTH-DELETE*PJUMPCONST;
        IF CURTP@.TABLELIST=NIL THEN
          CURTP@.LENGTH:=CURTP@.LENGTH-PJUMPCONST
          +SBRCONTCNST;
        CURTP:=CURTP@.TABLELIST;
      END
    END
  ELSE
    BEGIN
      CLEAN (CURTP, DELETE);
      DELETE:=0; (* CALCULATE HAS THIS COVERED *)
      CURTP@.LENGTH:=LENGTH-DELETE*PJUMPCONST;
      CURTP:=CURTP@.TABLELIST
    END;
  UNTIL (CURTF = NIL);
END; ----- (* SET_LENGTH *)
(*-----*)
BEGIN
  CURCP:=BUILTCCDE;
  SBRTPT:=HDRPTR;
  WHILE SBRTPT <> NIL DO
    BEGIN
      PROCESS SERSEGTBL (CURCP, HDRPTR, SBRTPT);
      SET_LENGTH (BUILT_CODE, SBRTPT, LIMIT);
      SBRTPT:=SBRTPT@.SBRLIST
    END;
  END; ----- (* BLD_FINSEGTEL *)
(*-----*)
BEGIN
  IF BUILT_CODE_COUNT<=LIMIT THEN
    BEGIN
      NEW (SEGTBL, TABLE);
      SEGTBL@.TAG:=TABLE;
      SEGTBL@.TABLELIST:=NIL;
      SEGTBL@.SBRLIST:=NIL;
      SEGTBL@.COALESCED:=TRUE;
      SEGTBL@.INCLUDED:=FALSE;
      SEGTBL@.START_ADDR:=0;
      SEGTBL@.STOP_ADDR:=BUILT_CODE@.SEQ@.KEY;
      NEW (HDRPTR, TABLE);
      HDRPTR@.START_ADDR:=0;
      HDRPTR@.STOP_ADDR:=BUILT_CODE_COUNT;
      HDRPTR@.NEST:=0;
      HDRPTR@.TABLELIST:=NIL;
      HDRPTR@.SBRLIST:=NIL;
      HDRPTR@.PJUMPLIST:=NIL;
      HDRPTR@.COALESCED:=TRUE;
      HDRPTR@.INCLUDED:=FALSE;
      HDRPTR@.LENGTH:=BUILT_CODE_COUNT+1;
      SEGTBL@.TABLELIST:=HDRPTR;
    END
  ELSE
    BEGIN
      BLD_PRIMSEGTEL (BUILT_CODE, HDRPTR);
      BLD_ADVSEGTEL (HDRPTR);
      BLD_FINSEGTEL (BUILT_CODE, HDRPTR, LIMIT);
      SEGTBL:=HDRPTR;
    END;
  END;
END; ----- (* BLD_SEGTEL *)

```

```

*=====*)
* COALESCE: CCOALESCE THE SEG TABLE MAKING GOOD BRK. *
* ONLY LOSS OF EFFICIENCY IS WITH CROSS SEGMENT *
* FORWARD JUMPS. THESE MAY PRECLUDE THE COMBINING *
* OF A SEGMENT BECAUSE OF ADDED CODE FOR THE JUMP *
*=====*)
PROCEDURE COALESCE (VAR SBR:TBLPTR; LIMIT:INTEGER;
VAR GOOD SEGMENT:EOLEAN; VAR SBRINVNEST:INTEGER);
VAR CURSEG:TBLPTR;

*-----*)
* SERSUM: SUMS ALL SBRS ON A SBRLIST. SETS INCLUDED *
* TRUE. ADDS SBRCONST IF SBRBREAK IS ENCOUNTERED. *
*-----*)
PROCEDURE SERSUM (VAR SBRLST:TBLPTR; VAR SUMSBR:INTEGER);
VAR SERL,SER:TEIPTR;
BEGIN
SERL:=SBRLST;
WHILE SBRL<>NIL DO
BEGIN
SBR:=SBRL@.SER;
CASE SBR@.TAG OF
TABLE:
BEGIN
IF (SER@.COALESCED) THEN
IF NOT (SER@.INCLUDED) THEN
BEGIN
SBR@.INCLUDED:=TRUE;
SUMSBR:=SUMSBR+SBR@.LENGTH
-SBRCONST;
IF SBR@.SBRLIST<>NIL THEN
SERSUM (SBR@.SERLIST,SUMSBR);
END;
END;
SBRBREAK:
SUMSER:=SUMSBR+SBRCONST;
END;
SBRL:=SBRL@.NEXT_SBR;
END;
END;
(* SERSUM *)
*-----*)

*-----*)
* SERSUMLINK: ADDS ALL SBRS BELOW AN INVOKE NODE. *
* IT DOES NOT DO THE WHOLE SBRLIST. THAT IS SUMSBR *
*-----*)
PROCEDURE SERSUMLINK (VAR SBRPTR:TBLPTR; VAR SUMT:INTEGER);
VAR PTR:TBLPTR;
BEGIN
PTR:=SBRPTR@.SER;
CASE PTR@.TAG OF
TABLE:
IF (PTR@.COALESCED) AND (PTR@.TABLELIST=NIL) THEN
IF (NOT (PTR@.INCLUDED)) THEN
BEGIN
PTR@.INCLUDED:=TRUE;
SERSUM (PTR@.SERLIST,SUMT);
SUMT:=SUMT+PTR@.LENGTH-SBRCONST;
END
ELSE
IF (PTR@.COALESCED) AND (NOT (PTR@.INCLUDED)) THEN
SUMT:=SUMT+PTR@.LENGTH-SBRCONST;
SBRBREAK:
SUMT:=SUMT;
END;
END;
(* SERSUMLINK *)
*-----*)

```

```

*-----*
*  CHK SEGSIZE:  VERIFIES THAT THE SEGMENT OF NEST 0  *
*  TOGETHER WITH ITS REQUIRED SBR INVOKES IS WITHIN *
*  MEMORY CONSTRAINTS.  IF IT DOES NOT FIT THEN A  *
*  SBRBREAK IS INSERTED.  THIS ROUTINE ASSUMES THAT *
*  THE SEGMENT LENGTH IS OK.  SET LENGTH ENSURES THIS *
*  IF A SBR IS NOT COALESCED THIS ROUTINE MUTUALLY *
*  RECURSES WITH COALESCE.  DURING ITS CHECK, IT  *
*  WILL INCLUDE AND RESET INCLUDES.  *
*-----*
PROCEDURE CHK_SEGSIZE(VAR CURSEG:TBLPTR; LIMIT:INTEGER);
VAR SEG,SBR,SBFL:TBLPTR;
    CHKLENGTH,SUMSBR:INTEGER;

*-----*
*  INSERT_SBRBRK:  INSERTS THE SBRBREAK NODE  *
*-----*
PROCEDURE INSERT_SBRBRK(SERNODE:TBLPTR);
VAR INSEPT,CUR,PWD:TBLPTR;
BEGIN
    CUR:=SERNODE;
    IF CUR@.SBR@.TAG <> SBRBREAK THEN
        BEGIN
            PWD:=CUR@.SBR;
            NEW(INSEPT,SBRBREAK);
            INSEPT@.TAG:=SBRBREAK;
            INSEPT@.SBRZ:=CUR@.SBR;
            CUR@.SER:=INSEPT;
        END;
    END;
END;
(* INSERT_SBRBRK *)

*-----*
BEGIN
    SEG:=CURSEG;
    WHILE SEG<>NIL DO
        BEGIN
            CHKLENGTH:=SEG@.LENGTH;
            IF SEG@.TABLELIST=NIL THEN
                CHKLENGTH:=CHKLENGTH-SBRCONTCONST;
            IF SEG@.SERLIST<>NIL THEN
                BEGIN
                    SBRL:=SEG@.SBRLIST;
                    WHILE SBRL<>NIL DO
                        BEGIN
                            SER:=SBRL@.SBR;
                            CASE SBR@.TAG OF
                                TABLE:
                                    BEGIN
                                        IF (SBR@.COALESCED=FALSE) THEN
                                            COALESCE(SBR,LIMIT,
                                                GOOD_SEGMENT,SBRINVNEST);
                                        IF SER@.TABLELIST<>NIL THEN
                                            BEGIN
                                                INSERT_SBRBRK(SBRL);
                                                END;
                                                SUMSBR:=0;
                                                SBRSUMLINK(SBRL,SUMSBR);
                                                CHKLENGTH:=CHKLENGTH+SUMSBR;
                                                IF (CHKLENGTH>LIMIT) THEN
                                                    BEGIN
                                                        INSERT_SBRBRK(SBRL);
                                                        CHKLENGTH:=CHKLENGTH-SUMSER;
                                                        RESET_INCLUDED(SEG@.SERLIST);
                                                    END;
                                                END;
                                                SBABREAK:
                                                END;
                                                SBRL:=SBRL@.NEXT_SBR;
                                            END;
                                        END;
                                    END;
                            END;
                        END;
                    END;
                END;
            END;
        END;
    END;
END;

```

```

        IF SEG@.SERLIST<>NIL THEN
            RESET INCLUDED(SEG@.SERLIST);
            DIAGS_NEST:=SBRBK(TEMPFILE,SEG,GOOD_SEGMENT);
            SEG:=SEG1.TABLELIST;
        END;
    END;
    (*----- (* CHK_SEGSIZE *) -----*)

(*----- (*
* CCMBINE: TAKES THE CHECKED SEGTABLE AND COMBINES
* IT INTO A MAXIMIZED COMBINATION OF SEGMENTS AND
* SBR. BASICALLY, IT MERGES THE ADJACENT SEGMENTS
* IF THEY CAN BE MERGED.
*----- *)
PROCEDURE COMBINE(VAR CURSEG:TBLPTR; LIMIT:INTEGER);
VAR SUMSBRF,SUMSBRZ,SUMFWD,SUMCJR,SUMTOT:INTEGER;
    CUR,FWD:TBLPTR;
    DELSUMSBRZ,DELSUMFJMP:INTEGER;

(*----- (*
* MERGES: MERGES ADJACENT SEGMENTS TO INCLUDE THEIR
* RESPECTIVE PJUMPLISTS AND SBRLISTS. IT THEN
* USES CLEAN TO REMOVE ANY DUPLICATE JUMPS/SEES
* USES CLEAN'S DELETE FACILITY TO READJUST LENGTH
* THE LENGTH FOR THE SBRBREAK CODE IS NOT
* INCLUDED IN SEGMENT LENGTH, ROOM IS LEFT.
*----- *)
PROCEDURE MERGES(VAR CUR,FWD:TBLPTR);
VAR SBRTAIL,JMPTAIL:TBLPTR; DELETE:INTEGER;
    BEGIN
        IF FWD@.F_JUMPLIST<>NIL THEN
            IF CUR@.F_JUMPLIST<>NIL THEN
                BEGIN
                    JMPTAIL:=CUR@.F_JUMPLIST;
                    WHILE JMPTAIL@.NEXT_FJUMP<>NIL DO
                        JMPTAIL:=JMPTAIL@.NEXT_FJUMP;
                    JMPTAIL@.NEXT_FJUMP:=FWD@.F_JUMPLIST;
                END
            ELSE
                CUR@.F_JUMPLIST:=FWD@.F_JUMPLIST;
            IF FWD@.SBRLIST<>NIL THEN
                IF CUR@.SERLIST<>NIL THEN
                    BEGIN
                        SBRTAIL:=CUR@.SERLIST;
                        WHILE SBRTAIL@.NEXT_SBR<>NIL DO
                            SBRTAIL:=SBRTAIL@.NEXT_SBR;
                        SBRTAIL@.NEXT_SBR:=FWD@.SBRLIST;
                    END
                ELSE
                    CUR@.SERLIST:=FWD@.SBRLIST;
                CUR@.STOP_ACDR:=FWD@.STOP_ADDR;
                CUR@.TABLELIST:=FWD@.TABLELIST;
                CUR@.NEST:=0;
                DELETE:=0;
                CLEAN(CUR,DELETE);
                CUR@.LENGTH:=CUR@.LENGTH+FWD@.LENGTH-((DELETE+1)*
                    FJUMFCONST);
                DISPOSE(FWD,TABLE);
            END;
        END;
    (*----- (* MERGES *) -----*)

```

```

(*-----*)
(* MOD SUMTOTFJMP: SIMULATES THE COMBINE OF THE FJUMP*)
(* THIS IS A PREDICTION IF ADJACENT SEGMENTS WERE *)
(* MERGED. *)
(*-----*)
PROCEDURE MOD_SUMTOTFJMP (CUR, FWD: TBLPTR;
                          VAR DELSUMFJMP: INTEGER);
  VAR P, T: TBLPTR; FCOUNT: INTEGER;
  BEGIN
    FCOUNT:=0;
    P:=CUR@.F JUMPLIST;
    WHILE P<>NIL DO
      BEGIN
        IF P@.JUMP_ADDRTO<=FWD@.STOP_ADDR THEN
          FCOUNT:=FCOUNT+1;
          T:=FWD@.F JUMPLIST;
          WHILE T<>NIL DO
            BEGIN
              IF P@.JUMP_ADDRTO = T@.JUMP_ADDRTO THEN
                FCOUNT:=FCOUNT+1;
                T:=T@.NEXT_FJUMP;
              END;
            END;
          P:=P@.NEXT_FJUMP;
        END;
      DELSUMFJMP:=(FCOUNT+1)*FJUMPCONST;
    END;
  (*-----*)
  (* MOD_SUMTOTFJMP *)
  (*-----*)

```

```

(*-----*)
(* MOD SUMTOTSER: SIMULATES THE CHANGE TO TOTAL LEN *)
(* BECAUSE OF THE MERGING OF ADJACENT SEGMENTS. *)
(*-----*)
PROCEDURE MOD_SUMTOTSBR (CUR, FWD: TBLPTR;
                          VAR DELSUMSBRZ: INTEGER);
  VAR P, T, PP, TT: TBLPTR; SCOUNT: INTEGER;
  BEGIN
    SCOUNT:=0;
    P:=CUR@.SBRLIST;
    WHILE P<>NIL DO
      BEGIN
        T:=FWD@.SBRLIST;
        WHILE T<>NIL DO
          BEGIN
            PP:=P@.SBR;
            IF PP@.TAG = SBRBREAK THEN
              BEGIN
                TT:=T@.SER;
                IF TT@.TAG = SBRBREAK THEN
                  IF PP@.SBRZ = TT@.SBRZ THEN
                    SCOUNT:=SCOUNT+1;
                  END;
                T:=T@.NEXT_SBR;
              END;
            END;
            P:=P@.NEXT_SBR;
          END;
        DELSUMSBRZ:=SCOUNT*SBRCONST;
      END;
    (*-----*)
    (* MOD_SUMTOTSER *)
    (*-----*)

```

```

BEGIN
  CUR:=CURSEG;
  DIAGS_NEST1LENGTHCHK (TEMPFILE, CUR, GOOD_SEGMENT);
  SUMTOT:=0;
  SUMSBRZ:=0;
  IF CUR@.SBRLIST<>NIL THEN
    SERSUM (CUR@.SBRLIST, SUMSBRZ);
  SUMCUR:=CUR@.LENGTH+SUMSBRZ;
  WHILE (CUR@.TABLELIST<>NIL) DO

```

```

BEGIN
  FWD:=CUR@.TABLELIST;
  SUMSERF:=0;
  IF FWD@.SBRLIST<>NIL THEN
    SBRSUM(FWD@.SBRLIST,SUMSBRF);
    SUMPWD:=FWD@.LENGTH+SUMSBRF;
    SUMTOT:=SUMCUR+SUMPWD;
    MOD-SUMTOTFJMP(CUR,FWD,DELSUMFJMP);
    MOD-SUMTCTISBR(CUR,FWD,DELSUMSBRZ);
    SUMTCT:=SUMTOT-DELSUMFJMP-DELSUMSBRZ;
    IF SUMTOT<=LIMIT THEN
      BEGIN
        MERGES(CUR,FWD);
        CUR@.LENGTH:=CUR@.LENGTH-DELSUMSBRZ;
        SUMCUR:=SUMTOT;
      END
    ELSE
      BEGIN
        RESET_INCLUDED(FWD@.SBRLIST);
        RESET_INCLUDED(CUR@.SBRLIST);
        CUR:=FWD;
        DIAGS_NEST1LENGTHCHK(TEMPFILE,CUR,
                              GOOD_SEGMENT);
        SUMCUR:=CUR@.LENGTH;
        SUMSERC:=0;
        SBRSUM(CUR@.SBRLIST,SUMSBRZ);
        SUMCUR:=SUMCUR+SUMSBRZ;
      END;
    SUMTOT:=0;
  END;
  IF CUR@.SBRLIST<>NIL THEN
    RESET_INCLUDED(CUR@.SBRLIST); (*ALL INCLUDES PST*)
    CURSEG@.COALESCED:=TRUE;
  END;
  (*----- (* COMBINE *)-----*)
  BEGIN
    CURSEG:=SBR;
    SBRINVNEST:=SBRINVNEST+1;
    DIAGS_NEST6SBRINVCHK(TEMPFILE,CURSEG,GOOD_SEGMENT,
                          SBRINVNEST);
    CHK_SEGSIZE(CURSEG,LIMIT);
    CCMEINE(CURSEG,LIMIT);
    SBRINVNEST:=SBRINVNEST-1;
  END;
  (*===== (* COALESCE *)=====*)

```



```

*=====*)
* INSTRUCTIONS: PRINTS OUT THE SEGMENTED CODE TOGETHER *
* WITH OTHER INFORMATION TO USE THE SEGMENTED CODE. *
* DOES THIS BY FIRST CHECKING THE SEGMENT TABLE AND *
* ASSIGNING A MEMORY MODULE NUMBER TO SPECIFIC LOCATIONS *
* IN THE TABLE WHERE SBR BREAKS OCCUR. AN *
* IMPLIED ASSIGNMENT IS MADE TO THE FIRST SBR AS A *
* START POINT. OTHERS ARE INCLUDED IF THERE IS A BREAK *
* LEADING TO IT. ONCE MODULES ARE ASSIGNED THEN CODE *
* WITH PROMPTS ADDED ARE COPIED FROM THE ORIGINALS. *
* THESE COPIES ARE THEN PRINTED OUT. *
*=====*)
PROCEDURE INSTRUCTIONS (VAR OUTFILE, MESSAGEFILE,
                        TEMPPFILE:TEXT; BUILT_CODE:CODEPTR;
                        SEG_TBL:TBLPTR; PART_NUM:INTEGER;
                        PARTITION:REAL; GOOD_SEGMENT:BOOLEAN);
VAR HEAD_MODULE:TBLPTR;

*-----*)
* BLD_MODULE_NODES: BUILDS THE MODULES BASED ON THE *
* BREAKS ENCOUNTERED IN THE SEG_TBL *
*-----*)
PROCEDURE BLD_MODULE_NODES (SEG_TBL:TBLPTR;
                           VAR HEAD_MODULE:TBLPTR);
VAR TAIL_MODULE, SEG:TBLPTR;
    MEMCOUNT:INTEGER;

*-----*)
* INSERT_MODULE_NODES: INSERTS A MODULE_NODE INTO *
* THE MODULE_LIST. IT MUST FIRST CHECK TO SEE *
* THAT IT IS NOT ALREADY ACCOUNTED FOR AS THERE *
* MAY BE MULTIPLE INVOKES OF THE SAME BREAK. *
*-----*)
PROCEDURE INSERT_MODULE_NODES (SEG_TBL:TBLPTR;
                              VAR HEAD_MODULE, TAIL_MODULE:TBLPTR;
                              VAR MEMCOUNT:INTEGER);
VAR INSERT:TBLPTR;

*-----*)
* NOT_IN_MODULE_LIST: CHECKS TO SEE IF IN LIST *
*-----*)
FUNCTION NOT_IN_MODULE_LIST (SEG_TBL, HEAD_MODULE
                             :TBLPTR)
                             :BOOLEAN;

VAR S:TBLPTR;
BEGIN
    S:=HEAD_MODULE;
    NOT_IN_MODULE_LIST:=TRUE;
    IF S<>NIL THEN
        BEGIN
            WHILE S<>NIL DO
                BEGIN
                    IF S@.SEG_TBLS = SEG_TBL THEN
                        NOT_IN_MODULE_LIST:=FALSE;
                    S:=S@.NEXT;
                END;
            END;
        END;
    END;
END;
(* NOT_IN_MODULE_LIST *)
BEGIN
    IF NOT_IN_MODULE_LIST (SEG_TBL, HEAD_MODULE) THEN
        BEGIN
            NEW (INSERT, MODULE);
            INSERT@.TAG:=MODULE;
            INSERT@.MEMNUM:=MEMCOUNT;
            MEMCOUNT:=MEMCOUNT+1;
            INSERT@.OFFSET:=-SEG_TBL@.START_ADDR;
            INSERT@.HIGH_OFFSET:=- (SEG_TBL@.START_ADDR DIV
                                   100);
        END;
    END;
END;

```

```

INSERT@.LOWOFFSET:=- (SEG@.START_ADDR
                      100*INSERT@.HIGHOFFSET);
INSERT@.RETURNCODE_NEEDED:=FALSE;
INSERT@.CODELIST:=NIL;
INSERT@.SEG@:=SEG@;
INSERT@.NEXT:=NIL;
IF HEAD_MEMODULE = NIL THEN
  HEAD_MEMODULE:=INSERT;
IF TAIL_MEMODULE <> NIL THEN
  TAIL_MEMODULE@.NEXT:=INSERT;
TAIL_MEMODULE:=INSERT;
END;
END; (* INSERT_MEMODULENODES *)
(*-----*)

(*-----*)
(* RECURSE_BLD_MEMODULENODES: IS THE RECURSIVE PART *)
(* OF THE BLD_MEMODULE ROUTINE. GOES WITHIN THE *)
(* THE SBRFIRST THEN DOWN THE SBRLIST. *)
(* RECURSION IS USED TO TRAVERSE THE SEG@. *)
(*-----*)
PROCEDURE RECURSE_BLD_MEMODULE (SEG:TBLPTR;
                                VAR HEAD_MEMODULE, TAIL_MEMODULE:TBLPTR;
                                VAR MEMCOUNT:INTEGER);
VAR SEG_SBR, SER, SBRL:TBLPTR;
BEGIN
  SEG_SBR:=SEG@.TABLELIST;
  WHILE SEG_SBR<>NIL DO
    BEGIN
      INSERT_MEMODULENODES (SEG_SBR, HEAD_MEMODULE,
                           TAIL_MEMODULE, MEMCOUNT);
      SEG_SBR:=SEG_SBR@.TABLELIST;
    END;
  SEG_SBR:=SEG;
  WHILE SEG_SBR<>NIL DO
    BEGIN
      IF SEG_SBR@.SBRLIST<>NIL THEN
        BEGIN
          SBRL:=SEG_SBR@.SBRLIST;
          WHILE SBRL<>NIL DO
            BEGIN
              SBR:=SBRL@.SBR;
              IF SBR@.TAG = SBRBREAK THEN
                INSERT_MEMODULENODES (SBR@.SERZ,
                                       HEAD_MEMODULE, TAIL_MEMODULE, MEMCOUNT);
              IF SBR@.TAG = SBRBREAK THEN
                SBR:=SBR@.SBRZ;
              RECURSE_BLD_MEMODULE (SBR,
                                     HEAD_MEMODULE, TAIL_MEMODULE,
                                     MEMCOUNT);
              SBRL:=SBRL@.NEXT_SBR;
            END;
          END;
          SEG_SBR:=SEG_SBR@.TABLELIST;
        END;
      END;
    END;
  END; (* RECURSE_BLD_MEMODULE *)
(*-----*)
BEGIN
  HEAD_MEMODULE:=NIL;
  TAIL_MEMODULE:=NIL;
  MEMCOUNT:=1;
  SEG:=SEG@;
  INSERT_MEMODULENODES (SEG, HEAD_MEMODULE,
                       TAIL_MEMODULE, MEMCOUNT);
  RECURSE_BLD_MEMODULE (SEG, HEAD_MEMODULE,
                        TAIL_MEMODULE, MEMCOUNT);
END; (* BLD_MEMODULENODES *)
(*-----*)

```

```

{ *-----* }
{ * BLD MEMODULECODE: THIS ROUTINE BUILDS THE CODE FOR * }
{ * THE MEMORY MODULE CODELISTS. IT WILL JUSTIFY ALL * }
{ * THE ADDRESSES AND ADD BREAK CODE. * }
{ *-----* }
PROCEDURE BLD_MEMODULECODE (BUILT_CODE: CODEPTR;
                           VAR HEAD_MEMODULE: TBLPTR);
  VAR CURMEM: TBLPTR;

  { *-----* }
  { * BLD A MEMORY: THIS ROUTINE INITIALIZES THE * }
  { * RECURSIVE PROCESS THAT WILL BE DONE IN THE * }
  { * FORM_MEMORY ROUTINE. * }
  { *-----* }
  PROCEDURE BLD_A_MEMORY (BUILT_CODE: CODEPTR;
                        VAR HEAD_MEMODULE: TBLPTR;
                        CURMEM: TBLPTR);

    VAR ADDRESS: INTEGER;
        CODE_H, CODE_T: TBLPTR;
        SEG: TBLPTR;

    { *-----* }
    { * FORM_MEMORY: BUILDS A COMPLETE MEMORY MODULE * }
    { * CODE COMPLETE WITH BREAK CODE AND JUSTIFY. * }
    { * ROUTINE RECURSES ON EACH SBR ON THE SBRLISTS * }
    { *-----* }
    PROCEDURE FORM_MEMORY (BUILT_CODE: CODEPTR;
                        VAR HEAD_MEMODULE, CURMEM, SEG, CODE_H,
                        CODE_T: TBLPTR;
                        VAR ADDRESS: INTEGER);
      VAR CODE_HH, CODE_TT, SER, SBRL: TBLPTR;

      { *-----* }
      { * PROCESS SEG: TAKES CARE OF ONE SEGMENT IN THE * }
      { * SEGMENT TABLE'S WORTH OF CODE. * }
      { *-----* }
      PROCEDURE PROCESS_SEG (BUILT_CODE: CODEPTR;
                          VAR HEAD_MEMODULE,
                          CURMEM: TBLPTR; SEG: TBLPTR;
                          VAR CODE_HH, CODE_TT: TBLPTR;
                          VAR ADDRESS: INTEGER);

        VAR START, STOP: INTEGER;

        { *-----* }
        { * COPY_CODE: COPIES CODE FROM A START TO A * }
        { * STOP POINT OF THE BUILT_CODE * }
        { *-----* }
        PROCEDURE COPYCODE (BUILT_CODE: CODEPTR;
                          VAR CODE_HH, CODE_TT: TBLPTR;
                          VAR ADDRESS: INTEGER;
                          START, STOP: INTEGER);

          VAR INSERT, CURTP: TBLPTR;
              CURCP: CODEPTR;
          BEGIN
            CURCP := BUILT_CODE;
            WHILE CURCP@.AADDR <> START DO
              CURCP := CURCP@.SEQ;
              NEW (INSERT, CODE);
              INSERT@.TAG := CODE;
              INSERT@.ABS_ADDR := CURCP@.AADDR;
              INSERT@.ADDRESS := ADDRESS;
              ADDRESS := ADDRESS + 1;
              INSERT@.KEY CODE := CURCP@.KEY;
              INSERT@.SEQUENTIAL := NIL;
              CODE_HH := INSERT;
              CODE_TT := INSERT;
            REPEAT
              NEW (INSERT, CODE);
              CURCP := CURCP@.SEQ;
              INSERT@.TAG := CODE;

```

```

INSERT@.ABS ADDR:=CURCP@.AADDR;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=CURCP@.KEY;
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;
UNTIL (CURCP@.AADDR = STOP);
CODE_TT@.SEQUENTIAL:=NIL;
END; (* COPYCODE *)

```

```

(*-----*)
{* ADD RETURN CODE: ADDS SBR RETURN CODE TO THE *}
{* TAIL SEGMENT OF THE INVOKED SBR *}
(*-----*)
PROCEDURE ADD_RETURN CODE (VAR CODE TT:TELPTR;
                          VAR ADDRESSSS:INTEGER);

```

```

VAR INSERT:TBLPTR;
BEGIN
CODE TT@.KEYCODE:=STO; (*INVSBR CHG 2 STC*)
NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=DISPLAYREGSTORE; (*DISPLAY*)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=RCLIND; (* RCL IND *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=MANRTNREG; (* MAN RTN REG *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=OP; (* OP *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=30+MANRTNREG; (*+MANRTNREG*)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=RS; (* RUN/STCP *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

CODE_TT@.SEQUENTIAL:=NIL;
END; (* ADD_RETURNCODE *)
(*-----*)

```

```

(*-----*)
* AEDCODE SBRBRK MARKINVOKED: ADDS CODE FOR A *
* SBRBRK AND WILL MARK THE INVOKED RTN PCR *
* MANUAL RETURN CODE *
*-----*)
PROCEDURE ADDCODE_SBRBRK_MARKINVOKED
(VAR HEAD_MEMODULE,CURMEM:TBLPTR;
SEG:TBLPTR; VAR CODE TT:TBLPTR;
VAR ADDRESS:INTEGER);
VAR INSERT,SBR,SBRL:TBLPTR;
LBIADDR:INTEGER;

```

```

(*-----*)
* FIND_LBL: FINDS THE LABEL FROM A GIVEN *
* ADDRESS USING THE BUILT_CODE LIST *
*-----*)
FUNCTION FIND_LBL(BUILT_CODE:CODEPTR;
ADDRESS:INTEGER):INTEGER;
VAR C:CODEPTR;
BEGIN
C:=BUILT_CODE;
WHILE C@.AADDR<>ADDRESS DO
C:=C@.SEQ;
FIND_LBL:=C@.SEQ@.KEY;
END; (* FIND_LBL *)
(*-----*)

```

```

(*-----*)
* GENCCDESBR: ADDS CODE FOR A BRK SBR INVOK *
*-----*)
PROCEDURE GENCODESBR(VAR CODE TT:TBLPTR;
HEAD_MEMODULE,CURMEM,SBR:TBLPTR;
VAR ADDRESS:INTEGER);
VAR RELADDR,HUNDREDS,TENS,UNITS:INTEGER;
MEMPTR:TBLPTR;
BEGIN
MEMPTR:=HEAD_MEMODULE;
WHILE(MEMPTR@.SEGTBLS<>SBR@.SBRZ) DO
MEMPTR:=MEMPTR@.NEXT;
RELADDR:=SBR@.SBRZ@.START_ADDR+MEMPTR@.
OFFSET;
HUNDREDS:=RELADDR DIV 100;
TENS:=(RELADDR-(HUNDREDS*100)) DIV 10;
UNITS:=RELADDR-(HUNDREDS*100+TENS*10);
NEW(INSERT,CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESS;
ADDRESS:=ADDRESS+1;
INSERT@.KEYCODE:=LBL; (* LBL *)
IF CODE TT <> NIL THEN
CODE TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;
NEW(INSERT,CCODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESS;
ADDRESS:=ADDRESS+1;
INSERT@.KEYCODE:=FIND_LBL (* KEY LBL *)
(BUILT_CODE,SBR@.SBRZ@.START_ADDR);
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=STO; (* STORE *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=DISPLAYREGSTORE; (* DISP *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=OP; (* CP *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=20+MANRTNREG; (* INCR *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=CURMEM@.MEMNUM; (* MEM# *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=STOIND; (* IND STO *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=MANRTNREG; (* MAN RTN *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=MEMPTR@.MEMNUM; (* MEM #*)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;

```

```

INSERT@.KEYCODE:=DECIMAL; (* . *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=HUNDREDS; (* HUNDRECS *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=TENS; (* TENS *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=UNITS; (* UNITS *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=RS; (* R/S *)
INSERT@.SEQUENTIAL:=NIL;
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

END; (* GENCODE_SER *)
(*-----*)
BEGIN
IF SEG@.SBRLIST<>NIL THEN
BEGIN
SERL:=SEG@.SBRLIST;
WHILE SBRL<>NIL DO
BEGIN
SBR:=SBRL@.SBR;
IF SBR@.TAG=SBRBREAK THEN
GENCODESBR (CODE_TT, HEAD, MEMODULE,
CURMEM, SBR, ADDRESSSS);
SBRL:=SBRL@.NEXT_SBR;
END;
END;
END; (* ADDCODE_SBRBRK_MARKINVOKED *)
(*-----*)

```

```

(*-----*)
{* JUSTIFY_CODE: SETS ALL THE JUMPS AND ADDRS *}
{*-----*}
PROCEDURE JUSTIFY_CODE(SEG, CURMEM:TBLPTR;
                      VAR CODE_HH:TELETR);
  VAR T, F, JMPTR:TBLPTR;
      DELADDR, RELADDR, ABSADDR:INTEGER;

  (*-----*)
  (* ADVANCE_CODE_TBL: ADVANCES THE CODEPTRS *)
  (* OF THE 1, 2, AND 3-STEP INSTRUCTIONS. *)
  (*-----*}
  PROCEDURE ADVANCE_CODE_TBL(VAR F:TBLPTR);
    BEGIN
      IF F@.KEYCODE IN STEP 3 THEN
        F:=F@.SEQUENTIAL@.SEQUENTIAL@.SEQUENTIAL
          @.SEQUENTIAL
      ELSE
        IF F@.KEYCODE IN STEP 2 THEN
          F:=F@.SEQUENTIAL@.SEQUENTIAL@.
            SEQUENTIAL
        ELSE
          IF F@.KEYCODE IN (STEP 1+(.71,76.))
            THEN F:=F@.SEQUENTIAL@.SEQUENTIAL
          ELSE
            F:=F@.SEQUENTIAL;
        END;
      (* ADVANCE_CODE_TBL *)
    END;

  BEGIN
    F:=CODE_HH;
    WHILE F<>NIL DO
      BEGIN
        IF F@.KEYCODE IN (STEP_3+STEP_2) THEN
          BEGIN
            IF F@.KEYCODE IN STEP 3 THEN
              T:=F@.SEQUENTIAL@.SEQUENTIAL
            ELSE
              T:=F@.SEQUENTIAL;
            ABSADDR:=T@.KEYCODE*100;
            ABSADDR:=ABSADDR
              +T@.SEQUENTIAL@.KEYCODE;
            IF SEG@.STOP_ADDR >= ABSADDR THEN
              BEGIN
                DELADDR:=ABSADDR-T@.ABS_ADDR;
                RELADDR:=T@.ADDRESS+DELADDR;
                T@.KEYCODE:=RELADDR DIV 100;
                T@.SEQUENTIAL@.KEYCODE:=
                  RELADDR-(100*T@.KEYCODE);
              END
            ELSE
              BEGIN
                JMPTR:=SEG@.F_JUMPLIST;
                WHILE JMPTR@.JUMP_ADDR TO <>
                  ABSADDR DO
                  JMPTR:=JMPTR@.NEXT_FJUMP;
                  T@.KEYCODE:=JMPTR@.
                    JUMP_INTADDRTC1;
                  T@.SEQUENTIAL@.KEYCODE:=JMPTR@.
                    JUMP_INTADDRTC2;
                END;
              END;
            ADVANCE_CODE_TBL(F);
          END;
        END;
      END;
    (* JUSTIFY *)
  (*-----*)

```



```

*-----*
* ADDCODE_FJMP: ADDS CODE FOR JUMP BREAK *
*-----*
PROCEDURE ADDCODE_FJMP(VAR HEAD MEMODULE,CUFMEM,
                      SEG, CODE TT:TBLPTR,
                      VAR ADDRESS:INTEGER);
VAR CUR,INSERT,JUMPTR,MEMPTR:TBLPTR;
    ADDCRT01,ADDRTO2, MEM_ADDRTO:INTEGER;
    DELTA_ADDRESS:INTEGER;

*-----*
* GEN JUMPCODE SETINTADDRS: GENERATES THE *
* JUMP CODE AND SETS THE INTADDR FIELDS *
* OF THE SEGMENT TABLE F JUMPLIST. *
* INTADDR ARE THE ADDRESSES LOCAL TO THAT *
* SPECIFIC PIECE OF CODE (THE PROMPT). *
*-----*
PROCEDURE GEN_JUMPCODE_SETINTADDRS
(VAR CODE_TT:TBLPTR;-VAR ADDRESS:INTEGER;
 VAR JUMPTR:TBLPTR);
VAR INSERT:TBLPTR;
    HUNDREDS,TENS,UNITS,RELADDRESS:INTEGER;
BEGIN
    HUNDREDS:=JUMPTR@.JUMP_ADDRTO1;
    TENS:=JUMPTR@.JUMP_ADDRTO2 DIV 10;
    UNITS:=JUMPTR@.JUMP_ADDRTO2-(10*TENS);

    NEW(INSERT, CODE);
    INSERT@.TAG:=CODE;
    INSERT@.ADDRESS:=ADDRESS;
    RELADDRESS:=ADDRESS;
    ADDRESS:=ADDRESS+1;
    INSERT@.KEYCODE:=STO; (* STO *)
    IF CODE_TT<>NIL THEN
        CODE_TT@.SEQUENTIAL:=INSERT;
    CODE_TT:=INSERT;

    NEW(INSERT, CODE);
    INSERT@.TAG:=CODE;
    INSERT@.ADDRESS:=ADDRESS;
    ADDRESS:=ADDRESS+1;
    INSERT@.KEYCODE:= DISPLAYREGSTORE; (*DISP*)
    CODE_TT@.SEQUENTIAL:=INSERT;
    CODE_TT:=INSERT;

    NEW(INSERT, CODE);
    INSERT@.TAG:=CODE;
    INSERT@.ADDRESS:=ADDRESS;
    ADDRESS:=ADDRESS+1;
    INSERT@.KEYCODE:=CE; (* CE *)
    CODE_TT@.SEQUENTIAL:=INSERT;
    CODE_TT:=INSERT;

    NEW(INSERT, CODE);
    INSERT@.TAG:=CODE;
    INSERT@.ADDRESS:=ADDRESS;
    ADDRESS:=ADDRESS+1;
    INSERT@.KEYCODE:= JUMPTR@.MEM_ADDR;
    CODE_TT@.SEQUENTIAL:=INSERT;
    CODE_TT:=INSERT;

    NEW(INSERT, CODE);
    INSERT@.TAG:=CODE;
    INSERT@.ADDRESS:=ADDRESS;
    ADDRESS:=ADDRESS+1;
    INSERT@.KEYCODE:= DECIMAL; (*.**)
    CODE_TT@.SEQUENTIAL:=INSERT;
    CODE_TT:=INSERT;

```

```

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=HUNDREDS; (* 100S *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=TENS; (* 10S *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=UNITS; (* 1S *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;

NEW (INSERT, CODE);
INSERT@.TAG:=CODE;
INSERT@.ADDRESS:=ADDRESSSS;
ADDRESSSS:=ADDRESSSS+1;
INSERT@.KEYCODE:=RS; (* R/S *)
CODE_TT@.SEQUENTIAL:=INSERT;
CODE_TT:=INSERT;
JUMPTR@.JUMP_INTADDRTO1:=RELADDRESS
DIV 100;
JUMPTR@.JUMP_INTADDRTO2:=RELADDRESS-100*
JUMPTR@.JUMP_INTADDRTO1;
END; (* GEN_JUMP CODE_SETINTADDR *)
(*-----*)
BEGIN
IF SEG@.TABLELIST<>NIL THEN
BEGIN
MEMPTR:=HEAD MEMODULE;
WHILE NOT (SEG@.STOP_ADDR+1 = MEMPTR@.
SEGTBLS@.START_ADDR) DO
MEMPTR:=MEMPTR@.NEXT;
ADDRTO1:=MEMPTR@.SEGTBLS@.START_ADDR
DIV 100;
ADDRTO2:=MEMPTR@.SEGTBLS@.START_ADDR
- (100*ADDRTO1);
ADDRTO1:=ADDRTO1+MEMPTR@.HIGHOFFSET;
ADDRTO2:=ADDRTO2+MEMPTR@.LOWOFFSET;
MEM_ADDRTO:=MEMPTR@.MEMNUM;
NEW (JUMPTR, FWD_JUMP);
WITH JUMPTR@ DO
BEGIN
JUMP_ADDRTO1:=ADDRTO1;
JUMP_ADDRTO2:=ADDRTO2;
MEM_ADDR:=MEM_ADDRTO;
END;
GEN_JUMP CODE_SETINTADDR
(CODE_TT, ADDRESSSS, JUMPTR);
DISPOSE (JUMPTR, FWD_JUMP);
END;
IF SEG@.F_JUMLIST<>NIL THEN
BEGIN
JUMPTR:=SEG@.F_JUMLIST;
WHILE JUMPTR<>NIL DO
BEGIN
MEMPTR:=HEAD_MEMODULE;

```

```

WHILE NOT ((JUMPTR@.JUMP_ADDRTO >=
MEMPTR@.SEGTBLS@.START_ADDR) AND
(JUMPTR@.JUMP_ADDRTO <=
MEMPTR@.SEGTBLS@.STOP_ADDR)) DO
MEMPTR:=MEMPTR@.NEXT;

DELTA_ADDRESS:=JUMPTR@.JUMP_ADDRTO-
MEMPTR@.SEGTBLS@.START_ADDR;
ADDRTO1:=DELTA_ADDRESS DIV 100;
ADDRTO2:=DELTA_ADDRESS-100*ADDRTO1;
MEM_ADDRTO:=MEMPTR@.MEMNUM;
JUMPTR@.JUMP_ADDRTO1:=ADDRTO1;
JUMPTR@.JUMP_ADDRTO2:=ADDRTO2;
JUMPTR@.MEM_ADDR:=MEM_ADDRTO;
GEN JUMPCODE SETINTADDRS
(CODE TT,ADDRESS, JUMPTR);
JUMPTR:=JUMPTR@.NEXT_FJUMP;
END;
ENC;
END; (* ADDCODE_FJMP *)
(*-----*)
BEGIN
START:=SEG@.START_ADDR;
STOP:=SEG@.STOP_ADDR;
COPYCODE(BUILT_CODE, CODE_HH, CODE_TT,
ADDRESS, START, STOP);
IF (CURMEM@.RETURNCODE_NEEDED) AND
(SEG=CURMEM@.SEGTBLS) THEN
ADD_RETURNCODE(CODE_TT, ADDRESS);
ADDCODE_FJMP(HEAD_MODULE, CURMEM, SEG, CODE_TT,
ADDRESS);
ADDCODE_SBRBK_MARKINVOKED(HEAD_MODULE,
CURMEM, SEG, CODE_TT, ADDRESS);
JUSTIFY_CODE(SEG, CURMEM, CODE_HH);
END; (* PROCESS_SEG *)
(*-----*)

```

```

(*-----*)
* MARKINVKCKED: MARKS MEMODULE OF SBR WHICH WAS *
* INVOLVED IN A BREAK FOR A MANUAL RETURN. *
* ACD_RETRUNCODE USES THIS MARK TO ADD CCDE. *
(*-----*)
PROCEDURE MARKINVOKED(VAR HEAD_MEMODULE:TBLPTR;
                      SBR:TBLPTR);
  VAR MEMPT5, SBR_BRKINCLUDE:TBLPTR;
  BEGIN
    SBR_BRKINCLUDE:=SBR@.SBRZ;
    MEMPTR:=HEAD_MEMODULE;
    WHILE SER_BRKINCLUDE@.TABLELIST<>NIL DO
      SBR_ERKINCLUDE:=SBR_BRKINCLUDE@.TABLELIST;
    WHILE MEMPTR@.SEGTBLS<>SBR_BRKINCLUDE DO
      MEMPTR:=MEMPTR@.NEXT;
    MEMPTR@.RETURN_CODE_NEEDED:=TRUE;
  END;
  (*-----*)
  (* MARKINVOKED *)
  BEGIN
    IF (SEG@.INCLUDED = FALSE) THEN
      BEGIN
        PROCESS_SEG (BUILT_CODE, HEAD_MEMODULE, CURMEM,
                    SEG, CODE_HH, CODE_TT, ADDRESSSS);
        IF CCDE_H=NIL THEN
          CODE_H:=CODE_HH;
        IF CCDE_T<>NIL THEN
          CODE_T@.SEQUENTIAL:=CODE_HH;
        SEG@.INCLUDED:=TRUE;
        CODE_T:=CODE_TT;
        IF (SEG@.SBRLIST<>NIL) THEN
          BEGIN
            SBRL:=SEG@.SBRLIST;
            WHILE SBRL<>NIL DO
              BEGIN
                SBR:=SBRL@.SBR;
                CASE SBR@.TAG OF
                  SBRBREAK:
                    MARKINVOKED (HEAD_MEMODULE, SBR);
                  TABLE:
                    FORM_MEMORY (BUILT_CODE,
                                HEAD_MEMODULE, CURMEM, SBR, CODE_H,
                                CODE_T, ADDRESSSS);
                END; (* CASE *)
                SBRL:=SBRL@.NEXT_SBR;
              END;
            END;
          END;
        END;
      END;
    END;
  END;
  (*-----*)
  (* FORM_MEMORY *)
  BEGIN
    SEG:=CURMEM@.SEGTBLS;
    ADDRESSSS:=0; CODE_H:=NIL; CODE_T:=NIL;
    FORM_MEMORY (BUILT_CODE, HEAD_MEMODULE, CURMEM, SEG,
                CODE_H, CODE_T, ADDRESSSS);
    CURMEM@.CODELIST:=CCDE_H;
  END;
  (*-----*)
  (* BLD_A_MEMORY *)
  BEGIN
    CURMEM:=HEAD_MEMODULE;
    WHILE CURMEM<>NIL DO
      BEGIN
        BLD_A_MEMORY (BUILT_CODE, HEAD_MEMODULE, CURMEM);
        CURMEM@.SEGTBLS@.INCLUDED:=FALSE;
        IF CURMEM@.SEGTBLS@.SBRLIST<>NIL THEN
          RESET_INCLUDED (CURMEM@.SEGTBLS@.SBRLIST);
        CURMEM:=CURMEM@.NEXT;
      END;
    END;
  END;
  (*-----*)
  (* BLD_MEMODULECODE *)
  (*-----*)

```

```

(*-----*)
PROCEDURE OUTPUT_INSTR(VAR OUTFILE,MESSAGEFILE,
                      TEMPFILE:TEXT; BUILT_CODE:CODEPTR;
                      HEAD_MEMODULE:TBLPTR; PART_NUM:INTEGER;
                      PARTITION:REAL;
                      GOOD_SEGMENT:BCOLEAN);

{ * CUTPUT_GOODSEGS: PRINTS OUT GOOD SEGMENT INSTRUC * }
{ *-----* }
PROCEDURE CUTPUT_GOODSEG(VAR OUTFILE,MESSAGEFILE:
                        TEXT; HEAD_MEMODULE:TBLPTR; PART_NUM:INTEGER;
                        PARTITION:REAL);

{ * CUTPUT_MSGF1: PRINTS OUT GOOD SEGMENT GENERAL * }
{ * INSTRUCTIONS AND THE PROGRAM LISTING BY MEM * }
{ * MODULES. * }
{ *-----* }
PROCEDURE OUTPUT_MSGF1(VAR OUTFILE,MESSAGEFILE:TEXT;
                      HEAD_MEMODULE:TBLPTR; PARTNUM:INTEGER;
                      PARTITION:REAL);

VAR P,T:TBLPTR;
    STEP_TYPE:INTEGER;

{ * SET_STEP: SETS THE KEY CODE COUNTER FOR THE * }
{ * PRINTLINE PROCEDURE. USED FOR CORRECT OUT- * }
{ * PUT OF THE TI-59 CODE (3,2,1,0 STEP INSTR) * }
{ *-----* }
{ * 3 STEP      2 STEP      71&76      1      0 * }
{ * AAA          AAA          AAA      AAA      AAA * }
{ * X            XXX          AAA      XX      (.1.) * }
{ * XXX          XXX          (.5,4.)/(.3,2.) * }
{ * XXX          (.8,7,6.) * }
{ * (.12,11,10,9.) * }
{ *-----* }
PROCEDURE SET_STEP(T:TBLPTR;
                  VAR STEP_TYPE:INTEGER);
BEGIN
  IF STEP_TYPE IN (.0,1,3,5,8.) THEN
    IF T@.KEYCODE IN STEP_3 THEN
      STEP_TYPE:=12
    ELSE
      IF T@.KEYCODE IN STEP_2 THEN
        STEP_TYPE:=8
      ELSE
        IF T@.KEYCODE IN STEP_1 THEN
          STEP_TYPE:=3
        ELSE
          IF T@.KEYCODE IN (.71,76.) THEN
            STEP_TYPE:=5
          ELSE
            STEP_TYPE:=1;
  END;
(*-----* (* SET_STEP *)

```

```

(*-----*)
{* PRINT_LINE: PRINTS OUT ONE LINE OF CODE AT A *}
{* TIME. *}
(*-----*)
PROCEDURE PRINT_LINE(VAR OUTFILE:TEXT; VAR T:
                    TBLPTR; VAR STEP_TYPE:INTEGER);

(*-----*)
{* WHITELELS: WRITES TBLPTR LBLs IN WHOLE KEY *}
{* CODE FORMAT IE.  NNN NN  AAA *}
(*-----*)
PROCEDURE WRITELBLS(VAR OUTFILE:TEXT; T:TBLPTR);
BEGIN
  WRITE(OUTFILE, ' ');
  WRITE LEADZER0(OUTFILE, T@.ADDRESS, 3);
  WRITE(OUTFILE, ' ');
  WRITE LEADZER0(OUTFILE, T@.KEYCODE, 2);
  WRITE(OUTFILE, ' ');
  WRITELBL(OUTFILE, T@.KEYCODE);
END; (* WRITELBLS *)
(*-----*)

(*-----*)
{* WRITENUMS: WRITES OUT A LINE OF DIGITS OF *}
{* TI-59 CODE (KEYCODES ARE DIGITS NOT LBLs) *}
(*-----*)
PROCEDURE WRITENUMS(VAR OUTFILE:TEXT; T:TBLPTR);
BEGIN
  WRITE(OUTFILE, ' ');
  WRITE LEADZER0(OUTFILE, T@.ADDRESS, 3);
  WRITE(OUTFILE, ' ');
  WRITE LEADZER0(OUTFILE, T@.KEYCODE, 2);
  WRITE(OUTFILE, ' ');
  WRITE LEADZER0(OUTFILE, T@.KEYCODE, 2);
  WRITELN(OUTFILE);
END; (* WRITENUMS *)
(*-----*)
BEGIN
  IF STEP_TYPE IN (.1, 3, 4, 5, 8, 12.) THEN
    WRITELBLS(OUTFILE, T)
  ELSE
    WRITENUMS(OUTFILE, T);
  T:=T@.SEQUENTIAL;
  STEP_TYPE:=STEP_TYPE-1;
END; (* PRINT_LINE *)
(*-----*)
BEGIN
  PRINTLN MSG(OUTFILE, MESSAGEFILE, BAXINSTR);
  PRINT MSGLINE1(OUTFILE, MESSAGEFILE, RTNRGTOP);
  WRITELN(CUTFILE, REGCOUNT:3);
  PRINT MSGLINE1(OUTFILE, MESSAGEFILE, STOINRG);
  WRITELN(CUTFILE, MANRTNREG:3);
  WRITELN(CUTFILE);
  PRINT MSGLINE1(OUTFILE, MESSAGEFILE, PGMPARTIS);
  WRITELN(CUTFILE, PARTITION:4:2);
  WRITELN(CUTFILE);
  PRINT MSGLINE1(OUTFILE, MESSAGEFILE, PARTNUMIS);
  WRITELN(CUTFILE, PARTNUM:1);
  WRITELN(CUTFILE); WRITELN(OUTFILE);
  P:=HEAD MEMODULE;
  WHILE P<>NIL DO
    BEGIN
      T:=P@.CODELIST;
      STEP_TYPE:=0;
      WRITELN(OUTFILE); WRITELN(OUTFILE);
      WRITELN(OUTFILE);
      PRINT_MSGLINE1(OUTFILE, MESSAGEFILE, MODN);
    END
  END

```

```

WRITEIN(OUTFILE,P@.MEMNUM:1);
PRINTLN_MSG(OUTFILE,MESSAGEFILE,CARD1);
PRINTLN_MSG(OUTFILE,MESSAGEFILE,SIDE1);
WHILE T<>NIL DO
  BEGIN
    IF T@.ADDRESS=240 THEN
      BEGIN
        WRITELN(OUTFILE);WRITELN(OUTFILE);
        PRINTLN_MSG(OUTFILE,MESSAGEFILE,SIDE2);
      END;
    IF T@.ADDRESS = 480 THEN
      BEGIN
        WRITELN(OUTFILE);WRITELN(OUTFILE);
        PRINTLN_MSG(OUTFILE,MESSAGEFILE,CARD2);
        PRINTLN_MSG(OUTFILE,MESSAGEFILE,SIDE1);
      END;
    SET_STEP(T,STEP_TYPE);
    PRINT_LINE(OUTFILE,T,STEP_TYPE);
  END;
  P:=P@.NEXT;
END;
END;
(* OUTPUT_MSGF1 *)

```

```

(*-----*)
(* OUTPUT MSGF2: OUTPUTS SPECIFIC PROMPTS AND *)
(* SPECIAL PROGRAM INSTRUCTIONS *)
(*-----*)
PROCEDURE OUTPUT_MSGF2 (VAR OUTFILE,MESSAGEFILE:TEXT;
                        HEAD_MODULE:TBLPTR);
  VAR SBRL,SEP,SM,P,F,SEG:TBLPTR;
  IS_SBRERK:BOOLEAN;
  BEGIN
    P:=HEAD_MODULE;
    PRINTLN_MSG(OUTFILE,MESSAGEFILE,SPECIFICS);
    WHILE P<>NIL DO
      BEGIN
        SEG:=P@.SEG:TBL;
        F:=SEG@.F:JUMPLIST;
        SBRL:=SEG@.SBRLIST;
        WRITEIN(OUTFILE);
        PRINT_MSGLINE1(OUTFILE,MESSAGEFILE,MODPROMPTS);
        WRITEIN(OUTFILE,P@.MEMNUM:1);
        PRINTLN_MSG(OUTFILE,MESSAGEFILE,PFWDJ);
        IF P=NIL THEN
          PRINTLN_MSG(OUTFILE,MESSAGEFILE,NONE)
        ELSE
          BEGIN
            WHILE F<>NIL DO
              BEGIN
                PRINT_MSGLINE1(OUTFILE,MESSAGEFILE,
                              ASTER);
                WRITE(OUTFILE,P@.MEM_ADDR:1,'. ');
                WRITE_LEADZERO(OUTFILE,(P@.
                                         JUMP_ADDRTO1*100
                                         P@.JUMP_ADDRTO2),3);
                WRITELN(OUTFILE);
                P:=P@.NEXT_PJUMP;
              END;
            END;
          END;
        PRINTLN_MSG(OUTFILE,MESSAGEFILE,PSBRINV);
        IF SERI=NIL THEN
          PRINTLN_MSG(OUTFILE,MESSAGEFILE,NONE)
        END;
      END;
    END;
  END;

```

```

ELSE
  BEGIN
    IS_SBRBRK:=FALSE;
    WHILE SBRL<>NIL DO
      BEGIN
        SBR:=SFRL@.SBR;
        IF SBR@.TAG = SBRBREAK THEN
          BEGIN
            SM:=HEAD MEMODULE;
            SBR:=SBR@.SBRZ;
            WHILE SM@.SEGTBLS<>SBR DO
              SM:=SM@.NEXT;
            IS_SBRBRK:=TRUE;
            PRINT_MSGLINE1(OUTFILE,MESSAGEFILE,ASTER);
            WRITELN(OUTFILE,SM@.MEMNUM:1,
              '.000');
          END;
          SBRL:=SBRL@.NEXT_SBR;
        END;
        IF NOT IS_SBRBRK THEN
          PRINTLN_MSG(OUTFILE,MESSAGEFILE,NCNE)
        END;
        PRINTLN_MSG(OUTFILE,MESSAGEFILE,PMANRTN);
        IF P@.RETURN_CODE_NEEDED THEN
          PRINTLN_MSG(OUTFILE,MESSAGEFILE,YES)
        ELSE
          PRINTLN_MSG(OUTFILE,MESSAGEFILE,NONE);
          PRINTLN_MSG(OUTFILE,MESSAGEFILE,PSEQ);
          IF SEG@.TABLELIST <> NIL THEN
            BEGIN
              PRINT_MSGLINE1(OUTFILE,MESSAGEFILE,
                ASTER);
              WRITELN(OUTFILE,P@.NEXT@.
                MEMNUM:1, '.000')
            END
          ELSE
            PRINTLN_MSG(OUTFILE,MESSAGEFILE,NONE);
            P:=P@.NEXT;
          END;
          WRITELN(OUTFILE); WRITELN(OUTFILE);
          PRINTLN_MSG(OUTFILE,SCRATCH,DATAREAD);
          WRITELN(OUTFILE); WRITELN(OUTFILE);
          PRINTLN_MSG(OUTFILE,SCRATCH,REGMAP);
          WRITELN(OUTFILE); WRITELN(OUTFILE);
        END;
      END;
    END;
  END;
  (*-----(* OUTPUT_MSG2 *)-----*)
  BEGIN
    OUTPUT_MSGF1(OUTFILE,MESSAGEFILE,HEAD MEMODULE,
      PART_NUM,PARTITION);
    OUTPUT_MSGF2(OUTFILE,MESSAGEFILE,HEAD MEMODULE);
    WRITELN(OUTFILE);
    PRINTLN_MSG(OUTFILE,MESSAGEFILE,ENDBL);
  END;
  (*-----(* OUTPUT_GOODSEG *)-----*)

```



```

(*-----*)
{* OUTPUT_BADSEG: HANDLES BAD SEG INSTRUCTIONS *}
(*-----*)
PROCEDURE OUTPUT_BADSEG (VAR OUTFILE, MESSAGEFILE,
                        TEMPFILE:TEXT; BUILT_CODE:CODEPTR);
  BEGIN
    WRITELN (OUTFILE) : WRITELN (OUTFILE);
    PRINTLN MSG (OUTFILE, MESSAGEFILE, FAILINSTR);
    RESET (TEMPFILE);
    PRINTLN MSG (OUTFILE, TEMPFILE, 9);
    WRITELN (OUTFILE) : WRITELN (OUTFILE);
    PRINTLN MSG (OUTFILE, MESSAGEFILE, UNSEGCODLBL);
    WRITELN (OUTFILE) : WRITELN (OUTFILE);
    PRINT CODELIST (OUTFILE, BUILT_CODE);
    WRITELN (OUTFILE);
    PRINTLN_MSG (OUTFILE, MESSAGEFILE, ENDLBL);
  END;
(*-----*)
(* OUTPUT_BADSEG *)
BEGIN
  IF NOT GOOD_SEGMENT THEN
    OUTPUT_BADSEG (OUTFILE, MESSAGEFILE, TEMPFILE,
                  BUILT_CODE)
  ELSE
    OUTPUT_GOODSEG (OUTFILE, MESSAGEFILE, HEAD_MEMODULE,
                   PART_NUM, PARTITION);
  END;
(*-----*)
(* OUTPUT_INSTR *)
BEGIN
  BLD_MEMODULENODES (SEGTBL, HEAD_MEMODULE);
  BLD_MEMODULECODE (BUILT_CODE, HEAD_MEMODULE);
  OUTPUT_INSTR (OUTFILE, MESSAGEFILE, TEMPFILE, BUILT_CODE,
               HEAD_MEMODULE, PART_NUM, PARTITION,
               GOOD_SEGMENT);
  END;
(*-----*)
(* INSTRUCTIONS *)
(*-----*)

```

```
(*****  
{*                               MAIN DRIVER                               *}  
*****)
```

```
BEGIN  
  INIT SETS(TEMPFILE,STEP 0,STEP 1,STEP 2,STEP 3,  
            GOOD_SEGMENT,MESSAGEFILE,TILBL,SBRINVNEST);  
  DET_LIMIT(REGCOUNT,LIMIT,NUMBANKS,PART_NUM,PARTITION);  
  INPUT (SCRATCH,BUILT_CODE,BUILT_CODE_COUNT);  
  SETJMPS (BUILT_CODE);  
  BUILD_SEGTBL (BUILT_CODE,SEGTBL,LIMIT,BUILT_CODE_COUNT);  
  COALESCE (SEGTBL@.TABLELIST,LIMIT,GOOD_SEGMENT  
            SBRINVNEST);  
  WRITEIN (TEMPFILE,'$9'); (* CLOSES TEMPFILE 'DIAG' FILE *)  
  INSTRUCTIONS (OUTFILE,MESSAGEFILE,TEMPFILE,  
                BUILT_CODE,SEGTBL@.TABLELIST,PART_NUM,PARTITION,  
                GOOD_SEGMENT);  
  REWRITE (TEMPFILE);      (* ERASES TEMPFILE DIAG FILE *)  
END.
```

```
(*****  
{*                               END OF PROGRAM                               *}  
*****)
```

APPENDIX J
MESSAGEFILE FILE--LINKER MESSAGES

\$81

=====

EAX59 PROGRAM INSTRUCTIONS: VERSION 1.0

=====

* CONGRATULATIONS, YOU HAVE JUST COMPILED A BASIC PROGRAM INTO A TI-59 PROGRAM. IN SO DOING IT IS VERY POSSIBLE THAT YOUR PROGRAM IS LARGER THAN THE MEMORY OF THE CALCULATOR. IF THIS IS THE CASE THEN THE PROGRAM HAS BEEN SEGMENTED AND PROMPTING CODE INSERTED TO GUIDE YOUR CALCULATOR PROGRAM DURING ITS EXECUTION. THE REMAINDER OF THIS OUTPUT CONSISTS OF TI-59 CODE LISTINGS AND OTHER INFORMATION TO AID YOU IN YOUR PROGRAM EXECUTION.

* THE FOLLOWING DEFINITIONS ARE PROVIDED AS AN AID TO READING THE PROGRAM LISTING FILE.

* DEFINITIONS:

- * MODULE: A MODULE IS DEFINED TO BE ALL THE MEMORY DEDICATED TO PROGRAM STEPS. THE SIZE IS VARIABLE AND IS DEPENDENT ON THE REGISTER REQUIREMENT. VALUES RANGE FROM 0 TO 239, 479 OR 719 DEPENDING ON THE AMOUNT OF REGISTERS USED BY THE PROGRAM.
- * CARD: A CARD IS DEFINED TO BE ONE MAGNETIC CARD. A CARD HOLDS 480 PROGRAM STEPS. THESE STEPS ARE NOT CONTIGUOUS BUT ARE ARRANGED ON THE TWO SIDES OF THE CARD.
- * SIDE: A SIDE IS ONE HALF OF A CARD. IT CONTAINS UP TO 240 STEPS. WHEN ONE SIDE OF A CARD IS READ BY THE CALCULATOR 240 PROGRAM STEPS ARE FILLED IN MEMORY. THESE BLOCKS OF 240 STEPS ARE REFERED TO AS "BANKS" IN THE MANUFACTURER LITERATURE. WHEN LOADING A CARD YOU WILL LOAD ONLY BANK NUMBER 1 AND/OR 2 FOR PROGRAM STEPS.
- * PARTITION: THIS IS DEFINED TO BE THE CURRENT SETTING OF CALCULATOR MEMORY AS APPLIED TO THE AMOUNT OF MEMORY DEDICATED TO STORAGE REGISTERS AND THE AMOUNT DEDICATED TO PROGRAM STEPS. WE WILL BE DEALING WITH 3 PARTITIONS. THESE ARE:

3	719.29
4	479.59
5	239.69

FORMAT: X YYY.ZZ

WHERE X STANDS FOR PARTITION NUMBER
YYY STANDS FOR PROGRAM STEPS (0-YYY)
ZZ STANDS FOR REGISTERS (0-ZZ).

=====

TI-59 PROGRAM LISTING BY MODULE/CARD/SIDE:

=====

- * THE FOLLOWING IS YOUR PROGRAM LISTING. THE PROGRAM IS LISTED ACCORDING TO MODULE NUMBER AND ITS ASSOCIATED CARDS AND CARD SIDES.
- * REFER TO THE TI-59 PROGRAMMER'S GUIDE ON HOW TO INPUT A PROGRAM AND WRITE IT TO MAGNETIC CARDS.
- * CAUTION: ENSURE THAT THE CORRECT CALCULATOR PARTITION IS SET BEFORE INPUTTING A PROGRAM AND WRITING TO MAGNETIC CARDS.
- * CAUTION: ENSURE THAT YOU DO NOT CONFUSE BANK NUMBERS WITH CARD/MODULE OR SIDE NUMBERS. THE NUMBERS WHICH REFER TO THE LISTING ARE AKIN TO A VIRTUAL ADDRESS AND DO NOT REPRESENT THE ACTUAL BANK NUMBER. IF IN DOUBT, REMEMBER TO USE THE TABLE BELOW TO TRANSLATE VIRTUAL TO ACTUAL BANK NUMBERS.

VIRTUAL BANK	ACTUAL BANK
MODULE #	
CARD1	
SIDE1 -----	BANK1
MODULE #	
CARD1	
SIDE2 -----	BANK2
MODULE #	
CARD2	
SIDE1 -----	BANK3

=====

TI-59 LISTING

=====

\$81

=====

 TI-59 PROGRAM SPECIFIC INSTRUCTIONS:

 =====

- * THE FOLLOWING INFORMATION WILL TELL YOU HOW TO RUN YOUR PROGRAM.
- * YOU MUST ENTER YOUR PROGRAM MANUALLY INTO THE CALCULATOR AND WRITE THE PROGRAM TO MAGNETIC CARDS. THIS STEP ONLY NEEDS TO BE ACCOMPLISHED ONCE. AFTER THAT, THE PROGRAM IS ENTERED USING THE MAGNETIC CARD FACILITY OF THE CALCULATOR. SEE THE MANUFACTURER'S LITERATURE ON ENTERING A PROGRAM AND WRITTING IT TO MAGNETIC CARDS. YOU WILL NEED TO PARTITION MEMORY.

- * HOW TO PARTITION THE MEMORY
 - * KEY SEQUENCE:

X
2ND
OP
17

- * X IS THE PARTITION NUMBER GIVEN IN THE LISTING OF YOUR PROGRAM.
- * WHEN TO PARTITION THE MEMORY
 - * ONCE BEFORE READING IN CARDS.
 - * ONCE BEFORE MANUALLY ENTERING PROGRAM IN ORDER TO WRITE TO CARDS.

- * HOW TO START AND RUN YOUR PROGRAM
 - * TURN ON CALCULATOR
 - * PARTITION CALCULATOR
 - * LCALL ALL MODULE 1 CARDS
 - * OPTIONAL STEP: IF YOU SELECTED THE MANUAL DATA INPUT DENOTED IN YOUR BASIC PROGRAM BY USING THE "DATA" AND "READ" STATEMENTS THEN YOU MUST MANUALLY ENTER YOUR DATA INTO THE CALCULATOR MEMORY. THIS IS DONE BY REFERRING TO "INPUT DATA TO READ" TABLE PROVIDED AT THE END OF THIS LISTING. MANUALLY ENTER THE GIVEN DATA INTO THE REGISTERS USING THE FOLLOWING KEYSTROKES:

DATA
STO
XX

- * WHERE XX IS THE DESIRED REGISTER NUMBER.
- * INITIALIZE THE MANUAL SBR RETURN CONTROL STACK WITH THE FOLLOWING KEYSTROKES:

XX
STO
08

- * WHERE XX IS THE MANUAL RETURN REGISTER STACK TOP. (THIS IS GIVEN WITH THE PROGRAM LISTING NEAR THE PARTITION INFORMATION.)
- * PRESS "A" TO START.
- * CLICK DISPLAY PROMPTS.
 - * DEFINITIONS:
 - * RUN-TIME PROMPTS: ARE DEFINED TO BE CALCULATOR PROMPTS DISPLAYED IN THE CALCULATOR WINDOW IN THE FORM OF A 4 DIGIT DECIMAL, 2 DIGIT INTEGER OR A 1 DIGIT INTEGER. EACH PROMPT

```

IS OUTLINED BELOW:
* 4 DIGIT DECIMAL
  * FORMAT: X.YYY
    * X STANDS FOR MODULE NUMBER (1-9)
    * YYY STANDS FOR STARTING ADDRESS
  * ACTIONS:
    * LOAD ALL MODULE X CARDS.
    * PRESS FOLLOWING KEY SEQUENCE TO INITIALIZE:
      RCL
      00
      GTO
      Y
      Y
      Y
    * PRESS R/S TO CONTINUE IN NEW MOD.
* 2 DIGIT INTEGER
  * FORMAT: XX WHERE XX STANDS FOR A REGISTER NUMBER.
  * ACTIONS:
    * LOOK UP IN REGISTER MAP PROVIDED THE BASIC NAME THAT CORRESPONDS TO THE XX NUMBER.
    * ENTER THE BASIC VARIABLE VALUE.
    * PRESS R/S TO CONTINUE WITH THE ENTERED VALUE.
* 1 DIGIT INTEGER
  * FORMAT: X WHERE X IS A MODULE NUMBER.
  * ACTIONS:
    * LOAD ALL MODULE X CARDS.
    * PRESS FOLLOWING SEQUENCE TO INITIALIZE:
      RCL
      00
      INV
      SBR
    * PRESS R/S TO CONTINUE IN NEW MOD.
* PAUSE IN DISPLAY
  * AN UNFORMATTED DIGIT FLASHES IN THE DISPLAY BEFORE BEING DISPLAYED. THIS IS AN ANSWER THAT CORRESPONDS TO A REQUESTED ANSWER IN THE BASIC PROGRAM USING THE BASIC PRINT STATEMENT. THESE ANSWERS OCCUR IN THE SAME ORDER AS THEY WERE REQUESTED IN THE BASIC PROGRAM.
  * ACTIONS: NOTE ANSWER AND PRESS R/S.
* 888 IN DISPLAY
  * SPECIFIC PROMPT THAT INDICATES THAT THE PROGRAM HAS STOPPED EXECUTION.
  * ACTIONS: IF DESIRED FIND ANSWERS IN THE CALCULATOR MEMORY USING THE "TI-59 REGISTER TO NAME MAPPING" AT THE END OF THE INSTRUCTIONS.

```

* EXPECTED CONTROL FLOW PROMPTS BY MODULE FOLLOW:

\$82
\$83

```

=====
END BAX59 SEGMENTATION/INSTRUCTION: VERSION 1.0
=====

```

\$83
\$84

```

=====
BAX59 PROGRAM INSTRUCTIONS: VERSION 1.0
**** SEGMENTER FAILURE **** PROGRAM FAILURE****
BAX59 DIAGNOSTICS FOLLOW:
=====

```

* SEGMENTER FAILURES:

- * THE SEGMENTOR COULD NOT SEGMENT THE COMPILED PROGRAM IN A SATISFACTORY MANNER. POSSIBLE REASONS FOR THE FAILURE ARE GIVEN BELOW.
- * THERE ARE TWO TYPES OF SEGMENT BREAKS:
 - * A JUMP BREAK OCCURS THROUGH AN ABSOLUTE JUMP TO SOME PORTION OF CODE IN ANOTHER MODULE.
 - * A SUBROUTINE BREAK OCCURS THROUGH A SUBROUTINE INVOKE TO A SUBROUTINE WHOSE DEFINITION RESIDES IN ANOTHER MODULE.
- * SEGMENT FAILURE OCCURS WHEN ONE OF THE ABOVE BREAKS OCCURS INSIDE A BACKWARD JUMPING LOOP THAT COVERS MORE PROGRAM STEPS THAN IS AVAILABLE IN THE CALCULATOR MEMORY. SEGMENTATION IS NOT ALLOWED IN A LOOP AS IT IS IMPRACTICAL TO KEEP READING IN CARDS EVERY TIME THE PROGRAM LOOPS BACK OVER A BREAK (IMAGINE A 1 TO 1000 LOOP OVER SUCH A BREAK). TO AVOID SUCH A PROBLEM YOU MUST STRUCTURE YOUR BASIC PROGRAM TO AVOID LARGE BACKWARD-JUMPING LOOPS.
- * PROGRAM FAILURES: POSSIBLE PROGRAM FAILURE OCCURS WHEN SUBROUTINE CALLS ARE NESTED GREATER THAN SIX DEEP. THE CALCULATOR ONLY HAS SIX SUBROUTINE RETURN REGISTERS.
- * BELOW ARE DIAGNOSTICS INDICATING THE SIZES OF THE LOOPS IN TI-59 PROGRAM STEPS AND THE TYPES OF BREAKS OCCURRING WITHIN THESE LOOPS. DIAGNOSTICS ARE GIVEN IN ABSOLUTE CODE. SUBROUTINE NESTING LEVEL DIAGNOSTICS ARE GIVEN FOR INVOKED ROUTINE DEFINITION.

\$84
\$5

=====

EAX59 VERSION 1.0

UNSEGMENTED ABSOLUTE COMPILED TI59 CODE FOLLOWS

=====

\$5
\$6

* SEQUENTIAL CONTINUATION: 4 DIGIT REAL CODE.

\$6
\$7

* MANUAL RETURN FROM A SUBROUTINE: 1 DIGIT CODE.

\$7
\$8

* FORWARD JUMP CONTINUATION: 4 DIGIT REAL CODE.

\$8
\$9

* SUBROUTINE INVOKE: 4 DIGIT REAL CODE.

\$9
\$99

0		1		2		3
4		5		6		7
e		9		2ND E'		A
e		C		D		E
2ND A'		2ND B'		2ND C'		2ND D'
2ND CLR		INV		LN X		CE
CLR		2ND INV		2ND LOG		2ND CP
2ND TAN		X<=>T		Y**2		2ND SORT(X)
1/X		2ND PGM		2ND P=>R		2ND SIN
2ND CCS		2ND INC		STO		RCL
SUM		Y**X		2ND CMS		2ND EXC
2ND PRD		X		EE		(
2ND INT		/		2ND ENG		2ND FIX
2ND EXC	2ND IND	2ND DEG		GTO		2ND PGM 2ND IND
2ND X=T		2ND PRC 2ND IND		2ND OP		2ND PAUSE
SBR		2ND NCF		2ND RCL 2ND IND		2ND RAD
-		STO 2ND IND		SUM 2ND IND		2ND SUMMATION
X-BAR		2ND LEL		2ND X>=T		GTO 2ND IND
2ND OF 2ND IND		2ND GRAD		RST		2ND IFFLG
2ND D.MS		+		2ND STFLG		2ND R/S
INV SBR		2ND PI		2ND LIST		=
2ND WRITE		2ND DSZ		+/-		2ND PRT
				2ND ADV		

```

$99
$100      * EXPECTED PROMPTS FOR MODULE # $
$100
$101      * NONE
$101
$102      * $
$102
$103      * YES
$103
$104
$104      * MANUAL RETURN REGISTER TOP IS $
$104
$105      STORE IN REGISTER: $
$105
$106      * PROGRAM PARTITION IS $
$106
$107      * PARTITION NUMBER IS $
$107
$108      *MODULE # $
$108
$109      CARD #1
$109
$110      CARD #2
$110
$111      SIDE #1
$111
$112      SIDE #2
$112

```


APPENDIX K
ARTILLERY TEST PROGRAM SOURCE CODE

```

00005 CPTION 0 5
00010 REM *****
00011 REM
00012 REM
00014 REM
00015 REM
00016 REM
00017 REM
00018 REM
00019 REM
00020 REM
00022 REM
00023 REM
00024 REM
00025 REM
00026 REM
00027 REM
00050 REM *****
00055 REM ***** DATA SECTION M109 *****
00100 REM
00120 REM
00130 DATA
00140 DATA
00150 DATA
00160 REM
00170 DATA
00180 DATA
00190 DATA
00200 REM
00210 DATA
00220 DATA
00230 DATA
00240 REM
00250 DATA
00260 DATA
00270 DATA
00280 REM
00290 REM
00300 DATA
00310 REM
00320 REM
00330 DATA
00340 REM
00350 REM
00360 DATA
00370 REM
00380 REM
00390 DATA
00400 REM
00410 REM
00420 DATA
00430 REM
00440 REM
00450 DATA
00460 REM
00470 REM
00471 DATA
00472 REM
00473 REM

```

BAX59 TEST PROGRAM NUMBER1

*THIS TEST PROGRAM IS AN ADAPTATION OF THE PROGRAM USED BY THE FIELD ARTILLERY IN THE COMPUTATION OF FIRING DATA FOR THEIR GUNS. THE ORIGINAL PROGRAM WAS WRITTEN FOR THE TI-59 CALCULATOR. THIS TEST WAS CHOSEN NOT ONLY TO EVALUATE THE COMPILER AND SEGMENTOR BUT TO COMPARE THE RELATIVE EFFICIENCY OF THE TRANSLATED WEASIC PROGRAM WITH THAT OF A HUMAN CODED PROGRAM. BOTH PROGRAMS ACCOMPLISH THE SAME TASK.

***** CHARGE CONSTANTS M109A1 SELF PROPELLED *****

*CHARGE 4
-0.0133670, 21.2691, -105.7
-0.0001499, -0.06630, -0.41
.77, .01314, .00001720

*CHARGE 5
-0.0149331, 24.3439, 64.7
-0.0001420, -0.07069, .06
1.26, -0.1508, .00001678

*CHARGE 7
-0.0173835, 29.8741, 2255.2
-0.0001668, .08487, 3.29
1.3, .02713, .00001306

*CHARGE 8
-0.0182137, 32.3731, 4107.4
-0.0001668, .09272, 5.74
1.36, .02891, .00001410

*M109 MAX RANGE OF CURVE FIT BY CHARGE
5700, 7000, 10800, 17600

*M109 HIGH ANGLE CROSS OVER POINT MILS
715

*BATTERY DATA/ BTRYZ, BTRYN, BTRYA, BTRYL
0, 0, 0, 800

*REGISTRATION DATA/RNGK, DFCOR
1.0, 0

*TARGET DATA/OBSERVOR LOCATION (DUAL MEANING)
4000, 4000, 10

*OBSERVOR DATA
4000, -400, 10

*SPECIFIC CORRECTION FACTORS DATA
1018.5924, 1600, 3200

```

00480 REM ***** VARIABLE READ INITIALIZATION *****
00490 REM
00495 REM
00500 REM      *M109 BALLISTIC CONSTANTS BY CHARGE
00510 READ      *CHARGE 4
00520 READ      A24,A14,A04
00530 READ      C24,C14,C04
00540 READ      B04,B14,B24
00550 REM      *CHARGE 5
00560 READ      A25,A15,A05
00570 READ      C25,C15,C05
00580 READ      B05,B15,B25
00590 REM      *CHARGE 7
00600 READ      A27,A17,A07
00610 READ      C27,C17,C07
00620 READ      B07,B17,B27
00630 REM      *CHARGE 8
00640 READ      A28,A18,A08
00650 READ      C28,C18,C08
00660 READ      B08,B18,B28
00670 REM
00680 REM      *M109 MAX RANGE OF CURVE FIT VARIABLES
00690 READ      CHG4MAX,CHG5MAX,CHG7MAX,CHG8MAX
00700 REM
00710 REM      *M109 HIGH ANGLE CROSS OVER VARIABLE
00720 READ      HACRCSS
00730 REM
00740 REM      *BATTERY VARIABLES
00750 READ      ETRYE,ETRYN,BTRYA,BTRYL
00760 REM
00770 REM      *REGISTRATION VARIABLES
00780 READ      RGK,DFCOR
00790 REM
00800 REM      *TARGET VARIABLES OR OBSERVOR INIT LOCATION
00810 READ      GRIDE,GRIDN,GRIDA
00820 REM
00830 REM      *OBSERVOR VARIABLES
00840 READ      OT,LATDEV,RGDEV
00850 REM
00860 REM      *SPECIFIC CORRECTION FACTORS VARIABLES
00870 READ      MILRAD,ROTCOR,REFDEF
00880 REM
00885 REM      *****
00890 REM      ***** MAIN PROGRAM BEGINS *****
00900 REM      START
00910 REM      *COMPUTE TARGET GRID
00920 READ      GOSUE 1050
00940 REM      *COMPUTE GUN RANGE, AZIMUTH
00950 READ      GOSUE 1130
00970 REM      *COMPUTE FIRING DATA
00980 READ      GOSUE 1240
00990 REM
01000 REM      STOP
01010 REM      ***** MAIN STOP *****
01020 REM
01022 REM
01030 REM      ***** SUBROUTINES *****
01035 REM
01040 REM      *****
01050 REM      *** COMPUTE NEW TARGET GRID FROM SHIFTS ***
01051 REM      *****
01055 REM      START
01060 REM      GRIDN = GRIDN+(RGDEV*SIN((ROTCOR-OT)/MILRAD))-LATDEV*COS((ROTCOR-OT)/MILRAD) &
01070 &
01080 REM      GRIDE = GRIDE+(RGDEV*CCS((ROTCOR-OT)/MILRAD))+LATDEV*SIN((ROTCOR-OT)/MILRAD) &
01090 &
01100 REM      RETURN
01110 REM      *****
01115 REM
01120 REM      *****

```

```

01130 REM **** COMPUTE GUN RANGE, AZIMUTH *****
01131 REM **** *****
01135 REM START
01140 & TGTRG = SQR((GRIDE-BTRYE)**2+
01145 & (GRIDN-BTRYN)**2)
01150 & TGTAZ = ASIN((GRIDN-BTRYN)/TGTRG)*MILRAD
01160 IF GRIDE >= BTRYE
01170 TGTAZ = ROTCOR - TGTAZ
01180 ELSE
01190 TGTAZ = 3*ROTCOR + TGTAZ
01200 ENDIF
01210 RETURN
01220 REM *****
01223 REM *****
01230 REM *****
01240 REM **** FIRING DATA COMPUTATION ROUTINE *****
01245 REM **** *****
01250 REM START
01260 IF TGTRG <= CHG4MAX
01270 & INVCKE= FN_FD(A24,A14,A04,C24,C14,C04,
01275 & B24,E14,B04)
01280 ELSEIF TGTRG <= CHG5MAX
01290 & INVCKE= FN_FD(A25,A15,A05,C25,C15,C05,
01295 & B25,E15,B05)
01300 ELSEIF TGTRG <= CHG7MAX
01310 & INVCKE= FN_FD(A27,A17,A07,C27,C17,C07,
01315 & B27,E17,B07)
01320 ELSEIF TGTRG <= CHG8MAX
01330 & INVOKE= FN_FD(A28,A18,A08,C28,C18,C08,
01335 & B28,E18,B08)
01340 ELSE
01350 PRINT TGTRG
01360 ENDIF
01370 RETURN
01380 REM *****
01390 REM *****
01395 REM *****
01400 REM *** FIRING DATA COMPUTATION FUNCTION *****
01405 REM **** *****
01410 REM START
01420 DEP FN_FD(A2,A1,A0,C2,C1,C0,B2,B1,B0)
01430 & EL = (-A1+SQR(A1**2-(4*A2*(A0-TGTRG*RGK)
01435 & ))) / (2*A2)
01440 IF EL > HACROSS
01450 PRINT TGTAZ,TGTRG
01460 ELSE
01470 PRINT C0+C1*EL+C2*EL**2
01480 & PRINT REFDEF+DFCOR+(BTRYL-TGTAZ)+
01481 & (B0+B1*EL+B2*EL**2)
01490 PRINT EL+((GRIDA-BTRYA+20)/TGTRG*1000)
01500 ENDIF
01510 FNEND
01515 REM *****
01516 REM END FAX59 TEST PROGRAM NUMBER ONE
01518 REM **** *****

```

APPENDIX L
 TEST PROGRAM LISTING FILE (LISTF)

=====
 WBASIC PROGRAM LISTING
 =====

```

00005 OPTICN 0 5
00010 REM *****
00011 REM
00012 REM          BAX59 TEST PROGRAM NUMBER1
00014 REM
00015 REM *THIS TEST PROGRAM IS AN ADAPTATION OF THE PROGRAM
00016 REM USED BY THE FIELD ARTILLERY IN THE COMPUTATION
00017 REM OF FIRING DATA FOR THEIR GUNS. THE ORIGINAL
00018 REM PROGRAM WAS WRITTEN FOR THE TI-59 CALCULATOR.
00019 REM THIS TEST WAS CHOSEN TO NOT ONLY EVALUATE THE
00020 REM THE COMPILER AND SEGMENTOR BUT TO COMPARE THE
00022 REM THE RELATIVE EFFICIENCY OF THE TRANSLATED WEASIC
00023 REM PROGRAM WITH THAT OF A HUMAN CODED PROGRAM. BOTH
00024 REM PROGRAMS ACCOMPLISH THE SAME TASK.
00025 REM
00026 REM *****
00027 REM
00050 REM ***** DATA SECTION M109 *****
00055 REM
00100 REM *CHARGE CONSTANTS M109A1 SELF PROPELLED
00120 REM *CHARGE 4
00130 DATA -.0133670, 21.2691, -105.7
00140 DATA -.00001499, .06630, -.41
00150 DATA .77, .01314, .00001720
00160 REM *CHARGE 5
00170 DATA -.0149331, 24.3439, 64.7
00180 DATA -.00001420, .07069, .06
00190 DATA 1.26, .01508, .00001678
00200 REM *CHARGE 7
00210 DATA -.0173835, 29.8741, 2255.2
00220 DATA -.00001668, .08487, 3.29
00230 DATA 1.3, .02713, .00001306
00240 REM *CHARGE 8
00250 DATA -.0182137, 32.3731, 4107.4
00260 DATA -.00001668, .09272, 5.74
00270 DATA 1.36, .02891, .00001410
00280 REM
00290 REM *M109 MAX RANGE OF CURVE FIT BY CHARGE
00300 DATA 5700, 7000, 10800, 17600
00310 REM
00320 REM *M109 HIGH ANGLE CROSS OVER POINT MILS
00330 DATA 715
00340 REM
00350 REM *BATTERY DATA/ BTRYE, BTRYN, BTRYA, BTRYL
00360 DATA 0, 0, 0, 800
00370 REM
00380 REM *REGISTRATION DATA/RNGK, DFCOR
00390 DATA 1.0, 0
00400 REM
00410 REM *TARGET DATA/CBSERVOR LOCATION (EUAL MEANING)
00420 DATA 4000, 4000, 10
00430 REM
00440 REM *OESERVOR DATA
00450 DATA 4000, -400, 10
00460 REM
  
```

```

00470 REM      *SPECIFIC CORRECTION FACTORS DATA
00471 DATA    1018.5924,1600,3200
00472 REM
00473 REM
00480 REM      ***** VARIABLE READ INITIALIZATION *****
00490 REM
00495 REM      *M109 FALLISTIC CONSTANTS BY CHARGE
00500 REM      *CHARGE 4
00510 READ     A24,A14,A04
00520 READ     C24,C14,C04
00530 READ     B04,E14,B24
00540 REM      *CHARGE 5
00550 READ     A25,A15,A05
00560 READ     C25,C15,C05
00570 READ     B05,E15,B25
00580 REM      *CHARGE 7
00590 READ     A27,A17,A07
00600 READ     C27,C17,C07
00610 READ     B07,E17,B27
00620 REM      *CHARGE 8
00630 READ     A28,A18,A08
00640 READ     C28,C18,C08
00650 READ     B08,E18,B28
00660 REM
00670 REM      *M109 MAX RANGE OF CURVE FIT VARIABLES
00680 READ     CHG4MAX,CHG5MAX,CHG7MAX,CHG8MAX
00690 REM
00700 REM      *M109 HIGH ANGLE CROSS OVER VARIABLE
00710 READ     HACRCSS
00720 REM
00730 REM      *BATTERY VARIABLES
00740 READ     ETRYE,ETRYN,BTRYA,BTRYL
00750 REM
00760 REM      *REGISTRATION VARIABLES
00770 READ     RGK,DFCOR
00780 REM
00790 REM      *TARGET VARIABLES OR OBSERVOR INIT LOCATION
00800 READ     GRIDE,GRIDN,GRIDA
00810 REM
00820 REM      *OBSERVOR VARIABLES
00830 READ     OT,LATDEV,RGDEV
00840 REM
00845 REM      *SPECIFIC CORRECTION FACTORS VARIABLES
00860 READ     MILRAD,ROTCOR,REPDEF
00870 REM
00875 REM      *****
00880 REM      ***** MAIN PROGRAM BEGINS *****
00900 REM      START
00910 REM      *COMPUTE TARGET GRID
00920 REM      GOSUE 1050
00940 REM      *COMPUTE GUN RANGE,AZIMUTH
00950 REM      GOSUE 1130
00970 REM      *COMPUTE FIRING DATA
00980 REM      GOSUE 1240
00990 REM
01000 REM      STOP
01010 REM      ***** MAIN STOP *****
01020 REM
01022 REM
01030 REM      ***** SUBROUTINES *****
01035 REM
01040 REM      *****
01050 REM      *** COMPUTE NEW TARGET GRID FROM SHIFTS ***
01051 REM      *****
01055 REM      START
01060 REM      GRIDN = GRIDN+(RGDEV*SIN((ROTCOR-OT)/
&
01070 &          MILRAD)-LATDEV*COS((ROTCOR-OT)/MILRAD)
01080 REM      GRIDE = GRIDE+(RGDEV*COS((ROTCOR-OT)/
&
01090 &          MILRAD)+LATDEV*SIN((ROTCOR-OT)/MILRAD)

```

```

01100          RETURN
01110 REM *****
01115 REM *****
01120 REM *****
01130 REM ***** COMPUTE GUN RANGE, AZIMUTH *****
01131 REM *****
01135 REM          START
01140          TGTRG = SQR((GRIDE-BTRYE)**2+
01145 &          (GRIDN-BTRYN)**2)
01150          TGTAZ = ASIN((GRIDN-BTRYN)/TGTRG)*MILRAD
01160          IF GRIDE >= BTRYE
01170             TGTAZ = ROTCOR - TGTAZ
01180          ELSE
01190             TGTAZ = 3*ROTCOR + TGTAZ
01200          ENDIF
01210          RETURN
01220 REM *****
01223 REM *****
01230 REM *****
01240 REM ***** FIRING DATA COMPUTATION ROUTINE *****
01245 REM *****
01250 REM          START
01260          IF TGTRG <= CHG4MAX
01270             INVCKE= FN_FD(A24,A14,A04,C24,C14,C04,
01275 &             B24,E14,B04)
01290          ELSEIF TGTRG <= CHG5MAX
01290             INVCKE= FN_FD(A25,A15,A05,C25,C15,C05,
01295 &             B25,E15,B05)
01300          ELSEIF TGTRG <= CHG7MAX
01310             INVCKE= FN_FD(A27,A17,A07,C27,C17,C07,
01315 &             B27,E17,B07)
01320          ELSEIF TGTRG <= CHG8MAX
01330             INVCKE= FN_FD(A28,A18,A08,C28,C18,C08,
01335 &             B28,E18,B08)
01340          ELSE
01350             PRINT TGTRG
01360          ENDIF
01370          RETURN
01380 REM *****
01390 REM *****
01395 REM *****
01400 REM ***** FIRING DATA COMPUTATION FUNCTION *****
01405 REM *****
01410 REM          START
01420          DEF FN_FD(A2,A1,A0,C2,C1,C0,B2,B1,B0)
01430          EI = (-A1+SQR(A1**2-(4*A2*(A0-TGTRG*RGK)
01435 &          )))/(2*A2)
01440          IF EI > HACROSS
01450             PRINT TGTAZ,TGTRG
01460          ELSE
01470             PRINT C0+C1*EI+C2*EI**2
01480             PRINT REFDEF+DFCOR+(BTRYL-TGTAZ)+
01481 &             (B0+B1*EI+B2*EI**2)
01490             PRINT EI+((GRIDA-BTRYA+20)/TGTRG*1000)
01500          ENDFN
01510          F2END
01515 REM *****
01516 REM          END EAX59 TEST PROGRAM NUMBER ONE
01518 REM *****

```

```

=====
                        COMPILATION SUMMARY
=====

```

0 FATAL ERRORS.
0 WARNING MSGS.

81 IS NEXT AVAILABLE REGISTER
TOTAL REGISTERS RESERVED = 11

TOTAL REGISTERS USED = 70
 TOTAL LABELS USED = 5

----- COMPILATION TERMINATES -----

=====
 TI-59 CCDE TRANSLATED FROM WBASIC
 (UNSEGMENTED)
 =====

\$0	ADDR	CODE	
	000	76	2ND LEL
	001	11	A
	002	71	SBR
	003	12	E
	004	68	2ND NCP
	005	71	SBR
	006	13	C
	007	68	2ND NCP
	008	71	SBR
	009	14	D
	010	68	2ND NCP
	011	24	CE
	012	08	8
	013	08	8
	014	08	8
	015	91	R/S
	016	76	2ND LBL
	017	12	B
	018	53	{
	019	43	RCL
	020	59	59
	021	53	+
	022	53	{
	023	43	RCL
	024	63	63
	025	53	*
	026	53	{
	027	53	{
	028	43	RCL
	029	65	65
	030	75	-
	031	43	RCL
	032	61	61
	033	54)
	034	55	/
	035	43	RCL
	036	64	64
	037	54)
	038	58	2ND SIN
	039	75	-
	040	43	RCL
	041	62	62
	042	55	*
	043	53	{
	044	53	{
	045	43	RCL
	046	65	65
	047	75	-
	048	43	RCL
	049	61	61
	050	54)
	051	55	/
	052	43	RCL
	053	64	64
	054	54)
	055	39	2ND CCS

0056		}
0057		STO
0058		59
0059		(
0060		RCL
0061		58
0062		+
0063		(
0064		RCL
0065		63
0066		*
0067		(
0068		RCL
0069		65
0070		-
0071		RCL
0072		61
0073)
0074		/
0075		RCL
0076		64
0077)
0078		2ND
0079		CCS
0080		+
0081		RCL
0082		62
0083		*
0084		(
0085		RCL
0086		65
0087		-
0088		RCL
0089		61
0090)
0091		/
0092		RCL
0093		64
0094)
0095		2ND
0096		SIN
0097)
0098		STO
0099		58
100		INV
101		SER
102		2ND
103		IEL
104		C
105		(
106		RCL
107		58
108		-
109		RCL
110		52
111)
112		Y**X
113		2
114		+
115		(
116		RCL
117		59
118		-
119		RCL
120		53
121)
122		Y**X
123		
124		

125	02	2
126	54)
127	34	SQRT (X)
128	54)
129	42	STO
130	67	67
131	53	{
132	33	{
133	53	RCL
134	53	59
135	59	-
136	53	RCL
137	43	53
138	54)
139	55	/
140	55	RCL
141	43	67
142	67)
143	54	2ND INV
144	27	2ND SIN
145	38	*
146	55	RCL
147	43	64
148	64)
149	54	STO
150	42	68
151	68	RCL
152	43	58
153	58	X<=>T
154	22	RCL
155	43	52
156	52	X<=>T
157	32	INV
158	22	2ND X>=T
159	77	01
160	01	74
161	74	(
162	53	RCL
163	43	65
164	65	-
165	75	RCL
166	43	68
167	68)
168	54	STO
169	42	68
170	68	GT0
171	61	01
172	01	85
173	55	(
174	53	3
175	53	*
176	65	RCL
177	43	65
178	65	+
179	85	RCL
180	43	68
181	68)
182	54	STO
183	42	68
184	68	INV SER
185	92	2ND IRL
186	76	D
187	14	RCL
188	43	67
189	67	X<=>T
190	32	RCL
191	43	47
192	47	INV
193	22	

194	77	2ND	X>=T
195	02	02	
196	44	44	
197	53	(
198	43	RCL	
199	11	11	
200	42	STO	
201	71	71	
202	43	RCL	
203	12	12	
204	42	STO	
205	72	72	
206	43	RCL	
207	13	13	
208	42	STO	
209	73	73	
210	43	RCL	
211	14	14	
212	42	STO	
213	74	74	
214	43	RCL	
215	15	15	
216	42	STO	
217	75	75	
218	43	RCL	
219	16	16	
220	42	STO	
221	76	76	
222	43	RCL	
223	19	19	
224	42	STO	
225	77	77	
226	43	RCL	
227	18	18	
228	42	STO	
229	78	78	
230	43	RCL	
231	17	17	
232	42	STO	
233	79	79	
234	53	(
235	71	SBR	
236	15	E	
237	54)	
238	54)	
239	42	STO	
240	69	69	
241	61	GTO	
242	04	04	
243	16	16	
244	43	RCL	
245	67	67	
246	32	X<=>T	
247	43	RCL	
248	48	48	
249	22	INV	
250	77	2ND	X>=T
251	03	03	
252	00	00	
253	53	(
254	43	RCL	
255	20	20	
256	42	STO	
257	71	71	
258	43	RCL	
259	21	21	
260	42	STO	
261	72	72	
262	43	RCL	

263	22	STO
264	42	73
265	73	RCL
266	43	23
267	23	STO
268	42	74
269	74	RCL
270	43	24
271	24	STO
272	42	75
273	75	RCL
274	43	25
275	25	STO
276	42	76
277	76	RCL
278	43	28
279	28	STO
280	42	77
281	77	RCL
282	43	27
283	27	STO
284	42	78
285	78	RCL
286	43	26
287	26	STO
288	42	79
289	79	(
290	53	SBR
291	71	5
292	15)
293	54	STO
294	54	69
295	42	GTO
296	69	04
297	61	16
298	04	RCL
299	16	67
300	43	X<=>T
301	67	RCL
302	43	49
303	49	INV
304	22	2ND X>=T
305	77	03
306	03	56
307	56	(
308	53	RCL
309	43	29
310	29	STO
311	42	71
312	42	RCL
313	71	30
314	43	STO
315	30	72
316	42	RCL
317	72	31
318	43	STO
319	31	73
320	42	RCL
321	73	32
322	43	STO
323	32	74
324	32	RCL
325	42	33
326	74	STO
327	43	75
328	33	RCL
329	42	34
330	75	
331	43	
	34	

400	42	STO
999	76	RCL
999	43	37
999	37	STO
999	42	77
999	77	RCL
999	43	36
999	36	STO
999	42	78
999	78	RCL
999	43	35
999	35	STO
999	42	79
999	79	(
999	43	SBR
999	35	E
999	42	}
999	42	STO
999	69	GTO
999	61	04
999	64	16
999	16	RCL
999	43	67
999	67	X<=>T
999	43	RCL
999	50	50
999	22	INV
999	77	2ND X>=T
999	04	04
999	12	12
999	53	(
999	43	RCL
999	38	38
999	42	STO
999	71	71
999	43	RCL
999	39	39
999	42	STO
999	72	72
999	43	RCL
999	40	40
999	42	STO
999	73	73
999	43	RCL
999	41	41
999	42	STO
999	74	74
999	43	RCL
999	42	42
999	42	STO
999	75	75
999	43	RCL
999	43	43
999	42	STO
999	76	76
999	43	RCL
999	46	46
999	42	STO
999	77	77
999	43	RCL
999	45	45
999	42	STO
999	78	78
999	43	RCL
999	44	44
999	42	STO

401	79	
402	53	(
403	71	SBR
404	15	E
405	54)
406	54)
407	42	STO
408	69	69
409	61	GTO
410	04	04
411	16	16
412	43	RCL
413	67	67
414	99	2ND FRT
415	98	2ND ADV
416	92	INV SER
417	76	2ND LFL
418	15	E
419	00	0
420	42	STO
421	70	70
422	53	{
423	53	{
424	43	RCL
425	72	72
426	94	+/-
427	85	+
428	53	(
429	43	RCL
430	72	72
431	45	Y**X
432	02	2
433	75	-
434	53	(
435	04	4
436	05	*
437	43	RCL
438	71	71
439	65	*
440	53	(
441	43	RCL
442	73	73
443	75	-
444	43	RCL
445	67	67
446	65	*
447	43	RCL
448	56	56
449	54)
450	44)
451	44)
452	44	SQRT (X)
453	44)
454	55	/
455	53	{
456	02	2
457	65	*
458	43	RCL
459	71	71
460	54)
461	54)
462	42	STO
463	80	80
464	43	RCL
465	80	80
466	32	X<=>T
467	43	RCL
468	51	51
469	77	2ND X>=T

470	04	
471	82	
472	RCL	
473	68	
474	2ND	PRT
475	RCL	
476	67	
477	2ND	FRT
478	2ND	ADV
479	GTO	
480	05	
481	65	
482	(
483	RCL	
484	76	
485	+	
486	RCL	
487	75	
488	*	
489	RCL	
490	80	
491	+	
492	RCL	
493	74	
494	*	
495	RCL	
496	80	
497	Y**X	
498	2	
499)	
500	2ND	PRT
501	2ND	ADV
502	(
503	RCL	
504	66	
505	+	
506	RCL	
507	57	
508	+	
509	(
510	RCL	
511	55	
512	-	
513	RCL	
514	68	
515)	
516	+	
517	(
518	RCL	
519	79	
520	+	
521	RCL	
522	78	
523	*	
524	RCL	
525	80	
526	+	
527	RCL	
528	77	
529	*	
530	RCL	
531	80	
532	Y**X	
533	2	
534)	
535	2ND	PRT
536	2ND	ADV
537	(
538		

```

539 43 RCL
543 80 +
544 1 1 {
545 2 2 RCL
546 3 3 60
547 4 4 75
548 5 5 43
549 6 6 54
550 7 7 85
551 8 8 00
552 9 9 20
553 0 0 4
554 1 1 5
555 2 2 6
556 3 3 7
557 4 4 6
558 5 5 5
559 6 6 4
560 7 7 3
561 8 8 2
562 9 9 1
563 0 0 0
564 1 1 5
565 2 2 4
566 3 3 9
567 4 4 9
568 5 5 8
569 6 6 4
570 7 7 3
571 8 8 2
572 9 9 2

```

```

RCL
80
+
{
RCL
60
-
RCL
54
+
2
0
)
/
RCL
67
*
1
0
0
0
}
2ND FRT
2ND ADV
RCL
70
INV SER

```

```

=====
BAX59 SYMBOL TABLE DUMP
=====

```

BUCKET	CONTENTS	REG	TYP
03	REFDEF	66	GLOBAL VAR
05	LOG10	..	QUICK FN
06	GRIDA	60	GLOBAL VAR
08	EL	80	GLOBAL VAR
10	GRIDE	58	GLOBAL VAR
13	FP	..	QUICK FN
13	ABS	..	QUICK FN
14	FN_FD	70	PARAMETER FN
16	PI	..	CONSTANT
16	IP	..	QUICK FN
16	CSC	..	QUICK FN
18	SEC	..	QUICK FN
20	DPCOR	57	GLOBAL VAR
24	LOG	..	QUICK FN
26	GRIDN	59	GLOBAL VAR
26	RGK	56	GLOBAL VAR
26	RND	10	LONG FN
28	ACOS	..	QUICK FN
30	MILRAD	64	GLOBAL VAR

33	TAN	..	QUICK FN
	ATN	..	QUICK FN
	ASIN	..	QUICK FN
35	COS	..	QUICK FN
36	COT	..	QUICK FN
38	RGDEV	63	GLOBAL VAR
40	SIN	..	QUICK FN
41	OT	61	GLOBAL VAR
43	EXP	..	QUICK FN
53	LATDEV	62	GLOBAL VAR
59	SQR	..	QUICK FN
63	BTRYA	54	GLOBAL VAR
64	INVOKE	69	GLOBAL VAR
67	BTRYE	52	GLOBAL VAR
69	TGTRG	67	GLOBAL VAR
71	HACROSS	51	GLOBAL VAR
74	CHG4MAX	47	GLOBAL VAR
75	CHG5MAX	48	GLOBAL VAR
77	CHG7MAX	49	GLOBAL VAR
	A04	13	GLOBAL VAR
78	CHG8MAX	50	GLOBAL VAR
	A05	22	GLOBAL VAR
	E04	17	GLOBAL VAR
	A14	12	GLOBAL VAR
79	TGTAZ	68	GLOBAL VAR
	B05	26	GLOBAL VAR
	A15	21	GLOBAL VAR
	E14	18	GLOBAL VAR
	C04	16	GLOBAL VAR
	A24	11	GLOBAL VAR
80	A07	31	GLOBAL VAR
	B15	27	GLOBAL VAR
	C05	25	GLOBAL VAR
	A25	20	GLOBAL VAR
	B24	19	GLOBAL VAR
	C14	15	GLOBAL VAR
81	BTRYL	55	GLOBAL VAR
	A08	40	GLOBAL VAR
	E07	35	GLOBAL VAR
	A17	30	GLOBAL VAR
	B25	28	GLOBAL VAR
	C15	24	GLOBAL VAR
	C24	14	GLOBAL VAR
82	E08	44	GLOBAL VAR
	A18	39	GLOBAL VAR
	B17	36	GLOBAL VAR
	C07	34	GLOBAL VAR

	A27	29	GLOBAL VAR
	C25	23	GLOBAL VAR
83	-----	-----	-----
	BTRY N	53	GLOBAL VAR
	E 18	45	GLOBAL VAR
	C08	43	GLOBAL VAR
	A 28	38	GLOBAL VAR
	E27	37	GLOBAL VAR
	C 17	33	GLOBAL VAR
84	-----	-----	-----
	ROTCOR	65	GLOBAL VAR
	E28	46	GLOBAL VAR
	C 18	42	GLOBAL VAR
	C27	32	GLOBAL VAR
85	-----	-----	-----
	C28	41	GLOBAL VAR
	-----	-----	-----

APPENDIX M
 TEST PROGRAM NAME MAPPING FILE (NAMEP)

=====

 TI-59 REGISTER TO NAME MAPPING

 =====

REG#	BASIC NAME
11	A24
12	A14
13	A04
14	C24
15	C14
16	C04
17	B04
18	B14
19	B24
20	A25
21	A15
22	A05
23	C25
24	C15
25	C05
26	B05
27	B15
28	B25
29	A27
30	A17
31	A07
32	C27
33	C17
34	C07
35	B07
36	B17
37	B27
38	A28
39	A18
40	A08
41	C28
42	C18
43	C08
44	B08
45	B18
46	B28
47	CHG4MAX
48	CHG5MAX
49	CHG7MAX
50	CHG8MAX
51	HACROSS
52	BTRYE
53	BTRYN
54	BTRYA
55	BTRYL
56	RGK
57	DFCOR
58	GRIDE
59	GRIDN
60	GRIDA
61	OT
62	LATDEV
63	RGDEV

```

64 MILRAD
65 ROTCOR
66 REFDEF
67 TGTRG
68 TGTAZ
69 INVOKS
70 FN FD
71 (FN PARAMETER)
72 (FN PARAMETER)
73 (FN PARAMETER)
74 (FN PARAMETER)
75 (FN PARAMETER)
76 (FN PARAMETER)
77 (FN PARAMETER)
78 (FN PARAMETER)
79 (FN PARAMETER)
80 EL

```

APPENDIX N
 TEST PROGRAM DATA/READ MAPPING FILE (READF)

===== INPUT DATA TO READ MAPPING =====

DATA	REG	NAME
- .0133670	11	A24
21.2691	12	A14
-105.7	13	A04
- .00001499	14	C24
.06630	15	C14
-.41	16	C04
.77	17	B04
.01314	18	B14
.00001720	19	B24
- .0149331	20	A25
24.3439	21	A15
64.7	22	A05
- .00001420	23	C25
.07069	24	C15
.06	25	C05
1.26	26	B05
.01508	27	B15
.00001678	28	B25
- .0173835	29	A27
29.8741	30	A17
2255.2	31	A07
- .00001668	32	C27
.08487	33	C17
3.29	34	C07
1.3	35	B07
.02713	36	B17
.00001306	37	B27
- .0182137	38	A28
32.3731	39	A18
4107.4	40	A08
- .00001668	41	C28
.09272	42	C18
5.74	43	C08
1.36	44	B08
.02891	45	B18
.00001410	46	B28
5700	47	CHG4 MAX
7000	48	CHG5 MAX
10800	49	CHG7 MAX
17600	50	CHG8 MAX

715
0
0
0
800

1.0
0

4000
4000
10

4000
-400
10

1018.5924
1600
3200

51 HACROSS
52 BTRYE
53 BTRYN
54 BTRYA
55 BTRYL

56 RGK
57 DFCOR

58 GRIDE
59 GRIDN
60 GRIDA

61 OT
62 LATDEV
63 RGDEV

64 MILRAD
65 ROTCOR
66 REFDEF

APPENDIX O
 TEST PROGRAM LINK INTERFACE FILE (SCRATCH)

\$1
 81 IS NEXT AVAILABLE REG.
 \$1
 \$2

=====

 TI-59 CODE TRANSLATED FROM WBASIC

 (UNSEGMENTED)

 =====

\$0	AADR	CODE		
	C00	76	2ND	LEL
	001	11	A	
	002	71	SBR	
	003	12	B	
	004	68	2ND	NOP
	C05	71	SBR	
	006	13	C	
	C07	68	2ND	NOP
	008	71	SBR	
	C09	14	D	
	010	68	2ND	NOP
	011	24	CE	
	012	088	8	
	C13	088	8	
	014	088	8	
	015	91	R/S	
	016	76	2ND	LEL
	017	12	B	
	018	53	(
	C19	43	RCL	
	020	59	59	
	021	55	+	
	022	53	(
	023	43	RCL	
	024	63	63	
	025	55	*	
	026	53	{	
	027	43	RCL	
	028	65	65	
	029	75	-	
	C30	43	RCL	
	031	61	61	
	032	11)	
	033	55	/	
	034	43	RCL	
	035	64	64	
	036	44	}	
	037	58	2ND	SIN
	038	75	-	
	039	43	RCL	
	040	62	62	
	041	55	*	
	042	43	{	
	043	53	RCL	
	044	53	RCL	
	045	53	RCL	

046		65
047		-
048		RCL
049		61
050)
051		/
052		RCL
053		64
054)
055		2ND CCS
056)
057)
058		STO
059		59
060		(
061		RCL
062		58
063		+
064		(
065		RCL
066		63
067		*
068)
069		RCL
070		65
071		-
072		RCL
073		61
074)
075		/
076		RCL
077		64
078)
079		2ND CCS
080		+
081		RCL
082		62
083		*
084)
085		RCL
086		65
087		-
088		RCL
089		61
090)
091		/
092		RCL
093		64
094)
095		2ND SIN
096)
097)
098		STO
099		58
100		INV
101		2ND SER
102		LBL
103		C
104)
105)
106)
107		RCL
108		58
109		-
110		RCL
111		52
112)
113)
114		Y**X

115	0	2
116	0	+
117	0	(
118	0	RCL
119	0	59
120	0	-
121	0	RCL
122	0	53
123	0)
124	0	Y**X
125	0	2
126	0)
127	0	SQRT (X)
128	0)
129	0	STO
130	0	67
131	0	{
132	0	{
133	0	{
134	0	RCL
135	0	59
136	0	-
137	0	RCL
138	0	53
139	0)
140	0	/
141	0	RCL
142	0	67
143	0)
144	0	2ND INV
145	0	2ND SIN
146	0	*
147	0	RCL
148	0	64
149	0)
150	0	STO
151	0	68
152	0	RCL
153	0	58
154	0	X<=>T
155	0	RCL
156	0	52
157	0	X<=>T
158	0	INV
159	0	2ND X>=T
160	0	01
161	0	74
162	0	(
163	0	RCL
164	0	65
165	0	-
166	0	RCL
167	0	68
168	0)
169	0	STO
170	0	68
171	0	GTO
172	0	01
173	0	85
174	0	{
175	0	{
176	0	*
177	0	RCL
178	0	65
179	0	+
180	0	RCL
181	0	68
182	0)
183	0	STO

184	68	58	
185	92	INV	SBR
186	76	2ND	LEL
187	14	D	
188	43	RCL	
189	67	67	
190	32	X<=>T	
191	43	RCL	
192	47	47	
193	22	INV	
194	77	2ND	X>=T
195	02	02	
196	44	44	
197	53	(
198	43	RCL	
199	11	11	
200	42	STO	
201	71	71	
202	43	RCL	
203	12	12	
204	42	STO	
205	72	72	
206	43	RCL	
207	13	13	
208	42	STO	
209	73	73	
210	43	RCL	
211	14	14	
212	42	STO	
213	74	74	
214	43	RCL	
215	15	15	
216	42	STO	
217	75	75	
218	43	RCL	
219	16	16	
220	42	STO	
221	76	76	
222	43	RCL	
223	19	19	
224	42	STO	
225	77	77	
226	43	RCL	
227	18	18	
228	42	STO	
229	78	78	
230	43	RCL	
231	17	17	
232	42	STO	
233	79	79	
234	53	(
235	71	SBR	
236	15	=	
237	54)	
238	54	}	
239	42	STO	
240	69	69	
241	61	GTO	
242	04	04	
243	16	16	
244	43	RCL	
245	67	67	
246	32	X<=>T	
247	43	RCL	
248	48	48	
249	22	INV	
250	77	2ND	X>=T
251	03	03	
252	00	00	

253	43	(
254	20	RCL
255	42	STO
256	71	71
257	43	RCL
258	21	21
259	42	STO
260	72	72
261	43	RCL
262	22	22
263	42	STO
264	73	73
265	43	RCL
266	23	23
267	42	STO
268	74	74
269	43	RCL
270	24	24
271	42	STO
272	75	75
273	43	RCL
274	25	25
275	42	STO
276	76	76
277	43	RCL
278	28	28
279	42	STO
280	77	77
281	43	RCL
282	27	27
283	42	STO
284	78	78
285	43	RCL
286	26	26
287	42	STO
288	79	79
289	53	(
290	71	SB R
291	15)
292	54)
293	54)
294	42	STO
295	69	69
296	61	GTO
297	04	04
298	16	16
299	43	RCL
300	67	67
301	32	X<=>T
302	43	RCL
303	49	49
304	22	INV
305	77	2ND X>=T
306	03	03
307	56	56
308	53	(
309	43	RCL
310	29	29
311	42	STO
312	71	71
313	43	RCL
314	30	30
315	42	STO
316	72	72
317	43	RCL
318	31	31
319	42	STO
320	73	73
321	73	73

322	43	RCL
323	42	STO
324	74	74
325	43	RCL
326	33	33
327	42	STO
328	75	75
329	43	RCL
330	34	34
331	42	STO
332	76	76
333	43	RCL
334	37	37
335	42	STO
336	77	77
337	43	RCL
338	36	36
339	42	STO
340	78	78
341	43	RCL
342	35	35
343	42	STO
344	79	79
345	53	(
346	71	SBR
347	15	F
348	54)
349	42	STO
350	69	69
351	61	GTO
352	04	04
353	16	16
354	43	RCL
355	67	67
356	32	X<=>T
357	43	RCL
358	50	50
359	22	INV
360	77	2ND X>=T
361	04	04
362	12	12
363	53	(
364	43	RCL
365	38	38
366	42	STO
367	71	71
368	43	RCL
369	39	39
370	42	STO
371	72	72
372	43	RCL
373	40	40
374	42	STO
375	73	73
376	43	RCL
377	41	41
378	42	STO
379	74	74
380	43	RCL
381	42	42
382	42	STO
383	75	75
384	43	RCL
385	43	43
386	42	STO
387	76	76
388	43	RCL
389	43	43
390	43	RCL

46	46	STO	
42	77	RCL	
43	45	STO	
45	42	78	RCL
42	78	43	44
78	43	44	STO
43	44	44	79
44	42	79	(
40	01	53	SBR
40	02	71	E
40	03	15)
40	04	54	STO
40	06	54	69
40	07	42	GTO
40	08	69	04
40	09	61	16
41	10	04	RCL
41	11	16	67
41	12	43	2ND
41	13	67	ADV
41	14	99	SER
41	15	98	LBL
41	16	92	INV
41	17	76	2ND
41	18	15	E
41	19	00	0
42	20	42	STO
42	21	70	70
42	22	53)
42	23	53	RCL
42	24	43	72
42	25	72	+/-
42	26	94	+
42	27	85	(
42	28	53	RCL
42	29	43	72
43	00	72	Y**X
43	01	45	2
43	02	02	-
43	03	75	(
43	04	53	4
43	05	C4	*
43	06	65	RCL
43	07	43	71
43	08	71	*
43	09	65	(
44	00	53	RCL
44	01	43	73
44	02	73	-
44	03	75	RCL
44	04	43	67
44	05	67	*
44	06	65	RCL
44	07	43	56
44	08	56)
44	09	54	}
45	00	54	SQRT (X)
45	01	54)
45	02	34	/
45	03	44	(
45	04	55	2
45	05	55	*
45	06	02	RCL
45	07	65	71
45	08	43	
45	09	71	

460	54)
461	54	STO
462	42	80
463	80	RCL
464	43	80
465	80	X<=>T
466	32	RCL
467	43	51
468	51	2ND X>=T
469	77	04
470	04	82
471	82	RCL
472	43	68
473	68	2ND FRT
474	99	RCL
475	43	67
476	67	2ND FRT
477	99	2ND ADV
478	98	GTO
479	61	05
480	05	05
481	65	65
482	53	(
483	43	RCL
484	76	76
485	85	+
486	43	RCL
487	75	75
488	65	*
489	43	RCL
490	80	80
491	85	+
492	43	RCL
493	74	74
494	65	*
495	43	RCL
496	80	80
497	45	Y**X
498	02	2
499	54)
500	99	2ND PRT
501	98	2ND ADV
502	53	(
503	43	RCL
504	66	66
505	85	+
506	43	RCL
507	57	57
508	85	+
509	43	(
510	55	RCL
511	55	55
512	75	-
513	43	RCL
514	68	68
515	85)
516	85	+
517	53	(
518	43	RCL
519	79	79
520	85	+
521	43	RCL
522	78	78
523	65	*
524	43	RCL
525	80	80
526	85	+
527	43	RCL
528	77	77

5	6	*	
5	4	RCL	
5	3	80	
5	2	Y**X	
5	1	2	
5	0	}	
5	9	2ND	PRT
5	8	2ND	ADV
5	7	{	
5	6	RCL	
5	5	80	
5	4	+	
5	3	}	
5	2	RCL	
5	1	60	
5	0	-	
5	9	RCL	
5	8	54	
5	7	+	
5	6	2	
5	5	0	
5	4)	
5	3	/	
5	2	RCL	
5	1	67	
5	0	*	
5	9	1	
5	8	0	
5	7	0	
5	6	0	
5	5	0	
5	4	}	
5	3	2ND	PRT
5	2	2ND	ADV
5	1	RCL	
5	0	70	
5	9	INV	SER

-1 ----- END TI-59 CODE.

\$2

=====

TI-59 REGISTER TO NAME MAPPING

=====

REG#	BASIC NAME
11	A24
12	A14
13	A04
14	C24
15	C14
16	C04
17	B04
18	B14
19	B24
20	A25
21	A15
22	A05
23	C25
24	C15
25	C05
26	B05
27	B15
28	B25
29	A27
30	A17
31	A07
32	C27
33	C17
34	C07
35	B07
36	B17
37	B27
38	A28
39	A18
40	A08
41	C28
42	C18
43	C08
44	B08
45	B18
46	B28
47	CHG4MAX
48	CHG5MAX
49	CHG7MAX
50	CHG8MAX
51	HACROSS
52	BTRYE
53	BTRYN
54	BTRYA
55	BTRYL
56	RGK
57	DFCOR
58	GRIDE
59	GRIDN
60	GRI DA
61	OT
62	LATDEV
63	RGDEV
64	MILRAD
65	ROTCOR
66	REFDEP
67	TGTRG
68	TGTAZ
69	INVOKE
70	FN PD
71	{FN PARAMETER}
72	{FN PARAMETER}

73 (FN PARAMETER)
74 (FN PARAMETER)
75 (FN PARAMETER)
76 (FN PARAMETER)
77 (FN PARAMETER)
78 (FN PARAMETER)
79 (FN PARAMETER)
80

EL

\$3

\$4

=====
 INPUT DATA TO READ MAPPING
 =====

DATA	REG	NAME
-.0133670	11	A24
21.2691	12	A14
-105.7	13	A04
-.00001499	14	C24
.06630	15	C14
-.41	16	C04
.77	17	B04
.01314	18	B14
.00001720	19	B24
-.0149331	20	A25
24.3439	21	A15
64.7	22	A05
-.00001420	23	C25
.07069	24	C15
.06	25	C05
1.26	26	B05
.0150E	27	B15
.00001678	28	B25
-.0173835	29	A27
29.8741	30	A17
2255.2	31	A07
-.00001668	32	C27
.08487	33	C17
3.29	34	C07
1.3	35	B07
.02713	36	B17
.00001306	37	B27
-.0182137	38	A28
32.3731	39	A18
4107.4	40	A08
-.00001668	41	C28
.09272	42	C18
5.74	43	C08
1.36	44	B08
.02891	45	B18
.00001410	46	B28
5700	47	CHG4 MAX
7000	48	CHG5 MAX
10800	49	CHG7 MAX
17600	50	CHG8 MAX
715	51	HACROSS
0	52	BTRY E
0	53	BTRY N
0	54	BTRY A
800	55	BTRY L
1.0	56	RGK
0	57	DFCOR

4000
4000
10

400C
-400
10

1018.5924
1600
3200

\$4

58 GRIDE
59 GRIDN
60 GRIDA

61 OT
62 LATDEV
63 RGDEV

64 MILRAD
65 ROTCOR
66 REPDEF

APPENDIX P
TEST PROGRAM LINKER OUTPUT

=====

EAX59 PROGRAM INSTRUCTIONS: VERSION 1.0

=====

* CONGRATULATIONS, YOU HAVE JUST COMPILED A BASIC PROGRAM INTO A TI-59 PROGRAM. IN SO DOING IT IS VERY POSSIBLE THAT YOUR PROGRAM IS LARGER THAN THE MEMORY OF THE CALCULATOR. IF THIS IS THE CASE THEN THE PROGRAM HAS BEEN SEGMENTED AND PROMPTING CODE INSERTED TO GUIDE YOUR CALCULATOR PROGRAM DURING ITS EXECUTION. THE REMAINDER OF THIS OUTPUT CONSISTS OF TI-59 CODE LISTINGS AND OTHER INFORMATION TO AID YOU IN YOUR PROGRAM EXECUTION.

* THE FOLLOWING DEFINITIONS ARE PROVIDED AS AN AID TO READING THE PROGRAM LISTING FILE.

* DEFINITIONS:

- * MODULE: A MODULE IS DEFINED TO BE ALL THE MEMORY DEDICATED TO PROGRAM STEPS. THE SIZE IS VARIABLE AND IS DEPENDENT ON THE REGISTER REQUIREMENT. VALUES RANGE FROM 0 TO 239, 479 OR 719 DEPENDING ON THE AMOUNT OF REGISTERS USED BY THE PROGRAM.
- * CARD: A CARD IS DEFINED TO BE ONE MAGNETIC CARD. A CARD HOLDS 480 PROGRAM STEPS. THESE STEPS ARE NOT CONTIGUOUS BUT ARE ARRANGED ON THE TWO SIDES OF THE CARD.
- * SIDE: A SIDE IS ONE HALF OF A CARD. IT CONTAINS UP TO 240 STEPS. WHEN ONE SIDE OF A CARD IS READ BY THE CALCULATOR 240 PROGRAM STEPS ARE FILLED IN MEMORY. THESE BLOCKS OF 240 STEPS ARE REFERED TO AS "BANKS" IN THE MANUFACTURER LITERATURE. WHEN LOADING A CARD YOU WILL LOAD ONLY BANK NUMBER 1 AND/OR 2 FOR PROGRAM STEPS.
- * PARTITION: THIS IS DEFINED TO BE THE CURRENT SETTING OF CALCULATOR MEMORY AS APPLIED TO THE AMOUNT OF MEMORY DEDICATED TO STORAGE REGISTERS AND THE AMOUNT DEDICATED TO PROGRAM STEPS. WE WILL BE DEALING WITH 3 PARTITIONS. THESE ARE:

3	719.29
4	479.59
5	239.69

FORMAT: X YYY.ZZ

WHERE X STANDS FOR PARTITION NUMBER
 YYY STANDS FOR PROGRAM STEPS (0-YYY)
 ZZ STANDS FOR REGISTERS (0-ZZ).

=====

TI 59 PROGRAM LISTING BY MODULE/CARD/SIDE:

=====

- * THE FOLLOWING IS YOUR PROGRAM LISTING. THE PROGRAM IS LISTED ACCORDING TO MODULE NUMBER AND ITS ASSOCIATED CARDS AND CARD SIDES.
- * REFER TO THE TI-59 PROGRAMMER'S GUIDE ON HOW TO INPUT A PROGRAM AND WRITE IT TO MAGNETIC CARDS.
- * CAUTION: ENSURE THAT THE CORRECT CALCULATOR PARTITION IS SET BEFORE INPUTTING A PROGRAM AND WRITING TO MAGNETIC CARDS.
- * CAUTION: ENSURE THAT YOU DO NOT CONFUSE BANK NUMBERS WITH CARD/MODULE OR SIDE NUMBERS. THE NUMBERS WHICH REFER TO THE LISTING ARE AKIN TO A VIRTUAL ADDRESS AND DO NOT REPRESENT THE ACTUAL BANK NUMBER. IF IN DOUBT, REMEMBER TO USE THE TABLE BELOW TO TRANSLATE VIRTUAL TO ACTUAL BANK NUMBERS.

VIRTUAL BANK	ACTUAL BANK
MODULE #	
CARD1	
SIDE1 -----	BANK1
MODULE #	
CARD1	
SIDE2 -----	BANK2
MODULE #	
CARD2	
SIDE1 -----	BANK3

=====

TI-59 LISTING

=====

- * MANUAL RETURN REGISTER TOP IS 81
STORE IN REGISTER: 8
- * PROGRAM PARTITION IS 239.89
- * PARTITION NUMBER IS 9

*MODULE # 1
CARD #1
SIDE #1

000	76	2ND LBL
001	11	A
002	71	SBR
003	12	B
004	68	2ND NOP
005	71	SBR
006	13	C
007	68	2ND NOP
008	71	SBR
009	14	D
010	68	2ND NOP
011	24	CE
012	08	8
013	08	8

014	08
015	91
016	76
017	14
018	42
019	00
020	69
021	28
022	01
023	72
024	08
025	02
026	93
027	00
028	00
029	00
030	91
031	76
032	15
033	23
034	54
035	59
036	55
037	55
038	43
039	63
040	55
041	55
042	33
043	43
044	65
045	75
046	43
047	65
048	54
049	55
050	43
051	64
052	38
053	43
054	43
055	62
056	62
057	63
058	53
059	53
060	43
061	75
062	43
063	43
064	61
065	54
066	53
067	43
068	64
069	54
070	59
071	44
072	54
073	42
074	59
075	33
076	53
077	58
078	55
079	53
080	63
081	53
082	53

```

8
R/S
2ND LBL
D
STO
00
2ND OP
28
1
STO 2ND IND
08
2
.
0
0
R/S
2ND LBL
8
(
CL
59
+
RCL
63
X
)
RCL
65
-
RCL
61
)
/
RCL
64
)
2ND SIN
-
RCL
62
X
(
)
RCL
65
-
RCL
61
)
/
RCL
54
)
2ND COS
)
STO
59
(
CL
58
+
(
RCL
63
X

```

083	553	{
084	553	RCL
085	433	65
086	553	-
087	753	RCL
088	433	61
089	611)
090	554	/
091	553	RCL
092	433	64
093	644	}
094	554	2 ND COS
095	339	+
096	853	RCL
097	433	62
098	622	X
099	655	{
100	553	RCL
101	553	65
102	433	-
103	655	RCL
104	755	61
105	433)
106	611	/
107	554	RCL
108	555	64
109	433)
110	644	2 ND SIN
111	554	}
112	338	STO
113	554	58
114	554	INV
115	422	2 ND
116	558	LBL
117	922	C
118	766	{
119	133	RCL
120	553	58
121	553	-
122	553	RCL
123	433	53
124	558	6
125	753	-
126	433	RCL
127	552	52
128	554	}
129	455	**X
130	022	2
131	855	+
132	553	{
133	433	RCL
134	559	59
135	755	-
136	433	RCL
137	553	53
138	554	}
139	455	**X
140	502	2
141	554	}
142	334	SQRT (X)
143	554	}
144	422	STO
145	677	67
146	553	{
147	553	RCL
148	553	59
149	433	-
150	559	
151	755	

152
 153
 154
 155
 156
 157
 158
 159
 160
 161
 162
 163
 164
 165
 166
 167
 168
 169
 170
 171
 172
 173
 174
 175
 176
 177
 178
 179
 180
 181
 182
 183
 184
 185
 186
 187
 188
 189
 190
 191
 192
 193
 194
 195
 196
 197
 198
 199
 200

43
 53
 54
 55
 43
 67
 54
 27
 38
 65
 43
 64
 54
 42
 68
 43
 58
 32
 33
 52
 32
 22
 77
 01
 89
 53
 43
 65
 75
 43
 68
 54
 42
 68
 61
 02
 00
 53
 03
 65
 43
 65
 85
 43
 68
 54
 42
 68
 92

RCL
 53
)
 RCL
 67
)
 2ND INV
 2ND SIN
 X
 RCL
 64
)
 STO
 68
 RCL
 58
 X<=>T
 RCL
 52
 X<=>T
 INV
 2ND X>=I
 01
 89
 ()
 RCL
 65
 -
 RCL
 68
)
 STO
 68
 GTO
 02
 00
 ()
 X
 RCL
 65
 +
 RCL
 68
)
 STO
 68
 INV SBR

*MODULE # 2
 CARL #1
 SIDE #1

000	76	2ND LBL
001	14	D
002	43	RCL
003	67	57
004	32	X<=>T
005	43	RCL
006	47	47
007	22	INV
008	77	2ND X>=T
009	00	00
010	58	58
011	53	(
012	43	RCL
013	11	11
014	42	STO
015	71	71
016	43	RCL
017	12	12
018	42	STO
019	72	72
020	43	RCL
021	13	13
022	42	STO
023	73	73
024	43	RCL
025	14	14
026	42	STO
027	74	74
028	43	RCL
029	15	15
030	42	STO
031	75	75
032	43	RCL
033	16	16
034	42	STO
035	76	76
036	43	RCL
037	19	19
038	42	STO
039	77	77
040	43	RCL
041	18	18
042	42	STO
043	78	78
044	43	RCL
045	17	17
046	42	STO
047	79	79
048	53	(
049	71	S BR
050	15	E
051	54)
052	54)
053	42	STO
054	69	69
055	61	GTO
056	01	01
057	23	23
058	43	RCL
059	67	67
060	32	X<=>T
061	43	RCL
062	48	48
063	22	INV
064	77	2ND X>=T

065	01	01
066	32	32
067	53	(
068	43	RCL
069	20	20
070	42	STO
071	71	71
072	43	RCL
073	21	21
074	42	STO
075	72	72
076	43	RCL
077	22	22
078	42	STO
079	73	73
080	43	RCL
081	23	23
082	42	STO
083	74	74
084	43	RCL
085	24	24
086	42	STO
087	75	75
088	43	RCL
089	25	25
090	42	STO
091	76	76
092	43	RCL
093	28	28
094	42	STO
095	77	77
096	43	RCL
097	27	27
098	42	STO
099	78	78
100	43	RCL
101	26	26
102	42	STO
103	79	79
104	53	(
105	71	BR
106	15	E
107	54)
108	54)
109	42	STO
110	69	69
111	51	GTO
112	01	01
113	23	23
114	42	STO
115	00	00
116	24	CE
117	03	3
118	93	.
119	00	0
120	00	0
121	00	0
122	91	R/S
123	42	STO
124	00	00
125	24	CE
126	03	3
127	93	.
128	01	1
129	01	1
130	06	6
131	91	R/S
132	42	STO
133	00	00

0041	748
0042	433
0043	355
0044	422
0045	799
0046	533
0047	711
0048	155
0049	544
0050	544
0051	422
0052	699
0053	611
0054	016
0055	433
0056	433
0057	677
0058	322
0059	433
0060	500
0061	222
0062	777
0063	011
0064	122
0065	533
0066	433
0067	338
0068	422
0069	711
0070	433
0071	399
0072	422
0073	722
0074	433
0075	400
0076	422
0077	733
0078	433
0079	411
0080	422
0081	744
0082	433
0083	422
0084	422
0085	755
0086	433
0087	433
0088	422
0089	700
0090	433
0091	466
0092	422
0093	777
0094	433
0095	455
0096	422
0097	788
0098	433
0099	444
1000	422
1001	799
1002	533
1003	711
1004	155
1005	544
1006	544
1007	422
1008	666
1009	611

78	
RR	
CL	
55	
STO	
79	
BR	
BR	
STO	
69	
99	
STO	
00	
16	
CL	
67	
X<=>T	
CL	
50	
HIMV	
ND	X>=T
01	
12	
RR	
CL	
38	
STO	
71	
RR	
CL	
39	
STO	
72	
RR	
CL	
40	
STO	
73	
RR	
CL	
41	
STO	
74	
RR	
CL	
42	
STO	
75	
RR	
CL	
43	
STO	
76	
RR	
CL	
46	
STO	
77	
RR	
CL	
45	
STO	
78	
RR	
CL	
44	
STO	
79	
RR	
BR	
55	
STO	
69	
99	
STO	

110	01	01
111	16	16
112	43	RCL
113	67	57
114	99	2ND PRT
115	98	2ND ADV
116	42	STO
117	00	00
118	73	RCL 2ND IND
119	08	08
120	69	2ND OP
121	38	38
122	91	R/S
123	76	2ND LBL
124	15	E
125	42	STO
126	00	00
127	69	2ND OP
128	28	28
129	03	3
130	72	STO 2ND IND
131	08	08
132	04	4
133	93	.
134	00	0
135	00	0
136	00	0
137	91	R/S

*MODULE # 4
 CARD #1
 SIDE #1

000	76	2ND LBL
001	15	E
002	00	0
003	42	STO
004	70	70
005	53	{
006	53	{
007	43	RCL
008	72	72
009	94	+/-
010	85	+
011	53	{
012	43	RCL
013	72	72
014	45	Y**X
015	02	2
016	75	-
017	53	{
018	04	4
019	65	X
020	43	RCL
021	71	71
022	65	X
023	53	{
024	43	RCL
025	73	73
026	75	-
027	43	RCL
028	67	67
029	65	X
030	43	RCL
031	56	56
032	54	}
033	54	}
034	54	}

0035	34
0036	54
0037	55
0038	53
0039	02
0040	65
0041	43
0042	71
0043	54
0044	54
0045	42
0046	80
0047	43
0048	80
0049	32
0050	43
0051	51
0052	77
0053	00
0054	65
0055	43
0056	68
0057	99
0058	43
0059	67
0060	99
0061	98
0062	61
0063	01
0064	48
0065	53
0066	43
0067	76
0068	85
0069	43
0070	75
0071	65
0072	43
0073	80
0074	85
0075	43
0076	74
0077	65
0078	43
0079	80
0080	45
0081	02
0082	54
0083	99
0084	98
0085	53
0086	43
0087	66
0088	85
0089	43
0090	57
0091	85
0092	53
0093	43
0094	55
0095	55
0096	43
0097	68
0098	54
0099	85
1000	53
1001	43
1002	79
1003	85

```

S Q R T ( X )
)
/
{
X
RCL
71
}
S T O
80
RCL
80
X < = > T
RCL
51
2 N D X > = T
00
65
RCL
68
2 N D P R T
RCL
67
2 N D P R T
2 N D A D V
G T O
01
48
(
RCL
76
+
RCL
75
X
RCL
80
+
RCL
74
X
RCL
90
Y * * X
2
)
2 N D P R T
2 N D A D V
(
RCL
66
+
RCL
57
+
(
RCL
55
-
RCL
68
)
+
(
RCL
79
+

```

104	43	RCL
105	78	78
106	65	X
107	43	RCL
108	80	80
109	85	+
110	43	RCL
111	77	77
112	65	X
113	43	RCL
114	80	80
115	45	Y**X
116	02	2
117	54	}
118	54	}
119	99	2ND PRT
120	99	2ND ADV
121	53	{
122	43	RCL
123	80	80
124	85	+
125	53	{
126	53	{
127	43	RCL
128	60	60
129	75	-
130	43	RCL
131	54	54
132	85	+
133	02	2
134	00	0
135	54	}
136	55	/
137	43	RCL
138	67	67
139	65	X
140	01	1
141	00	0
142	00	0
143	00	0
144	54	}
145	54	}
146	99	2ND PRT
147	98	2ND ADV
148	43	RCL
149	70	70
150	42	STO
151	00	0
152	73	RCL 2ND IND
153	08	08
154	69	2ND OP
155	38	38
156	91	R/S

=====

TI-59 PROGRAM SPECIFIC INSTRUCTIONS:

=====

- * THE FOLLOWING INFORMATION WILL TELL YOU HOW TO RUN YOUR PROGRAM.
- * YOU MUST ENTER YOUR PROGRAM MANUALLY INTO THE CALCULATOR AND WRITE THE PROGRAM TO MAGNETIC CARDS. THIS STEP ONLY NEEDS TO BE ACCOMPLISHED ONCE. AFTER THAT, THE PROGRAM IS ENTERED USING THE MAGNETIC CARD FACILITY OF THE CALCULATOR. SEE THE MANUFACTURER'S LITERATURE ON ENTERING A PROGRAM AND WRITING IT TO MAGNETIC CARDS. YOU WILL NEED TO PARTITION MEMORY.

- * HOW TO PARTITION THE MEMORY
- * KEY SEQUENCE:

X
2ND
OP
17

* X IS THE PARTITION NUMBER GIVEN IN THE LISTING OF YOUR PROGRAM.

- * WHEN TO PARTITION THE MEMORY
- * ONCE BEFORE FEADING IN CARDS.
- * ONCE BEFORE MANUALLY ENTERING PROGRAM IN ORDER TO WRITE TO CARDS.

- * HOW TO START AND RUN YOUR PROGRAM

- * TURN ON CALCULATOR
- * PARTITION CALCULATOR
- * LOAD ALL MODULE 1 CARDS
- * OPTIONAL STEP: IF YOU SELECTED THE MANUAL DATA INPUT DENOTED IN YOUR BASIC PROGRAM BY USING THE "DATA" AND "READ" STATEMENTS THEN YOU MUST MANUALLY ENTER YOUR DATA INTO THE CALCULATOR MEMORY. THIS IS DONE BY REFERRING TO "INPUT DATA TO READ" TABLE PROVIDED AT THE END OF THIS LISTING. MANUALLY ENTER THE GIVEN DATA INTO THE REGISTERS USING THE FOLLOWING KEYSTROKES:

DATA
STO
XX

- * WHERE XX IS THE DESIRED REGISTER NUMBER.
- * INITIALIZE THE MANUAL SBR RETURN CONTROL STACK WITH THE FOLLOWING KEYSTROKES:

XX
STO
08

WHERE XX IS THE MANUAL RETURN REGISTER STACK TOP. (THIS IS GIVEN WITH THE PROGRAM LISTING NEAR THE PARTITION INFORMATION.)

```

* PRESS "A" TO START.
* FOLLOW DISPLAY PROMPTS.
* DEFINITIONS:
  * RUN-TIME PROMPTS: ARE DEFINED TO BE CALCULATOR
    PROMPTS DISPLAYED IN THE CALCULATOR WINDOW
    IN THE FORM OF A 4 DIGIT DECIMAL, 2 DIGIT
    INTEGER OR A 1 DIGIT INTEGER. EACH PROMPT
    IS OUTLINED BELOW:
    * 4 DIGIT DECIMAL
      * FORMAT: X.YYY
        * X STANDS FOR MODULE NUMBER (1-9)
        * YYY STANDS FOR STARTING ADDRESS
    * ACTIONS:
      * LOAD ALL MODULE X CARDS.
      * PRESS FOLLOWING KEY SEQUENCE TO
        INITIALIZE:
          RCL
          00
          GTO
          Y
          Y
          Y
        * PRESS R/S TO CONTINUE IN NEW MOD.
    * 2 DIGIT INTEGER
      * FORMAT: XX WHERE XX STANDS FOR A
        REGISTER NUMBER.
      * ACTIONS:
        * LOOK UP IN REGISTER MAP PROVIDED
          THE BASIC NAME THAT CORRESPONDS
          TO THE XX NUMBER.
        * ENTER THE BASIC VARIABLE VALUE.
        * PRESS R/S TO CONTINUE WITH THE
          ENTERED VALUE.
    * 1 DIGIT INTEGER
      * FORMAT: X WHERE IS A MODULE NUMBER.
      * ACTIONS:
        * LOAD ALL MODULE X CARDS.
        * PRESS FOLLOWING SEQUENCE TO
          INITIALIZE:
            RCL
            00
            INV
            SBR
          * PRESS R/S TO CONTINUE IN NEW MOD.
    * PAUSE IN DISPLAY
      * AN UNFORMATTED DIGIT FLASHES IN THE
        DISPLAY BEFORE BEING DISPLAYED.
        THIS IS AN ANSWER THAT CORRESPONDS
        TO A REQUESTED ANSWER IN THE BASIC
        PROGRAM USING THE BASIC PRINT
        STATEMENT. THESE ANSWERS OCCUR IN
        THE SAME ORDER AS THEY WERE
        REQUESTED IN THE BASIC PROGRAM.
      * ACTIONS: NOTE ANSWER AND PRESS R/S.
    * 888 IN DISPLAY
      * SPECIFIC PROMPT THAT INDICATES THAT
        THE PROGRAM HAS STOPPED EXECUTION.
      * ACTIONS: IF DESIRED FIND ANSWERS IN
        THE CALCULATOR MEMORY USING THE
        "TI-59 REGISTER TO NAME MAPPING"
        AT THE END OF THE INSTRUCTIONS.
* EXPECTED CONTROL FLOW PROMPTS BY MODULE FOLLOW:

```



```

* EXPECTED PROMPTS FOR MODULE # 1
* FORWARD JUMP CONTINUATION: 4 DIGIT REAL CODE.
* NONE
* SUBROUTINE INVOKE: 4 DIGIT REAL CODE.
* 2.000
* MANUAL RETURN FROM A SUBROUTINE: 1 DIGIT CODE.
* NONE
* SEQUENTIAL CONTINUATION: 4 DIGIT REAL CODE.
* NONE

* EXPECTED PROMPTS FOR MODULE # 2
* FORWARD JUMP CONTINUATION: 4 DIGIT REAL CODE.
* 3.116
* 3.000
* SUBROUTINE INVOKE: 4 DIGIT REAL CODE.
* 4.000
* MANUAL RETURN FROM A SUBROUTINE: 1 DIGIT CODE.
* NONE
* SEQUENTIAL CONTINUATION: 4 DIGIT REAL CODE.
* 3.000

* EXPECTED PROMPTS FOR MODULE # 3
* FORWARD JUMP CONTINUATION: 4 DIGIT REAL CODE.
* NONE
* SUBROUTINE INVOKE: 4 DIGIT REAL CODE.
* 4.000
* MANUAL RETURN FROM A SUBROUTINE: 1 DIGIT CODE.
* YES
* SEQUENTIAL CONTINUATION: 4 DIGIT REAL CODE.
* NONE

* EXPECTED PROMPTS FOR MODULE # 4
* FORWARD JUMP CONTINUATION: 4 DIGIT REAL CODE.
* NONE
* SUBROUTINE INVOKE: 4 DIGIT REAL CODE.
* NONE
* MANUAL RETURN FROM A SUBROUTINE: 1 DIGIT CODE.
* YES
* SEQUENTIAL CONTINUATION: 4 DIGIT REAL CODE.
* NONE

```

=====

INPUT DATA TO READ MAPPING

=====

DATA	REG	NAME
- .0133670	11	A24
21.2691	12	A14
-105.7	13	A04
-.00001499	14	C24
.06630	15	C14
-.41	16	C04
.77	17	B04
.01314	18	B14
.00001720	19	B24
-.0149331	20	A25
24.3439	21	A15
64.7	22	A05
-.00001420	23	C25
.07069	24	C15
.06	25	C05
1.26	26	B05
.01508	27	B15
.00001678	28	B25
-.0173835	29	A27
29.8741	30	A17
2255.2	31	A07
-.00001668	32	C27
.08487	33	C17
3.29	34	C07
1.3	35	B07
.02713	36	B17
.00001306	37	B27
-.0182137	38	A28
32.3731	39	A18
4107.4	40	A08
-.00001668	41	C28
.09272	42	C18
5.74	43	C08
1.36	44	B08
.02891	45	B18
.00001410	46	B28
5700	47	CHG4 MAX
7000	48	CHG5 MAX
10800	49	CHG7 MAX
17600	50	CHG8 MAX
715	51	HACROSS
0	52	BTRYE
0	53	BTRYN
0	54	BTRYA
800	55	BTRYL
1.0	56	RGK
0	57	DPCOR

4000
4300
10

4000
-400
10

1018.5924
1600
3200

58 GRIDE
59 GRIDN
60 GRIDA

61 OT
62 LATDEV
63 RGDEV

64 MILRAD
65 ROTCOR
66 REPDEF

=====

 TI-59 REGISTER TO NAME MAPPING

 =====

REG #	BASIC NAME
11	A24
12	A14
13	A04
14	C24
15	C14
16	C04
17	B04
18	B14
19	B24
20	A25
21	A15
22	A05
23	C25
24	C15
25	C05
26	B05
27	B15
28	B25
29	A27
30	A17
31	A07
32	C27
33	C17
34	C07
35	B07
36	B17
37	B27
38	A28
39	A18
40	A08
41	C28
42	C18
43	C08
44	B08
45	B18
46	B28
47	CHG4MAX
48	CHG5MAX
49	CHG7MAX
50	CHG8MAX
51	HACROSS
52	BTRYE
53	BTRYN
54	BTRYA
55	BTRYL
56	RGK
57	DFCOR
58	GRIDE
59	GRI DN
60	GRI DA
61	OT
62	LATDEV
63	RGDEV
64	MILRAD
65	ROTCOR
66	REFDEF
67	FGT RG
68	FGT AZ
69	INV OKE
70	FN PD
71	{ FN PARAMETER }
72	{ FN PARAMETER }
73	{ FN PARAMETER }

```
74      (FN PARAMETER F)
75      {FN PARAMETER R}
76      {FN PARAMETER E}
77      {FN PARAMETER R}
78      {FN PARAMETER R}
79      {FN PARAMETER R}
80      EL
```

```
=====
END BAX59 SEGMENTATION/INSTRUCTION: VERSION 1.0
=====
```

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