



\* TI PPC NOTES \*  
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V7n6, 1982

## NEWSLETTER OF

## THE TI PROGRAMMABLE CALCULATOR CLUB

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Lanham MD 20706, USA.

Many of you have written me lately, asking me to print a simple, yet efficient index for the two years of our existence. To tell you the truth, I have had such an index all along, stored on diskettes for my TI-99/4A computer. One list was sorted according to subject matter, which I called ARTICLE and a second one according to AUTHOR. Most of the time, that is all what is needed to find a specific article written by a specified author. The computer doesn't take more than three seconds to find what I want. So, with this issue you'll find the print-out of both indices, one for 1980 and one for 1981. This way, you can file each loose leaf at the head of its corresponding volume.

The actual printing was done on an Olympia electronic typewriter (and so is this editorial) pasted up and subsequently reduced to 66% of full size, so as to fit on an 8.5 by 11 inches sheet. (to my European friends that is a DIN A4 sheet. And you thought I didn't know about those things, did you?)

Instead of the usual 16 we have 18 pages this time. This to accomodate the extended treatment of the now famous Zimmerman puzzle. This challenge turned out to be a real success: so far (May 26th) I received 31 solutions, and all of them excellent! Starting on page 15 you will find the first 14 of them.

There is lots of expectation and many rumors about the National Computer Conference to be held in Houston, Texas at about mid-June. If TI shows something that we think we could make good use of, I'll publish an extra issue, entirely dedicated to "it".

Attention weekend sailors and navigators: Robert M. Elliott, 29 Ox Hill Road, Norwich CT, 06360 USA has developed a super program called CELESTIAL RUNNING FIX. This program provides the latitude and longitude of a running fix, using any two of up to four celestial observations. No plotting is required. The program is good for any celestial body: sun, moon, planet or star, using either natural or artificial horizon.

The program uses two mag cards and the printer, but can easily be modified for calculator-only use. Robert spent about 1500 hours on the development of it over the last three years. It is "purer than Ivory soap" and it has been tested with over 100 actual computations.

Robert will send you a copy of the documentation, the instructions how to use it, worked out examples and THE PROGRAM RECORDED ON TWO MAG CARDS, all for \$ 4.00. The thing is too large to fit in the TI PPC NOTES, so write Robert for a copy, please.

We had the Survival Guide for Master Library reprinted. From my experience with two earlier reprintings, they go fast. So, if you want to assure yourself a copy, send \$ 10.00 to the club as soon as possible. Members abroad please add \$ 2.00 for postage.

Peter Poloczec has accomplished the impossible: Combine Fast and Graphics modes in one and the same program. More on it in next issue.

Maurice E.T. Swinnen.

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Dear Mr. Swinnen,

Enclosed you will find what I hope to be the first if not the second program submission for the new AF-1. I hope this program will be of interest to your readers.

I will like to make some comments about the new AF-1 programmable calculator. First if you wish to get in contact with OK Instruments you use the following telephone numbers, which is not a toll free number but now it has the complete area code: (642) 774-3665 2-(886) 683-6763. Second, I did not like about the AF-1, among other things, is it's weight, too heavy, and it's size especially the size of the magnetic card, that with the increase of the postal service will make it hard on the users to submit a program.

I hope that OK Instruments will manufacture accessories such as capability to play pac-man, interface capabilities.

I wish that you could get in contact with me, my two telephone numbers are .....but remember, be careful what the number AF-1 stands for.

I think it should have been model AF-3 since it is the third time you did it to us.

Sincerely,

José M. G. Chula Vista CA.

Dear José,

I appreciate your comments on the new one. I think it is going to be a best seller. Also, many thanks for the nice 5 by 20 inch mag card you sent me. It has acquired a permanent place on the wall in my office at Walter Reed, right above my collection of (semi-dressed) Guardian Angels pictures.

Maurice.

Dear Maurice,

Please change my mailing address to ..... Also send me five (5) of the new OK# Instruments calculators and have a happy 13 33 35 24 27 21 32 32 27 36 16 13 45.

Very Truly Yours,

Tim N. San Diego, CA.

Dear Tim,

On request 1: will do.

On request 2: will do. Are you a sailor (San Diego?) and could you use an extra anchor maybe?

Su atentísimo y seguro servidor,

Maurice.

Dear Mr. Swinnen,

Your prompt response and excellent newsletter are well worthwhile the extra \$ 5.00 (for first class mail) which is enclosed.

I am still in the dark about interfacing the TI-59 with the outside world. Do you know of anyone who has done this via an RS-232 interface or IEEE control system? Surely someone must have done this in the past 5 years... Please let me know in or with your next newsletter.

By the way, what is the next programming puzzle?

Yours Sincerely,

Dave H. Canada.

Dear Dave,

I still have to receive a well-documented article on interfacing the TI-59 to an outside device. I have a few hardware articles available, such as a TV interface to the TI-59, such that listings will appear on the screen. But nobody has told me yet if they actually built the device and if it worked satisfactorily.

Maurice.

Dear Maurice,

I received v7n3 yesterday. I liked them a lot, but cannot find any clever reason why you are spending 2 (and more) pages we pay to make the 1 April joke. Two nice program (game programs, why not having some of them) can fit much better on those pages, don't you think so? By the way, I have a TRS80 program that combines a telephone numbers to spell words. I tried your number and get only one meaning APRIL JOKE.

I had an opportunity of playing with Sharp PC1211. I agree with your editorial (v7n3p1) completely. There is only one way you can use it better than TI-59: on the exams. You can put quite a nice amount of text and formulas on it. Some of my friends pointed it.

I hope I'll receive something from you soon.

Cordially,

D.R. Yugoslavia.

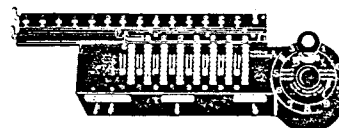
Dear Dejan,

Be a good sport now. Man doesn't live by bread alone. Once in a while he craves for something more "tasty."

You are still young. Don't take life too seriously or you'll soon turn into a 7687 7877.

Cordially,

Maurice.



BLACKJACK TUTOR,- by R<sup>2</sup> C<sup>2</sup> SNOW<sup>2</sup>. This program will TEACH you how to play BLACKJACK.  
----- It will deal cards for this well known "21" game and will allow the player to determine the hands. The player is never given a hand that totals 8 or more, as this would obviously mean a HIT. The Tutor will respond with a RIGHT for the correct strategy. It will also provide the correct strategy if the player chooses the incorrect one. If you press B, the Tutor will always provide you with the correct strategy.

The algorithm this program is based on is as follows:

1. When the dealer has a small card (2, 3, 4, 5, 6) STAND on hands of 13, 14, 15, 16.
2. When the dealer has a large card (7, 8, 9, 10, ace) DRAW until a count of 17 or greater has been reached.
3. DOUBLE DOWN when you have: a) hard 10, except when the dealer has 10 or ace, b) hard 9, except when the dealer has 7 through ace, c) hard 11, always, d) ace-2 through ace-5, when the dealer has 4, 5, 6, e) ace-6 when dealer has 2 through 6, f) ace-7 when dealer has 3 through 7.
4. SPLIT PAIRS: a) when you have 2s and the dealer has 3 through 7, b) when you have 3s and the dealer has 4 through 7, c) when you have 6s and the dealer has 2 through 6, d) when you have 7s and the dealer has 2 through 7, e) when you have 9s and the dealer has anything but ace-7-10, f) always when you have aces or 8s.
5. For ace-2 through ace-6, DRAW a card if not able to double down.
6. When holding ace-7: a) STAND if dealer has ace-2-7-8, b) DOUBLE DOWN if dealer has 3-4-5-6, c) DRAW if dealer has 9 or 10.
7. Always STAND on ace-8, ace-9.

INSTRUCTIONS: 1. Initialize, press E, enter seed, press E'.

2. To ask for help, press B. Correct strategy will be printed.

3. To play you own strategy: to STAND, press R/S, to HIT, press A, to SPLIT PAIRS, press C, to DOUBLE DOWN, press D.

RECORDING INSTRUCTIONS: Press 7 OP 17 and key in program and data (alpha) registers 09 through 69. Press 6 OP 17 and record four card sides.

Suggestion for R69: 10 STO 69 25131526 EE 11 +/- SUM 69 INV EE. Same procedure for R60.

<pre> 22222. 09 222222. 10 222222222. 11 33300. 12 33333. 13 33333. 14 33333. 15 33333. 16 333333333. 17 333333333. 18 333333333. 19 0. 20 0. 21 111110. 22 111110. 23 0. 24 22222222. 25 0. 26 11111. 27 1111111111. 28 3311311111. 29 3333333333. 30 1111111110. 31 0. 32 22200. 33 22200. 34 22200. 35 22200. 36 22200. 37 3000022223. 38 3333333333. 39 3333333333. 40 </pre>	<pre> 1617132717. DEALE 41 3536004133. RS UP 42 15133516. CARD 43 45324135. YOUR 44 1427131526. BLACK 45 2513152600. JACK 46 3637133116. STAND 47 2324370000. HIT 48 1417363700. BEST 49 14173700. BET 50 3633272437. SPLIT 51 33132435. PAIR 52 1632411427. DOUBL 53 16324331. DOWN 54 1731371735. ENTER 55 36171716. SEED 56 3524222337. RIGHT 57 3741373235. TUTOR 58 7.00000008 59 11.00000131517 60 2.00000003 61 3.00000004 62 4.00000005 63 5.00000006 64 6.00000007 65 7.00000008 66 8.00000009 67 9.00000012 68 10.00025131526 69 </pre>	<pre> BLACKJACK TUTOR  R/S    STAND A      HIT B      BEST BET C      SPLIT PAIR D      DOUBL DOWN  ENTER SEED  456123.  DEALERS UP CARD 7 YOUR CARD 8 17. 9 BEST STAND DEALERS UP CARD JACK YOUR CARD 7 16. 9 BEST HIT DEALERS UP CARD 4 YOUR CARD 5 9. 4 RIGHT DEALERS UP CARD 9 YOUR CARD 8 15. 7 BEST HIT </pre>
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000	83	GD*	063	87	IFF	126	32	X:T	188	43	RCL	250	00	0	312	03	3
001	05	05	064	06	06	127	02	2	189	01	01	251	48	EXC	313	69	DP
002	76	LBL	065	00	00	128	01	1	190	54	>	252	02	02	314	01	01
003	18	C'	066	68	68	129	67	EQ	191	22	INV	253	44	SUM	315	43	RCL
004	00	0	067	18	C'	130	02	02	192	59	INT	254	02	02	316	48	48
005	48	EXC	068	43	RCL	131	66	66	193	65	*	255	95	=	317	17	B'
006	04	04	069	02	02	132	22	INV	194	01	1	256	22	INV	318	01	1
007	76	LBL	070	91	R/S	133	87	IFF	195	00	0	257	87	IFF	319	04	4
008	17	B'	071	61	GTD	134	04	04	196	95	=	258	01	01	320	69	DP
009	69	DP	072	00	00	135	01	01	197	59	INT	259	01	01	321	01	01
010	03	03	073	36	36	136	45	45	198	67	EQ	260	56	56	322	43	RCL
011	69	DP	074	86	STF	137	01	1	199	02	02	261	01	1	323	50	50
012	05	05	075	05	05	138	01	1	200	18	18	262	02	2	324	69	DP
013	69	DP	076	43	RCL	139	77	GE	201	32	X:T	263	61	GTD	325	04	04
014	00	00	077	41	41	140	01	01	202	00	0	264	01	01	326	43	RCL
015	00	0	078	69	DP	141	45	45	203	67	EQ	265	56	56	327	49	49
016	92	RTN	079	01	01	142	22	INV	204	02	02	266	43	RCL	328	17	B'
017	76	LBL	080	43	RCL	143	86	STF	205	25	25	267	45	45	329	01	1
018	19	D'	081	42	42	144	04	04	206	01	1	268	69	DP	330	05	5
019	43	RCL	082	69	DP	145	32	X:T	207	67	EQ	269	02	02	331	69	DP
020	43	43	083	02	02	146	69	DP	208	02	02	270	43	RCL	332	01	01
021	42	STD	084	19	D'	147	06	06	209	29	29	271	46	46	333	43	RCL
022	04	04	085	75	-	148	69	DP	210	02	2	272	17	B'	334	52	52
023	87	IFF	086	01	1	149	00	00	211	67	EQ	273	81	RST	335	69	DP
024	05	05	087	95	=	150	87	IFF	212	02	02	274	76	LBL	336	04	04
025	00	00	088	22	INV	151	01	01	213	37	37	275	15	E	337	43	RCL
026	63	63	089	28	LOG	152	02	02	214	43	RCL	276	02	2	338	51	51
027	01	1	090	42	STD	153	49	49	215	47	47	277	85	+	339	17	B'
028	01	1	091	01	01	154	43	RCL	216	17	B'	278	76	LBL	340	01	1
029	82	HIR	092	18	C'	155	02	02	217	81	RST	279	10	E'	341	06	6
030	44	44	093	43	RCL	156	91	R/S	218	43	RCL	280	07	7</			

000	76	LBL	062	09	9	123	67	67	184	01	01	245	90	90	306	39	39
001	19	D'	063	95	=	124	50	I×I	185	88	88	246	22	INV	307	29	CP
002	53	(	064	61	GTD	125	32	X:T	186	01	1	247	87	IFF	308	22	INV
003	50	I×I	065	00	00	126	00	0	187	94	+/-	248	05	05	309	59	INT
004	85	+	066	88	88	127	67	EQ	188	82	HIR	249	01	01	310	22	INV
005	01	1	067	76	LBL	128	03	03	189	03	03	250	90	90	311	67	EQ
006	85	+	068	11	A	129	41	41	190	87	IFF	251	05	5	312	03	03
007	34	FX	069	32	X:T	130	01	1	191	05	05	252	32	X:T	313	39	39
008	23	LN×	070	07	7	131	52	EE	192	02	02	253	43	RCL	314	73	RC*
009	59	INT	071	69	DP	132	01	1	193	08	08	254	64	64	315	68	68
010	65	×	072	17	17	133	03	3	194	29	CP	255	67	EQ	316	50	I×I
011	02	2	073	32	X:T	134	22	INV	195	82	HIR	256	02	02	317	55	+
012	54	)	074	82	HIR	135	77	GE	196	13	13	257	87	87	318	01	1
013	92	RTN	075	04	04	136	02	02	197	77	GE	258	04	4	319	52	EE
014	76	LBL	076	42	STD	137	70	70	198	02	02	259	32	X:T	320	01	1
015	18	C'	077	68	68	138	01	1	199	08	08	260	43	RCL	321	06	6
016	50	I×I	078	42	STD	139	52	EE	200	86	STF	261	65	65	322	65	×
017	32	X:T	079	69	69	140	03	3	201	05	05	262	67	EQ	323	82	HIR
018	01	1	080	82	HIR	141	94	+/-	202	04	4	263	02	02	324	37	37
019	00	0	081	14	14	142	77	GE	203	00	0	264	87	87	325	01	1
020	00	0	082	91	R/S	143	02	02	204	18	C'	265	00	0	326	52	EE
021	49	PRD	083	76	LBL	144	70	70	205	61	GTD	266	18	C'	327	01	1
022	66	66	084	12	B	145	05	5	206	02	02	267	61	GTD	328	02	2
023	32	X:T	085	85	+	146	42	STD	207	39	39	268	02	02	329	95	=
024	44	SUM	086	01	1	147	64	64	208	43	RCL	269	51	51	330	22	INV
025	66	66	087	95	=	148	01	1	209	67	67	270	32	X:T	331	59	INT
026	97	DS2	088	22	INV	149	42	STD	210	55	+	271	75	-	332	55	+
027	64	64	089	44	SUM	150	65	65	211	82	HIR	272	52	EE	333	01	1
028	00	00	090	69	69	151	25	CLR	212	13	13	273	95	=	334	00	0
029	41	41	091	82	HIR	152	42	STD	213	22	INV	274	29	CP	335	00	0
030	05	5	092	14	14	153	66	66	214	28	LOG	275	67	EQ	336		

listed into banks 3 and 4 and let the listing program add 60 to it!

So, his user instructions say: Enter the low register number and press A. Enter the high register number and press B. To list press C.

The foregoing only when you want registers 00 through 59.

Now when you want to list registers 60 through 99 you first of all force those banks into banks 1 and 2 by means of -1 and -2. You then enter the correct low register and press D, enter the correct high register and press E, after which C will start your list.

In the example on the right you see in the left column data listed from registers 05 through 20, on the top listed by means of INV LIST and on the bottom by means of this program.

Now in the right hand column you see data which were originally in registers 60 through 73, forced into bank 1 and listed by means of INV LIST. On the bottom the same data listed by means of this program. Note the correct register numbers, achieved by entering 60 into D and 73 into E, after which C was pressed to start the listing.

5151513524.	05	2332430015.	00
4217350015.	06	3241271600.	01
3532363624.	07	4532410016.	02
3122515151.	08	3200000000.	03
1331004700.	09	3700230013.	04
4323320071.	10	37007100.	05
30133100.	11	1532312235.	06
21324400.	12	1337365700.	07
23173100.	13	4532410016.	08
1532353100.	14	2416002437.	09
6255555555.	15	4335323122.	10
5555555562.	16	17313735.	11
370100.	17	4500737373.	12
3001421700.	18	5555555555.	13
3335173636.	19		
6200000000.	20		
5151513524. ***RI	05	2332430015. HOW C	60
4217350015. VER C	06	3241271600. DULD	61
3532363624. ROSSI	07	4532410016. YOU D	62
3122515151. NG***	08	3200000000. D	63
1331004700. AN +	09	3700230013. T H A	64
4323320071. WHO ?	10	37007100. T ?	65
30133100. MAN	11	1532312235. CONGR	66
21324400. FOX	12	1337365700. ATS.	67
23173100. HEN	13	4532410016. YOU D	68
1532353100. CORN	14	2416002437. ID IT	69
6255555555. !(((	15	4335323122. WRONG	70
5555555562. (((	16	17313735. ENTR	71
370100. TO	17	4500737373. Y ???	72
3001421700. MOVE	18	5555555555. (((	73
3335173636. PRESS	19		
6200000000. ?	20		

FAST MODE and GRAPHICS MODE IN THE SAME PROGRAM? Peter Poloczec of Frankfurt, West Germany accomplished this feat recently and showed a method how it could be done. It requires rather a lot of button pushing, but, with time, that could be reduced to an absolute minimum and, who knows, maybe even made transparent to the user. Dejan Ristanović, who figures rather prominently in this issue, has used this technique to write a fascinating program called FAST SCRIBBLER. It permits you to draw pictures such as Frédéric De Mees' Mickey Mouse picture in Graphics Mode at the increased speed of Fast Mode. And, using Richard Snow's technique in the Stars and Stripes program, it prints out the required buttons to be pushed by means of a programmed LIST.

If I still have enough space left over in this issue, I will publish Dejan's FAST SCRIBBLER. Otherwise, it will appear in next issue.

SUPERCKECKSUM. It is my intention to use one of the many versions in existence to enable verification of correct keying-in of programs, by publishing the checksum with each published program, just as PPX is doing. Currently there is an evaluation of the many versions of checksum programs going on among reviewers in the US, in Denmark and in Sweden. Elmer Clausen enhanced Jules Bell's program quite a lot, while Lars Hedlund rewrote it in Fast Mode. That last one is a joy to use and seems to be the winner up to now. But then, Hans Peter Nielsen, the editor of PGM in Denmark says that it may be possible to even improve on the algorithm Jules used. Robert Prins in Holland also casts doubt on Jules' statement that the algorithm is absolutely fail-safe. I will keep you informed.

RANDOM NUMBERS,- Bertrand N. Bauer of Parameter Investigation, 2423 Ridgeway Ave., Evanston, IL, USA, writes in Keyboard, Sept-Oct., 1981(*Keyboard is an HP Journal for users of the 9825 desktop calculator, one of the few at HP that do not use RPN but rather the AOS system*) about a method of producing random numbers good enough for professional use. He bases his findings on an algorithm developed by Knuth in Donald E. Knuth, Seminumerical Algorithms, Vol.2, The Art of Computer Programming, Reading MA, Addison-Wesley, 2nd printing with revisions, 1971, pp. 82-97.

The algorithm passes from one 11-digit random fraction to the next by multiplying the current random fraction by 31 622 777 021, then adding  $10^{-11}$  and finally retaining only the fractional portion of the result. Or, as the mathematicians would write:

$$R' = \text{fra}((31\ 622\ 777\ 021)(R) + 10^{-11})$$

This algorithm has a lot of desirable characteristics: It generates every possible 11-digit fraction before it repeats itself. Thus, the "period" of this generator is  $10^{11}$  or 100 billion, that is the number of values of  $R'$  before repetition occurs. And this period is always 100 billion, no matter what the starting seed is.

Knuth also talks about a randomness test in the reference above. He calls it a spectral test and the result of it are three values:  $C_2$ ,  $C_3$  and  $C_4$ , which will determine if a certain algorithm to produce random numbers is good enough for specific categories of use. It so happens that the above algorithm has a  $C_2=1.22$ , a  $C_3=1.10$  and a  $C_4=3.84$ , which Knuth categorizes as good enough for professional use.

The limitations are that, if the seed is the same for each different run, the same series of random numbers will be produced. This might be desirable if you want to test a specific program, to compare results over several runs. If a randomized series of random numbers has to be produced, a "random" seed has to be entered each time to start a new run. On the TI58/59 this can be accomplished rather easily by writing at 000:

LBL E OP 29 RST ; press E and go drink a cup of coffee. When you come back, press R/S and RCL 09. You will find a suitably random number in R09. To convert it to a fraction, divide it by 1 EE 11.

From the above you might have gleaned a serious shortcoming of the TI-59: it should be possible to multiply two 11-digit numbers in order to execute the algorithm. Direct multiplication is, of course, not possible, but we can use "double precision arithmetic" to get to the required accuracy. The idea behind this is to write the 11-digit number as the sum of a 5-significant-digit number and a 6-significant-digit number, thus  $31\ 622\ (10^6) + 777\ 021$ . Then do the same to the random fraction. Next we apply the algebraic theorem  $(a+b)(c+d) = ac + ad + bc + bd$ . Individually, each of these four products has 12 significant digits maximum. We need only retain the digits to the right of the decimal point to obtain the desired accuracy. It turns out that the product  $(ac)$  never has any such digits, so we may discard this term.

So, now we may write a program according to the following procedure:

Step 1:  $Y = (\text{int}((R)(10^5)))(10^{-5})$

Step 2:  $Z = R - Y$

Step 3:  $U = \text{fra}(31\ 622)(10^6)(Z)$

Step 4:  $V = \text{fra}((777\ 021)(Y))$

Step 5:  $W = \text{fra}((777\ 021)(Z))$

Step 6:  $R' = \text{fra}(U + V + W + 10^{-11})$

Test data from a seed of .500 000 000 01 yield  $Y=.5$ ;  $Z=.000\ 000\ 000\ 01$ ;  
 $U=.316\ 22$ ;  $V=.5$ ;  $W=.000\ 007\ 770\ 21$  and  $R'=.816\ 227\ 770\ 22$ ;

With the program I have included a few more test data, which were verified on the 9825. As a 13-digit output I used Charlie Williamson's 13-digit printer from v7n4/5p11. I accessed it by leaving off the last R/S of my program and writing GTO NNN instead. The step NNN is the first step of his LBL A. Admittedly, you will have to wait 33 sec to print the final 13-digit output, but it was only done to have a nice print out. If you adopt this random generator program for one of your programs, it will, of course, not be necessary to also include a 13-digit printer.

Also, the PRT statements to print out  $Y$ ,  $Z$ ,  $U$ ,  $V$ , and  $W$  can be omitted and are used here only as debugging aids.

Maurice E.T. Swinnen

Random Numbers, M.Swinnen, listing.

1. -11	000 76 LBL	024 00 00	048 43 RCL	072 59 INT
0.	001 11 A	025 95 =	049 01 01	073 99 PRT
1. -11	002 99 PRT	026 94 +/-	050 65 X	074 85 +
0.31622	003 42 STD	027 42 STD	051 07 7	075 43 RCL
0.	004 00 00	028 02 02	052 07 7	076 05 05
.0000077702	005 65 X	029 99 PRT	053 07 7	077 85 +
3.16227702200 E-01	006 01 1	030 65 X	054 00 0	078 43 RCL
	007 52 EE	031 01 1	055 02 2	079 03 03
.3162277702	008 05 5	032 52 EE	056 01 1	080 85 +
0.31622	009 95 =	033 06 6	057 42 STD	081 01 1
.0000077702	010 59 INT	034 65 X	058 04 04	082 52 EE
0.89684	011 65 X	035 03 3	059 95 =	083 01 1
0.58062	012 01 1	036 01 1	060 22 INV	084 01 1
.0376241146	013 52 EE	037 06 6	061 59 INT	085 94 +/-
5.150841146300 E-01	014 05 5	038 02 2	062 99 PRT	086 95 =
	015 94 +/-	039 02 2	063 42 STD	087 22 INV
.5150841146	016 95 =	040 95 =	064 05 05	088 52 EE
0.51508	017 22 INV	041 22 INV	065 43 RCL	089 22 INV
.0000041146	018 52 EE	042 59 INT	066 02 02	090 59 INT
0.82986	019 42 STD	043 22 INV	067 65 X	091 42 STD
0.97668	020 01 01	044 52 EE	068 43 RCL	092 00 00
.1971539172	021 99 PRT	045 99 PRT	069 04 04	093 99 PRT
3.693917240000 E-03	022 75 -	046 42 STD	070 95 =	094 91 R/S
	023 43 RCL	047 03 03	071 22 INV	

HEX CODES,-Dejan Ristanović in Belgrade, Yugoslavia has been experimenting with different hex codes on step 16, this after the usual initialization PGM12 SBR 444 R/S... The experiments were done with the printer attached. Without the printer they may behave differently. Also, Dejan complains that most of the hex codes work differently in Trace mode versus Normal mode.

In the following list the character or digits in parenthesis following the hex codes is the code of the instruction to be put at step 16.

- " h62 (A) in Trace list as OSS\* and does some strange things. The next two digits make up the address of a register.
- h82 (31) lists as \*GT
- h73 (22) lists as \*ST. The next two digits are the address of a register. 12 Digits of that register, the integer part, are printed. This might be useful sometimes with the calculator used manually. Unfortunately, it works only in Trace.
- h93 (42) Has some meaning I cannot discover. It appears that this instruction needs a five-digit address. When you SST through it, it acts as a decimal point. Maybe it just hardens the display? The next five instructions are ignored if they are numbers, the same way as when you SST through a normal GTO instruction.
- h00 acts as List.
- h64 (13) acts as ABS(X) in Trace. Otherwise, it is all clear. It can be used as a Trace detector, but with fatal consequences.
- h84 (33) acts as RAD in Trace. Otherwise it acts as a strange RCL.
- h94 (43) enters EE 12. In Trace it acts as GRAD.
- h53 (2) is all clear. In Trace it is EE.
- h44 (93) acts as ln(0). In Trace it is TAN. This can be used as a nice trick to detect if the printer is in Trace or not. The only thing I cannot find is the exact way as HOW to use it.
- h75 (24) is DEG in Trace. Otherwise it acts as some kind of RCL. A Trace detector is simply h75 16. It jumps to LBL A' in Trace mode only.
- h85 (34) is RAD in Trace. Otherwise it also is a strange RCL.

It is a pity we don't have a Dr. W. Wickes in our club, (Dr. Wickes wrote THE book on synthetic programming for the HP-41C. Ed.) to write one titled "Synthetic Programming on the TI-59. After I read Dr. Wickes' book, the HP41C had no secrets for me anymore."



PERIODIC TABLE,- Lem Matteson of Kansas City, MO is the author of that program, which  
----- was published in v7n4/5p18 and 19.

The listing was done in the usual way, you know, TI-59 cum PC100A. Only Lem has this habit of blotting out unused zeros in GTO direct address instructions by means of a black pencil. Rather than trying to remove his pencil marks, I just used printers' white-out. This to satisfy numerous inquiries as to how those zeros had disappeared.

A few members told me that the program contained some typos, nothing big, though. Greg Hoen says that step 130 should be 4 instead of 5, because the author wanted to set flag 3 for numbers 1 to 4. With 5 there in step 130 the program sets flag 3 for numbers 1 through 5.

Additionally, step 141 should be 3 instead of 4. Also, atomic number 99 prints as O. Greg Hoen says that his chemistry reference material is ancient and that there could have easily been a name change. But it should be ES for Einsteinium and register 45 then should be -1736.254323288 vs -3256.2543289.

PROGRAM PROTECTION,- Over the years I have seen many card protection schemes. The most  
----- obvious one is to record your card with a -1 through -4. However, the protection flag can be bypassed by a variety of methods and the protection can be rather easily broken. Other schemes I have seen entail a password, a secret number that has to be matched to get access to the program. Simply listing the program will reveal, of course, the password. Therefore, in my opinion, it doesn't pay to try to protect a mag card.

On the other hand, it is rather simple to protect a CROM or module. Dave Leising sent me a few schemes he proposes:

1. Making any program longer than 960 steps, if needed by means of fake NOPs, will prevent it from being downloaded by means of OP 09. Very effective protection of individual programs in a module.

2. The last program in a CROM can be assigned a fictitiously high upper address in the internal CROM map, such that it cannot be downloaded, even though the program contains less than 960 steps.

3. Unauthorized execution of a protected module (CROM) can be prevented rather easily by means of the password scheme. Suppose the password is 93716. To get access to the entire CROM one would have to punch in: 93716 2nd PGM 15 A.

And at PGM 15 in the module we would have:

LBL A HIR 03 E'..... Now at LBL E' 93716 X:T HIR 13 EQ NNN 2nd SIN RTN

At NNN the program would continue if the correct password was entered. Otherwise the calculator would encounter the famous "2nd SIN" sequence and go to never-never land, from which a return is only possible by turning off the calculator.

E' calls could be distributed strategically along the entire module, preventing analysis by absolute address entry attempts.

The "2nd SIN" crash doesn't work on the TI-58, but I am sure there is a crash sequence that works on that calculator too.

HOW TO CHANGE YOUR TI-57 INTO A TI-57C. Dejan Ristanovic tells me that the French magazine "L'Ordinateur de Poche" says how to do it, in its No.4 issue of 1982. At the end of your program, step 48, place the following:  
2nd EXC SST LBL 1. Press LRN. Now, when you want to switch your calculator off, to save the memory, press GTO 2nd 48 R/S INV STO 3 +/- +/- . The display will go blank. Thus, the machine will not supply any more power to light the LED display. It uses even less energy than the TI-58C to save both data and program.  
When you want to switch the calculator back on, press INV FIX CLR.

52-NOTES,- Richard Vanderburgh, 9459 Taylorsville Road, Huber Heights OH, 45424, USA,  
----- the editor of that newsletter, now defunct, sends me the following information:  
34 Back issues of the 52-Notes are still available at \$ 1.50 each, \$ 2.00 if abroad, US funds. The first twelve issues cover mainly SR-52, SR-56 and PC100. The remaining 22 issues cover mainly TI-58, TI-59 and PC100C.

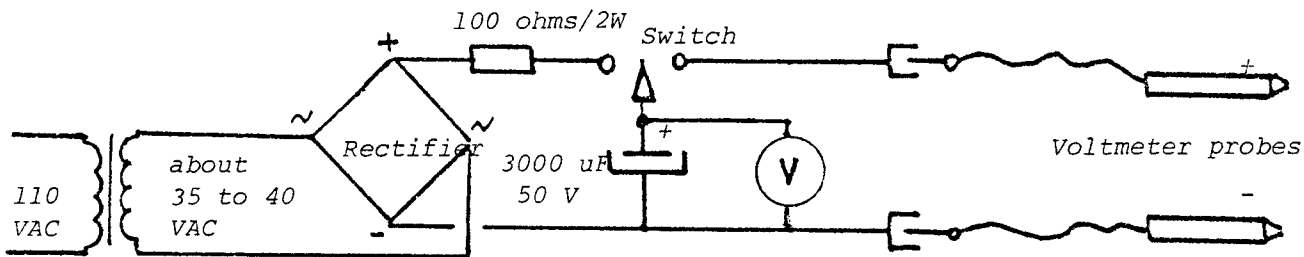
Note mainly the renaming of the "city" from "Dayton" to "Huber Heights" and the changed prices for each issue.

EMERGENCY POWER, Palmer Hanson Jr. writes: " In "Battery Power in an Emergency" (v6n8p14) ----- Bob Fruit discussed the use of a six volt lantern pack as an emergency power source for the TI-59. A more likely operation of the TI-59 with the power source outside the normal operating range occurs when one of the three cells of the BP-1A battery pack develops an internal short. Such a malfunction is typical of the nickel cadmium cells used in the battery pack. The result is that no amount of charging with either the model AC 9131 or with the charging circuit of the PC100 will eliminate the short. (The Zapper, described below, will. Ed.) The output of a "fully" charged battery pack becomes somewhat less than 3 volts. And with this low power output:

- When running with a battery pack only, the display will be somewhat dim. Calculations will be performed satisfactorily. Reading and writing of magnetic cards will be erratic at best, and often impossible. Unpredictable displays result when trying to read a card, for example, flashing two-digit numbers.
- With the defective battery pack installed in the PC100 charging receptacle, all reactions will be normal with the calculator installed on the PC100.
- With the defective battery pack installed and with the AC-9131 adaptor connected, the reactions will be the same as with the defective battery pack only.

Since originally indentifying the above malfunction I have helped several others with calculators which would read and write the magnetic cards erratically. Could this type of malfunction have been the source of "hundreds upon hundreds of misreads of my cards.." as reported in the "Mailbag" of v6n6/7p2 ? "

THE ZAPPER, Maurice E.T. Swinnen. In connection with the above, especially with the words ----- " when one of the three cells of the battery pack develops an internal short.." I had been annoyed several times, enough to start a "course of inquiry" as to the source of this annoyance. Dissecting a defective battery I found a well defined dendrite in its



insides. While the dendrite proved to be conductive, it definitely was identified as the villain who shorted my battery. A simple application of about 50 volts evaporated said dendrite. Thusly, the Zapper was born.

The schematic diagram above shows what is needed. A power supply, either one you put together from a discarded transformer and a bridge rectifier or an existing one, capable of supplying about 50 volts at a few milliamps will do. An electrolytic capacitor of at least 3000 uF at 50 volts has enough energy stored to zap even the most recalcitrant battery. A voltmeter (your VOM will do nicely if it has in the vicinity of 100,000 ohms per volt internal impedance so as not to load too heavily the charged capacitor) is handy to find out when your capacitor is charged up and ready for the next zap. A switch allows you to charge the capacitor slowly when switched to the left and to discharge the capacitor quickly across the NiCad cell when turned to the right. To this end, drill three small holes, about 2 mm, on either side of the battery pack, THROUGH THE PLASTIC ONLY, so as to expose the heads (metal) of each cell. Your voltmeter probes come in handy to apply "zapping power" through these small holes to each cell. Sometimes, a couple of applications are needed. It is good practice to zap a known shorted cell a couple of times, then to charge it for an hour or so in the PC100. If it still doesn't take a charge, zap again.

I have had success in about 90 % of all the cases, taken over more than forty batteries zapped, TI, HP and others.

PATENTS: A few more calculator-related patents have been received:

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 # 3,924,110 CALCULATOR SYSTEM FEATURING A SUBROUTINE REGISTER, Filed Sept. 13, 1973. Awarded Dec. 2, 1975 to Michael J. Cochran & Charles P. Grant.  
 # 3,916,169 CALCULATOR SYSTEM HAVING PRECHARGED VIRTUAL GROUND MEMORY, Filed Sept. 13, 1973. Awarded Oct. 28, 1975 to Michael J. Cochran & Charles P. Grant.  
 # 3,919,532 CALCULATOR SYSTEM HAVING AN EXCHANGE DATA MEMORY REGISTER, Filed Sept. 13, 1973. Awarded Nov. 11, 1975 to Michael J. Cochran & Charles P. Grant.  
 # 3,919,536 PRECHARGED DIGITAL ADDER AND CARRY CIRCUIT, Filed Sept. 13, 1973. Awarded Nov. 11, 1975 to Michael J. Cochran & Charles P. Grant.  
 # 3,922,538 CALCULATOR SYSTEM FEATURING RELATIVE PROGRAM MEMORY. Filed Sept. 13, 1973. Awarded Nov. 25, 1975 to Michael J. Cochran and Charles P. Grant.

It is curious to note, with respect to the last patent, that TI never used a relative memory in any of its calculators. But HP, on the other hand, has it incorporated in many of its programmables. Could it be that the first designs for the HP calculators were done at TI? I know that the very first ones were manufactured at TI.

For our newcomers: A relative memory means that you might have a command such as GTO RELATIVE. That would mean you would advance so many steps as the contents in a specific register would indicate. That would work from any program step. It would just as well work with a positive as with a negative number, meaning you would either go forward or backward relative to your present position. It is very easy to insert or delete steps in a RELATIVE ADDRESS mode as opposed to a DIRECT ADDRESS mode.

-----  
 PUZZLE,- Harald M. Otto of West Germany offers the following puzzle:

----- "Devise a routine that will invert the digits of a positive whole number comprising between 2 and 13 digits." Of course, only 10 digits will be shown in the display, but we now have enough routines to even print a 13-digit number and check if the reversing routine works correctly.

To make it absolutely unambiguous as to what is wanted, here are a few samples: 123456789 has to become 987654321 and 1122334455667 should become 7665544332211.

-----  
 FLASHTUBE CIRCUIT PROGRAM,- Eamon Murphy, Electro-Optical System Design, April 1982, pp. 41-45.

This TI-59/PC100 program will work primarily for linear flashtubes. It gives a simple design procedure to optimize single-mesh driving circuits. The program may be tuned to a particular manufacturer's tube by adjusting two values in the program.

All design equations are given plus a sample calculation. Author's address is Photodyne Inc. 948 Tourmaline Drive, Newbury Park, CA, 91320, USA.

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 PATENTS,- Still more patents arrived:

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 # 3,904,862 CALCULATOR SYSTEM HAVING A CONSTANT MEMORY, Filed Sept. 13, 1973, awarded Sept. 9, 1975 to Michael J. Cochran & Charles P. Grant.  
 # 3,904,863 CALCULATOR SYSTEM USING INSTRUCTION WORDS AS DATA, Filed Sept. 13, 1973, awarded Sept. 9, 1975 to Michael J. Cochran & Charles P. Grant.

-----  
 MODULE SELECTOR,- When ordering the module selector from Systems 7 in Houston TX at the greatly reduced price, as announced in v7n4/5p9, please be advised that this firm requires an extra \$ 4.00 US as handling and shipping charges per module switcher. Even at that, it is still considerably cheaper than directly from American Microproducts.

#### スピーク&スベルの特長

- 中学校で習う基礎英単語を、学習レベルに合わせて楽しく学べます。
- 人間の声に忠実な、電子音声による本格的な発音を聞きながら学べます。
- 英単語を使ったいろいろなゲームも楽しめます。



PROGRAMMING THE TI-55, by Stephen L. Snover and Mark A. Spikell, Prentice-Hall 1982,  
----- 116 pages, \$ 7.95 US.

Although I couldn't imagine why anyone would ever buy a TI-55 (\*) when for a few dollars more you can buy the so much more powerful TI-57, this book is well-written and could be a real help to, say, the high school student who just bought one of these programmables and finds out that the manual itself doesn't provide all of the answers.

Professors Stover and Spikell, old hands at this game, have milked the last drop out of this one too and have explained some of the function into such detail that it made me infortable. Could anyone be that dense to require such detail? But then, they are the educators, not I, and they probably know what they are doing.

The book contains a wealth of worked-out examples and procedures to teach somebody what programming is all about. Because the TI-55 lacks a decision-making test, the authors even devised one, based on the fact that if you count down a register in steps of one and use the contents each time as the divisor in a 1/x test, at zero it becomes undefined and the calculator stops with a flashing 999999. This sort of poor-man's DSZ (or should it rather be "decrement and stop on zero"?) is neat. It allows you to program finite loops now.

(\*) To justify my rather sour remark about the TI-55: it doesn't have decision-making test, nor does it have labels or even editing capabilities. Yes, if you forget to key in a step, you have to do it all over again. As a consolation, no program can be longer than 32 steps. A redeeming feature of the 55 is, that it has more special functions than any other calculator: conversions, statistics, even sinh, cosh and tanh.

If the TI-55 had come along in, say, 1972, it would have been a sensation. As it is, this is 1982 and we are getting jaded with respect to programmables. In my opinion, the TI-57 impacts rather heavily on the TI-55, with more programming capability at a slightly increased price.

HOW WELL DOES THAT LINE FIT, by George Wm. Thomson, Detroit, Michigan.

----- The built-in hard-wired functions of the TI-59 make the fitting of straight lines an easy task. It is assumed that the X-values are free from error and that the data points are of equal weight. The Method of Least Squares is used to minimize the sum of the squares of the deviations of the observed values from the calculated line. The only measure of "goodness of fit" of the points which is provided is the Correlation Coefficient (OP 13) which is a very poor kind of measure for most engineering and scientific work. Fortunately the much more useful Standard Error of FIT, sometimes called the Mean Square Error of Fit, can be obtained from the short subroutine below:

LBL A 1 - OP 13  $X^2 = X$  OP 11  $X$  RCL 03 = DIV ( RCL 03 - 2 =  $\sqrt{X}$  PRT RTN

The result is in the same units as the Y-values with n-2 degrees of freedom. The SE FIT is obtained by adding up the squares of the deviations of the calculated values of Y from the input values, dividing the sum by n-2 and taking the square root. Since none of the stored data in registers 01 through 06 is altered either by OP 12 to 15 or the above subroutine, adding or dropping data points from a set is easily done without reentering the whole set.

To use the subroutine, press A after at least three data points have been entered. The SE FIT will be printed. Points can then be added using the X:T and  $\Sigma^+$  keys.

The subroutine permits evaluation of the set of data at any stage.

This subroutine is almost identical with a routine I derived for the SR-51-II several years ago.

GAUSS REDUCTION OF MATRICES  $M \cdot N$ , - Saša Nick and Dejan Ristanović, the "Worthington & Regelman"'s of Belgrade, Yugoslavia have done it again. This time it is a mathematical program of class A-1. The program can solve any (and we mean ANY) system of  $m$  linear equations with  $n$  unknowns ( $m \cdot (n+1) < 85$ ). If you are familiar with linear algebra, you can use it to find the kernel ( $\ker f$ ) and the image ( $\text{Im } f$ ) of any linear transformation, and so on. The program is very short (only 240 steps), will therefore fit on one side of a mag card, and will accomodate the unusual number of  $m$  linear equations with  $n$  unknowns of  $m \cdot (n+1) = 84$ . The program could be made more user-friendly, but that would entail adding about 80 steps, which would reduce the size of the systems it being able to solve. The authors claim that they would be able to write such a "user-friendly" program, if members' interest warrants it.

The system is generally expressed as:

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n &= b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n &= b_2 \\ a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n &= b_m \end{aligned}$$

To solve such a system, you have to write it as a matrix  $m \cdot (n+1)$ . You now simply enter the coefficients of the matrix and start the program. After a short interval of computation you will receive another matrix that looks like:

$$\left| \begin{array}{cccc|c} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ a_{21} & a_{22} & \dots & a_{2n} & b_2 \\ \vdots & \vdots & & \vdots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} & b_m \end{array} \right| \sim \left| \begin{array}{cccc|c} 1 & a'_{12} & a'_{13} & \dots & a'_{1n} & b'_1 \\ 0 & 1 & a'_{23} & \dots & a'_{2n} & b'_2 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & 0 & \dots & a'_{mn} & b'_m \end{array} \right|$$

Once you have the second matrix it will be easy to solve the system. If there are more unknowns than equations, don't expect to find their exact values: some of the unknowns will be a function of other unknowns. If there are more equations than there are unknowns it is not entirely sure that the system has a solution at all. To check on this, you might solve the  $m \cdot n$  system and manually check if the other equations (there are  $m-n$  of them) are good. You can also use this program by imagining more unknowns ( $m-n$  new ones) and entering their coefficients as zero. Then reduce the matrix and find out if there is a line that looks like  $0x_j=5$ . If there is such a line, the system doesn't have a solution.

User instructions: Key in the program and save it on one side of a mag card in 6 OP 17 partitioning. To use it, enter  $n+1$  (number of columns) and press X:T. Enter the number of rows ( $m$ ) and press A. Enter  $a_{11}$  and press R/S. Enter  $a_{12} \dots a_{1n}$  and press R/S each time. Continue with this until you have entered all the coefficients.

Press B and wait for the computation to finish. The display will show  $a'_{11}$ . Press R/S and the display will show  $a'_{12}$  and so on. It is very important to realize that the calculator will NOT display all the elements. It will only display the triangle matrix. Thus, after  $a'_{1n}$  and  $b'_1$  the display will show  $a'_{22}$ . ( $a_{21}$  is always zero. The knowledge of this will save you some time, obviously)

For a new case, enter  $(n+1)$ , press X:T, and so on.

Operating limits and warnings: The calculator MUST NOT be in the ENG mode. You are responsible for checking if the number of rows and columns is alright. Otherwise, the calculator will flash after 85 coefficients have been entered.

The program cannot solve systems with some of the coefficients very large, say 10000000000 while others very small, say 0.0000000001. Also if all of the coefficients are too large (or too small) the authors suggest you transform the equations manually before you use the program.

See program on next page, please.

THIRTEEN DIGIT MODULO 30 SPEEDY FACTOR FINDER.- Palmer O. Hanson Jr. has completed the documentation on it and submitted it to PPX. It uses STF-h12 methods to enter into Fast Mode. Write to Texas Instruments Inc. PPX, P.O. BOX 53, Lubbock TX 79408 for a copy (\$ 4.00 + \$ 2.00 handling + tax) of this 17-page program.

GAUSS REDUCTION OF MATRICES M·N, Sasa Nick & Dejan Ristanović.

000	76	LBL	040	20	20	080	01	01	120	10	E'	160	16	16	200	75	-
001	10	E'	041	43	RCL	081	80	80	121	43	RCL	161	95	=	201	43	RCL
002	42	STD	042	00	00	082	19	D'	122	01	01	162	32	X:T	202	01	01
003	03	03	043	42	STD	083	43	RCL	123	22	INV	163	43	RCL	203	95	=
004	75	-	044	01	01	084	00	00	124	64	PI*	164	01	01	204	55	+
005	01	1	045	42	STD	085	10	E'	125	05	05	165	10	E'	205	82	HIR
006	95	=	046	04	04	086	73	RC*	126	69	DP	166	32	X:T	206	17	17
007	65	x	047	10	E'	087	05	05	127	35	35	167	22	INV	207	65	x
008	82	HIR	048	25	CLR	088	32	X:T	128	97	DSZ	168	74	SM*	208	53	(
009	17	17	049	73	RC*	089	43	RCL	129	02	02	169	05	05	209	22	INV
010	85	+	050	05	05	090	01	01	130	01	01	170	69	DP	210	59	INT
011	43	RCL	051	52	EE	091	10	E'	131	21	21	171	34	34	211	69	DP
012	04	04	052	22	INV	092	32	X:T	132	82	HIR	172	97	DSZ	212	10	10
013	85	+	053	67	EQ	093	63	EX*	133	18	18	173	02	02	213	85	+
014	05	5	054	01	01	094	05	05	134	42	STD	174	01	01	214	01	1
015	95	=	055	08	08	095	32	X:T	135	01	01	175	53	53	215	95	=
016	42	STD	056	29	CP	096	43	RCL	136	43	RCL	176	97	DSZ	216	22	INV
017	05	05	057	25	CLR	097	00	00	137	01	01	177	01	01	217	44	SUM
018	92	RTN	058	73	RC*	098	10	E'	138	32	X:T	178	01	01	218	01	01
019	76	LBL	059	05	05	099	32	X:T	139	43	RCL	179	36	36	219	61	GTD
020	19	D'	060	52	EE	100	72	ST*	140	00	00	180	82	HIR	220	01	01
021	82	HIR	061	22	INV	101	05	05	141	67	EQ	181	18	18	221	94	94
022	17	17	062	67	EQ	102	69	DP	142	01	01	182	32	X:T	222	76	LBL
023	42	STD	063	00	00	103	34	34	143	76	76	183	43	RCL	223	11	R
024	04	04	064	82	82	104	97	DSZ	144	42	STD	184	00	00	224	37	P/R
025	75	-	065	82	HIR	105	02	02	145	04	04	185	22	INV	225	09	9
026	43	RCL	066	17	17	106	00	00	146	32	X:T	186	77	GE	226	69	DP
027	00	00	067	44	SUM	107	83	83	147	10	E'	187	12	B	227	17	17
028	95	=	068	05	05	108	43	RCL	148	73	RC*	188	22	INV	228	47	CMS
029	42	STD	069	69	DP	109	00	00	149	05	05	189	58	FIX	229	06	6
030	02	02	070	21	21	110	42	STD	150	82	HIR	190	06	6	230	48	EXC
031	69	DP	071	43	RCL	111	04	04	151	06	06	191	42	STD	231	01	01
032	22	22	072	01	01	112	10	E'	152	19	D'	192	01	01	232	91	R/S
033	92	RTN	073	32	X:T	113	73	RC*	153	43	RCL	193	25	CLR	233	72	ST*
034	76	LBL	074	82	HIR	114	05	05	154	00	00	194	73	RC*	234	01	01
035	12	B	075	18	18	115	42	STD	155	10	E'	195	01	01	235	69	DP
036	58	FIX	076	77	GE	116	01	01	156	73	RC*	196	91	R/S	236	21	21
037	08	08	077	00	00	117	19	D'	157	05	05	197	69	DP	237	61	GTD
038	29	CP	078	56	56	118	43	RCL	158	65	x	198	21	21	238	02	02
039	69	DP	079	61	GTD	119	03	03	159	82	HIR	199	06	6	239	32	32

Texas Instruments Inc.- You might have heard rumors or seen newspaper articles telling you about financial troubles at TI. It simply concerns an internal reorganization which implies a lay-off of about 6000 people. But, on the other hand, TI is actively advertizing jobs in the Dallas area for specialists in certain IC operations.

Recently, Forbes Magazine, in its March 15, 1982 issue had an article called "Texas Instruments at mid-life", by Anne Bagamery. Ms. Bagamery didn't have all of her facts straight, though. At one point she writes: "For example, TI is not increasing its commitments to the consumer products division - \$ 19.95 watches and \$ 17.50 calculators mainly - but is phasing the division out, with the exception of personal computers and its Speak and Spell learning aids, which embody a technological advance of tremendous promise."

Now, J. Fred Bucy, TI's president, just a week before announced that they were increasing their commitment to the consumer division. When shown the article in question by one of our members, who is also a large stockholder in TI, Mr. Bucy said: "We are continuing development on advanced programmable calculators and will soon have products to announce. I am sure you can appreciate my not being able to be more explicit until these developments are ready for general announcement. We continue to be the largest supplier of calculators to the U.S. market in both unit and dollar volume, and it is our intention to remain the leader in this area."

ZIMMERMAN'S PUZZLE: In v7n4p12 we presented this puzzle. Please refer to that issue for a complete explanation as to what was required.

The first solution I received when the newsletter was barely a few days in the mail came from Paul Blair in Milan, Illinois. His solution is short and sweet. It is also fast, although I didn't want to make this a speed contest. As is usual though, speed, compactness and elegance of solution are three factors that get you points.

000 76 LBL	005 53 <	010 54 >	015 55 +	020 94 +/-	025 59 INT
001 11 R	006 24 CE	011 59 INT	016 05 S	021 85 +	026 65 *
002 85 +	007 55 +	012 55 +	017 95 =	022 01 1	027 01 1
003 05 S	008 01 1	013 02 2	018 22 INV	023 95 =	028 00 0
004 85 +	009 00 0	014 95 =	019 59 INT	024 22 INV	029 95 =
				---	030 91 R/S

Then the flood gates suddenly opened and I got heaps of solutions, mostly all good ones. Peter Brown in Denville, New Jersey sent me this one:

000 76 LBL	010 12 B	020 00 0	030 75 -	040 76 LBL	045 14 D
001 11 R	011 55 +	021 95 =	031 01 1	041 13 C	050 92 RTN
002 29 CP	012 01 1	022 22 INV	032 00 0	042 85 +	051 76 LBL
003 42 STD	013 00 0	023 44 SUM	033 54 >	043 43 RCL	052 14 D
004 00 00	014 95 =	024 00 00	034 22 INV	044 01 01	053 85 +
005 22 INV	015 59 INT	025 53 <	035 77 GE	045 95 =	054 01 1
006 67 EQ	016 42 STD	026 02 2	036 13 C	046 94 +/-	055 00 0
007 12 B	017 01 01	027 65 *	037 75 -	047 22 INV	056 95 =
008 92 RTN	018 65 *	028 43 RCL	038 09 9	048 77 GE	057 92 RTN
009 76 LBL	019 01 1	029 00 00	039 95 =		

Then I received one from John Adams in Pittsburg, Pennsylvania. John says that he observed the pattern for the input numbers 00 through 09 and the fact that the single digit numbers decrease by 1 as you move from left to right enabled him to write the needed formulas to accomplish the task: When B=0 then the single digit= 0-A; when B=1,2,3,4 then the digit =10-2B-A; when B=5,6,7,8,9 then the digit = 19-2B-A; if the resulting digit is negative then add 10 to the result.

000 76 LBL	010 00 0	020 95 =	029 85 +	038 13 B	047 29 CP
001 11 R	011 95 =	021 32 X/T	030 01 1	039 09 9	048 77 GE
002 59 INT	012 75 -	022 05 S	031 00 0	040 35 =	049 14 D
003 42 STD	013 59 INT	023 32 X/T	032 95 =	041 76 LBL	050 85 +
004 01 01	014 42 STD	024 67 EQ	033 29 CP	042 13 C	051 01 1
005 29 CP	015 02 02	025 12 B	034 77 GE	043 75 -	052 00 0
006 67 EQ	016 95 =	026 65 *	035 13 C	044 43 RCL	053 95 =
007 14 D	017 65 *	027 02 2	036 85 +	045 02 02	054 76 LBL
008 55 +	018 01 1	028 94 +/-	037 76 LBL	046 95 =	055 14 D
009 01 1	019 00 0				056 92 RTN

The same day I received one from Clyde Durbin in Dallas, Texas. Clyde says he spent only one half hour on it, so if someone wants to refine it, that is OK. It works and that is that. Clyde also says that one could stack all the control digits in one register. I think unpacking would be rather slow, although the program could be made very efficient.

000 76 LBL	010 01 1	020 00 0	029 87 IFF	039 59 INT	049 29 CP
001 11 R	011 00 0	021 94 +/-	030 32 X/T	040 65 *	050 77 GE
002 42 STD	012 95 =	022 95 =	031 43 RCL	041 01 1	051 88 DMS
003 00 00	013 22 INV	023 77 GE	032 00 00	042 00 0	052 85 +
004 29 CP	014 59 INT	024 87 IFF	033 55 +	043 95 =	053 01 1
005 67 EQ	015 75 -	025 85 +	034 01 1	044 59 INT	054 00 0
006 88 DMS	016 01 1	026 09 9	035 00 0	045 32 X/T	055 95 =
007 55 +	017 95 =	027 95 =	036 95 =	046 75 -	056 76 LBL
008 06 6	018 65 *	028 76 LBL	037 85 +	047 32 X/T	057 88 DMS
009 32 X/T	019 01 1		038 22 INV	048 95 =	058 91 R/S

The reason I publish ALL these solutions is not so much that you should USE them, but rather that we all can learn something from the different approaches different programmers take. Programming is NOT a science governed by rigid rules, but rather an art in which every player takes part according to his/her own emotions.

Sid Hack in Columbia, South Carolina offers this one. Sid says that absolute addressing could speed this one up:

000	76	LBL	018	43	RCL	036	02	2	053	05	5	070	75	-	087	95	=
001	11	R	019	01	01	037	00	0	054	95	=	071	02	2	088	22	INV
002	47	CM8	020	75	-	038	95	=	055	42	STD	072	65	X	089	59	INT
003	42	STD	021	05	5	039	22	INV	056	04	04	073	53	(	090	65	X
004	01	01	022	95	=	040	59	INT	057	43	RCL	074	43	RCL	091	01	1
005	01	1	023	55	+	041	65	X	058	03	03	075	02	02	092	00	0
006	00	0	024	02	2	042	02	2	059	22	INV	076	85	+	093	95	=
007	32	XIT	025	00	0	043	00	0	060	77	GE	077	43	RCL	094	91	R/S
008	43	RCL	026	95	=	044	95	=	061	39	SIN	078	04	04	095	76	LBL
009	01	01	027	59	INT	045	42	STD	062	08	8	079	95	=	096	30	TAN
010	65	X	028	42	STD	046	03	03	063	61	GTD	080	85	+	097	01	1
011	02	2	029	02	02	047	55	+	064	39	CD8	081	01	1	098	00	0
012	95	=	030	43	RCL	048	05	5	065	76	LBL	082	00	0	099	00	0
013	22	INV	031	01	01	049	95	=	066	38	SIN	083	95	=	100	44	SUM
014	77	GE	032	75	-	050	22	INV	067	02	2	084	55	+	101	01	01
015	30	TAN	033	05	5	051	59	INT	068	76	LBL	085	01	1	102	61	GTD
016	76	LBL	034	95	=	052	65	X	069	39	CD8	086	00	0	103	12	B
017	12	B	035	55	+												

The next one comes from John Huntington Lewis in Norfolk, Virginia. Hunt says that the problem lies in recognizing that there are two orderly patterns in the digit generation table: a rotating sequence in the rows and another rotating -even,odd- in the columns. Observation of any row shows that the control digit decrements by one each time the tens digit of the key number increases. The algorithm then is "CONTROL DIGIT = BASE + OFFSET." Note the unique way Hunt has provided for the output: The input number is preserved and is printed together with the control digit as a decimal. If that is not wanted just omit LBL X<sup>2</sup> and write LBL X<sup>2</sup> R/S.

000	76	LBL	015	02	02	030	77	GE	045	09	9	060	54	)	075	39	CP
001	11	R	016	43	RCL	031	38	SIN	046	75	-	061	77	GE	076	76	LBL
002	58	FIX	017	01	01	032	65	X	047	53	(	062	30	TAN	077	33	X <sup>2</sup>
003	01	01	018	22	INV	033	02	2	048	02	2	063	85	+	078	32	XIT
004	29	CP	019	59	INT	034	54	)	049	65	X	064	01	1	079	55	+
005	42	STD	020	65	X	035	94	+/-	050	43	RCL	065	00	0	080	01	1
006	00	00	021	01	1	036	85	+	051	01	01	066	54	)	081	00	0
007	55	+	022	00	0	037	01	1	052	54	)	067	76	LBL	082	54	)
008	01	1	023	54	)	038	00	0	053	54	)	068	30	TAN	083	85	+
009	00	0	024	42	STD	039	54	)	054	76	LBL	069	32	XIT	084	43	RCL
010	54	)	025	01	01	040	61	GTD	055	39	CD8	070	01	1	085	00	00
011	42	STD	026	05	5	041	39	CD8	056	29	CP	071	00	0	086	95	=
012	01	01	027	32	XIT	042	76	LBL	057	75	-	072	22	INV	087	99	PRT
013	59	INT	028	43	RCL	043	38	SIN	058	43	RCL	073	67	EQ	088	91	R/S
014	42	STD	029	01	01	044	01	1	059	02	02	074	33	X <sup>2</sup>			

Then Lester Tibbetts from Emporium, Pennsylvania came through with the following one. He described the development of his algorithm almost in exactly the same words as Hunt Lewis did above, although he was able to do it in about half as many steps.

000	76	LBL	007	00	0	014	02	2	021	77	GE	038	43	RCL	035	01	1
001	11	R	008	75	-	015	00	0	022	34	FX	029	01	01	036	00	0
002	29	CP	009	59	INT	016	75	-	023	09	9	030	95	=	037	95	=
003	67	EQ	010	42	STD	017	01	1	024	94	+/-	031	94	+/-	038	76	LBL
004	33	X <sup>2</sup>	011	01	01	018	00	0	025	85	+	032	77	GE	039	33	X <sup>2</sup>
005	55	+	012	95	=	019	85	+	026	76	LBL	033	33	X <sup>2</sup>	040	91	R/S
006	01	1	013	65	X	020	22	INV	027	34	FX	034	85	+			



000 76 LBL	007 54 )	014 00 00	021 59 INT	028 53 (	035 85 +
001 11 A	008 42 STD	015 85 +	022 65 *	029 43 RCL	036 77 GE
002 53 (	009 00 00	016 02 2	023 09 9	030 00 00	037 00 00
003 29 CP	010 53 (	017 49 PRD	024 75 -	031 75 -	038 40 40
004 55 +	011 59 INT	018 00 00	025 01 1	032 69 DP	039 01 1
005 01 1	012 94 +/-	019 43 RCL	026 00 0	033 10 10	040 00 0
006 00 0	013 44 SUM	020 00 00	027 65 *	034 54 )	041 54 )
					042 92 RTN

[illegible]

000	76	LBL	006	82	HIR	012	00	0	018	22	INV	024	59	INT	030	59	INT
001	11	A	007	04	04	013	93	.	019	59	INT	025	55	÷	031	65	×
002	55	÷	008	59	INT	014	09	9	020	65	×	026	01	1	032	01	1
003	01	1	009	94	+/-	015	75	-	021	02	2	027	00	0	033	00	0
004	00	0	010	85	+	016	82	HIR	022	02	2	028	95	=	034	95	=
005	95	=	011	03	3	017	14	14	023	95	=	029	22	INV	035	92	RTN

[illegible]

Nello Coda first sent this "conventional" solution. That is, one that conforms to the definition "computed" as stated in the original posing of the problem.

000 76 LBL	012 76 LBL	024 65 ×	036 11 11	048 59 INT	060 02 2
001 38 SIN	013 39 CDS	025 01 1	037 59 INT	049 65 ×	061 00 0
002 55 ÷	014 85 +	026 00 0	038 94 +/-	050 01 1	062 85 +
003 01 1	015 43 RCL	027 95 =	039 44 SUM	051 02 2	063 02 2
004 00 0	016 11 11	028 91 R/S	040 11 11	052 32 X:T	064 00 0
005 95 =	017 95 =	029 76 LBL	041 55 ÷	053 01 1	065 95 =
006 22 INV	018 55 ÷	030 11 A	042 01 1	054 00 0	066 77 GE
007 59 INT	019 01 1	031 55 ÷	043 00 0	055 95 =	067 38 SIN
008 65 ×	020 00 0	032 01 1	044 85 +	056 48 EXC	068 75 -
009 01 1	021 95 =	033 00 0	045 01 1	057 11 11	069 01 1
010 00 0	022 22 INV	034 95 =	046 95 =	058 94 +/-	070 61 GTD
011 95 =	023 59 INT	035 42 STD	047 22 INV	059 65 ×	071 39 CDS

But, as Nello points out, this one is only a special case, which is "computable." Suppose your control digits were arranged in such a random order that it would be virtually impossible to compute a table.

Now, before I hear a chorus of howls go up accompanied by cries of "foul" and "this does not fit the definition" let me entertain you with a short story:

My job entails "problem solving". I have been doing this for the last thirty years and I am slowly getting the hang of it. Now, the medical researchers I work for and with are supposed to define the problems for me so I can work them out. Invariably I get either a problem with severe restrictions, special cases or partial solutions. When I was still young and inexperienced I would fall for these ploys and give the researchers exactly what they wanted, a special case solution. After about one month the instrument I built for them would be returned with the plea "can you modify it for the following conditions." If I did, it would return after another two months, for more modifications. I once drew a picture of my own face, thick drops of bloody sweat dripping from it, with the caption NO MORE MODIFICATIONS, PLEASE. Thus, I learned that it is much better to give a GENERAL solution, EVEN IF ONLY A PARTIAL SOLUTION WAS REQUESTED.

User instructions for Nello's general solution: Put the calculator in 10 OP 17 and load the ten registers with the control digits as shown. Key in the program and record everything in 6 OP 17 on side 1 of a mag card. Enter the two-digit number and press A. You see now that by loading the control digits in a different order in the regs, any random sequence can be accomodated. It will also be rather simple to relocate this routine and even use different registers. You might have to change a few constants in the routine itself, such as the 90 in steps 006-007.

135792468.	90	000 76 LBL	010 11 11	020 95 =	030 65 ×
24681357.9	91	001 11 A	011 22 INV	021 22 INV	031 06 6
913570246.8	92	002 55 ÷	012 59 INT	022 28 LOG	032 69 DP
802469135.7	93	003 01 1	013 65 ×	023 35 1/X	033 17 17
791358024.6	94	004 00 0	014 01 1	024 65 ×	034 01 1
680247913.5	95	005 85 +	015 00 0	025 73 RC*	035 00 0
579136802.4	96	006 09 9	016 69 DP	026 11 11	036 95 =
468025791.3	97	007 00 0	017 17 17	027 95 =	037 59 INT
357914680.2	98	008 95 =	018 01 1	028 22 INV	038 91 R/S
246803579.1	99	009 42 STD	019 00 0	029 59 INT	

And this is where I stop for this issue. Not that I ran out of submissions: at the time of this writing, May 25th, I have received an additional 18 solutions. This is by far the most successful "utmaning" we ever ran in the TI PPC NOTES. If some of the other solutions prove to be really different from the ones published in this issue, I will put them in next issue. Otherwise, I will consider this poney beaten to death. To the many contributors: thanks. Please find Harald M. Otto's new puzzle somewhere in this issue. You will know what to do with. The best of luck.